

# Tutorial: The Correctness-by-Construction Approach to Programming Using CorC

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**Abstract**—The Correctness-by-Construction tutorial focuses on a structured programming approach for correct software development. Besides functional correctness, also non-functional properties such as security properties can be guaranteed using the CbC approach. In this tutorial, the participants learn a good practice to develop software that is midway between formal approaches and a “hack into correctness” style.

**Index Terms**—tutorial, formal methods, correctness-by-construction

## I. INTRODUCTION

The purpose of the tutorial is to influence the way participants approach the task of developing algorithms, with a view to improving code quality. The tutorial features: a step-by-step explanation of how to derive provably correct algorithms using small and tractable refinements rules; a detailed illustration of the methodology through a set of carefully selected examples of increasing complexity; a demonstration of how practical non-trivial algorithms have been derived. The focus is on bridging the gap between two extreme methods for developing software. On the one hand, some approaches are so formal that they scare off all, but the most dedicated theoretical computer scientists. On the other, there are some who believe that any measure of formality is a waste of time, resulting in software that is developed by following gut feelings and intuitions.

## II. CORRECTNESS-BY-CONSTRUCTION

Correctness-by-Construction (CbC) [KW12], [Mor94] is an approach to incrementally create formally correct programs guided by pre- and postcondition specifications. A program is created using refinement rules that guarantee the resulting implementation is correct with respect to the specification.

The CbC approach to program development begins with a Hoare triple comprising a precondition, an abstract statement, and a postcondition. Such a triple should be read as a total correctness assertion, if the precondition holds, and its abstract statement executes then the execution will terminate, and its postcondition will hold. This triple can be refined by using a set of refinement rules, i.e., the statement is replaced by more concrete statements. For example, a loop is introduced, or an abstract statement is replaced by an assignment. If no abstract statement remains, the code is fully specialized.

## III. CBC BY EXAMPLE

The tool CorC [RSC<sup>+</sup>19] is a hybrid textual and graphical IDE for the development of functional correct programs using CbC. In the tool, a starting Hoare triple specification  $\{P\}S\{Q\}$  with a precondition  $P$ , a postcondition  $Q$ , and an abstract program  $S$  is refined stepwise to a correct implementation. The refinement is done by applying refinement rules which guarantee that the refined program still satisfies its starting specification. CorC checks the side-conditions of each applied refinement automatically such that the programmer gets direct feedback whether a refinement cannot be proven.

In Fig. 1, we construct a linear search algorithm. Each node represents one Hoare triple. An arrow is the application of a refinement, connecting a Hoare triple with the refined Hoare triple. The linear search algorithm has a precondition  $P := \text{appears}(a, x, 0, a.length)$ . This predicate states that the element  $x$  appears in the array  $a$  between the boundaries 0 and the array length. The postcondition  $Q := \text{modifiable}(i); a[i] = x$  states that the element at index  $i$  in array  $a$  is equal to the searched element  $x$ . With `modifiable` we specify that only the variable  $i$  can be altered. In CorC, the used parameters and local variables are declared in a node on the top right. Below this node, global conditions are expressed. These conditions (e.g., invariants) have to be true in each step of the program.

To find the element in a linear search, we traverse the array from back to front. As invariant, we specify  $\text{appears}(a, x, i + 1, a.length)$ . We split the array in two parts: the part which we already examined and where  $x$  is not found; and the part that should still be examined. The program is constructed with four refinement steps. First, the composition rule [KW12] splits the starting Hoare triple into two triples with an intermediate condition which is equal to our invariant. The first statement *statement1* is refined to an assignment to start at the end of the array. The second statement *statement2* is refined to a loop. We repeat the loop until we found  $x$ . The loop body is refined in the last step. Here, we decrease the variable  $i$ . Each refinement is verified automatically by CorC to prove that the refined program satisfies the starting specification. For example, in the last refinement step, we prove that the invariant is preserved in

