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Computation Rate Maximization for Wireless-Powered Edge Computing With Multi-User Cooperation

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ABSTRACT The combination of mobile edge computing (MEC) and radio frequency-based wireless power transfer (WPT) presents a promising technique for providing sustainable energy supply and computing services at the network edge. This study considers a wireless-powered mobile edge computing system that includes a hybrid access point (HAP) equipped with a computing unit and multiple Internet of Things (IoT) devices. In particular, we propose a novel multi-user cooperation scheme to improve computation performance, where collaborative clusters are dynamically formed. Each collaborative cluster comprises a source device (SD) and an auxiliary device (AD), where the SD can partition the computation task into various segments for local processing, offloading to the HAP, and remote execution by the AD with the assistance of the HAP. Specifically, we aims to maximize the weighted sum computation rate (WSCR) of all the IoT devices in the network. This involves jointly optimizing collaboration, time and data allocation among multiple IoT devices and the HAP, while considering the energy causality property and the minimum data processing requirement of each device. Initially, an optimization algorithm based on the interior-point method is designed for time and data allocation. Subsequently, a priority-based iterative algorithm is developed to search for a near-optimal solution to the multi-user collaboration scheme. Finally, a deep learning-based approach is devised to further accelerate the algorithm's operation, building upon the initial two algorithms. Simulation results show that the performance of the proposed algorithms is comparable to that of the exhaustive search method, and the deep learning-based algorithm significantly reduces the execution time of the algorithm.

INDEX TERMS Mobile edge computing, multi-user collaboration, wireless power transfer, time and data allocation, deep learning.

I. INTRODUCTION

A. BACKGROUND

THE INDUSTRIAL Internet of Things (IIoT) is an application of Internet of Things (IoT) in the industrial sector [1]. It represents a new stage for the development of intelligence in traditional industries. IoT nodes with heterogeneous sensors are deployed in an IIoT network to perform various monitoring and control functions. The data collected by IoT nodes can be analyzed to enhance production effectiveness and reliability [2]. Generally, data is

collected through IoT nodes' access points and transmitted to a cloud computing center for processing and analysis. However, cloud computing introduces long transmission delays, rendering it unsuitable for real-time IIoT scenarios. Moreover, IIoT systems place a strong emphasis on the protection of private data, especially when dealing with confidential products that may involve sensitive information [3].

To enhance data processing and privacy protection in IIoT systems, empowering access points with mobile edge computing (MEC) capability is under consideration [4], [5].

© 2024 The Authors. This work is licensed under a Creative Commons Attribution-NonCommercial-NoDerivatives 4.0 License. For more information, see https://creativecommons.org/licenses/bv-nc-nd/4.0/ MEC enables low-latency operations by extending cloud computing and services to the network edge [6], [7], [8], [9], [10]. It also enhances the ability to protect private data, as data processing occurs exclusively within the plant's network. However, the computing capabilities of the access point and IoT nodes are limited. Additionally, data collected by different IoT nodes has varying levels of importance. Therefore, we aim to analyze as much high-importance data as possible while ensuring that each IoT node meets the minimum data processing requirement. This goal enables personalized and targeted data analysis in the IIoT system.

In the context of MEC, the literature extensively investigates two popular offloading policies [11], [12], [13], [14], [15]: partial offloading and binary offloading. The former is applicable to arbitrarily divisible computing tasks. In this case, each task is divided into two parts: one is offloaded to the MEC server, and the other is computed locally by the IoT device. The latter is suitable for indivisible computing tasks, such as simple atomic tasks. Therefore, each task is computed either by the IoT node or the MEC server. However, in the scenario we consider, these two strategies fail to harness the IIoT system's computational capacity fully. This study explores collaboration among multiple IoT nodes and the hybrid access point (HAP) to maximize the weighted sum computation rate (WSCR) of the IIoT system.

On the other hand, due to the stringent device size constraint and production cost consideration, an IoT device often carries a capacity-limited battery, which needs frequent battery replacement and charging. However, this requirement can result in significant labor costs and financial burdens for service providers, particularly in large-scale IIoT systems and harsh environments [16], [17]. Therefore, radio frequency (RF)-based wireless power transfer (WPT) has emerged as an effective solution to the problem of finite battery capacity [18]. Specifically, WPT employs a dedicated RF energy transmitter to charge the batteries of remote energyharvesting devices continuously [19]. According to reports, commercial WPT transmitters can continuously supply RF power of approximately 5W over distances exceeding 10 meters [20]. Deploying the information transceiver, energy transmitter, and computing unit on the HAP to provide both MEC services and WPT can yield their benefits simultaneously. Therefore, wireless-powered mobile edge computing has garnered significant interest as a viable solution to address computing capacity constraints and energy consumption issues in IIoT systems.

B. RELATED WORK

In the existing literature [11], [13], [16], [19], [21], [22], [24], [25], [26], [27], [28], [29], MEC and WPT have been extensively investigated in IoT systems. In [13], an MEC network was studied to minimize the execution delay of total tasks. The study proposed an optimized offloading framework that allows multiple IoT devices to offload tasks to multiple edge computing servers. In [29], the authors focused on minimizing the overall energy consumption of

IoT devices. The research investigated resource allocation techniques for multi-user MEC systems based on time division multiple access (TDMA) and orthogonal frequency division multiple access (OFDMA). However, the WPT was not involved in [13] and [29].

The optimal design of wireless-powered MEC networks differs significantly from traditional non-WPT-involved MEC networks due to the trade-off between WPT and wireless information transfer (WIT) [22]. To effectively leverage the benefits of MEC and RF-based WPT, several recent studies have delved into the investigation of wireless-powered MEC networks. In [19], a single-user wireless-powered MEC system was studied to minimize the total transmitted energy consumption of the energy transmitter (ET). The study developed heuristic online designs for joint energy and task allocation. In [21], the authors presented analytical expressions for the probability of successful computation and the average number of successfully computed bits in a singleuser wireless-powered MEC system, aiming to maximize the system's utility. However, [19] and [21] only considered the single-user network model and the derived results may not be readily applicable to multi-user scenarios.

In practice, IIoT systems commonly consist of a large number of IoT devices. As a result, several studies have explored multi-user wireless-powered MEC networks. However, most of these works primarily focus on the tradeoff between computing offloading and wireless charging, neglecting the aspect of collaboration among users. The study in [23] investigated a multi-user MEC network powered by WPT. In this network, each energy-harvesting node adheres to a binary computation offloading policy. A joint optimization method, utilizing the alternating direction method of multipliers (ADMM) decomposition technique, was proposed. This method aims to maximize the weighted sum computation rate of all nodes in the network by optimizing both individual computing mode selection and system transmission time allocation jointly. Work in [22] studied a wireless-powered multi-user MEC network in which the transmit power of HAP was minimized. The authors devised an algorithm based on the ADMM to optimize the time allocation and computational mode selection jointly. In [16], a wireless-powered multi-user hierarchical fog cloud computing network was studied to maximize the minimum energy balance among all users. To achieve the global optimal solution, the authors proposed a generalized Benders decomposition (GBD)-based approach. In [24], a rotary-wing unmanned aerial vehicle (UAV)-assisted multi-user wirelesspowered MEC network was studied, aiming to minimize the energy consumption of the UAV. The authors devised an iterative method based on successive convex approximation (SCA) theory to solve the non-convex problem, jointly optimizing the UAV trajectory, the offloading amount, and the computational resource allocation. In [25], an OFDMA-based multi-user wireless-powered MEC system was studied to maximize the weighted sum computation rate. The study proposed an efficient algorithm based on block coordinate descent to achieve near-optimal performance. In [11], a wireless-powered MEC network with multiple users was studied to maximize computational efficiency. The research considered and evaluated the offloading based on both TDMA and non-orthogonal multiple access (NOMA) and proposed two iterative algorithms and two alternative optimization algorithms, respectively.

The diversity among different IoT devices has led to considering user collaboration (UC) as an effective approach to enhance the system's computational capacity [28]. Consequently, numerous recent studies have explored the design of user collaboration in multi-user wireless-powered MEC networks. In [26], the authors investigated the application of UC and NOMA schemes in a two-user wireless-powered MEC network. They aimed to maximize computational efficiency by jointly designing energy beamforming, time, and power allocation. In [27], the authors proposed a NOMA-assisted UC scheme in a three-node WPT-MEC system. The system comprises a far user (FU), a near user (NU), and an access point (AP). The FU employs NOMA, enabling the simultaneous offloading of data to the NU and the AP, thereby overcoming the dual far-near effect of the FU. However, in [26] and [27], only two-user network models were considered, precisely the NUassisted FU scenario. The derived results may face challenges when extending to multi-user scenarios. In [28], a multiuser WPT-MEC system was studied. The authors proposed a multi-user collaboration scheme based on OFDMA to enhance computational efficiency. The scheme allows users to divide computational tasks into multiple parts for local computation, offloading to corresponding assistants, and the HAP for remote execution. However, the collaborative relationships among users are fixed in this work. In fact, the optimal collaborative relationships among users may change as the state of the network changes.

Table 1 summarizes the closely related works in wirelesspowered MEC networks. A comparison between these works and ours primarily focuses on the number of users, collaboration mode, objective, technical method, and contributions.

C. CONTRIBUTIONS

The main contributions of this paper are summarized as follows.

- *Firstly*, to maximize the WSCR of the system, we formulate a mixed-integer programming (MIP) problem for the multi-user wireless-powered MEC network. The problem jointly optimizes the collaboration strategy, time allocation, and data assignment among users and the HAP, subject to energy causality constraints and minimum processing data volume requirements.
- *Secondly*, we decompose the non-convex optimization problem into the time and data allocation sub-problem and the collaboration strategy design sub-problem. An interior-point method based algorithm is developed to solve the time and data allocation sub-problem given a fixed collaboration strategy.

- *Thirdly*, we propose a low-complexity priority-based iterative search algorithm to determine the multi-user collaboration strategy. This allows for finding a sub-optimal solution within a reduced search space.
- *Finally*, to further reduce algorithm execution time, we devise a deep learning-based algorithm building on the first two algorithms. Simulation results are performed to demonstrate the validity of the proposed algorithms.

It is worth noting that smart energy harvesting and edge computing are crucial technologies and challenges in the context of 6G [30]. Meanwhile, in the realm of 6G, the IIoT stands out as a typical application scenario within IoT. Therefore, the investigation of wireless-powered mobile edge computing within the context of IIoT holds significant importance.

D. ORGANIZATION

The organization of our work is demonstrated as follows. The system model is introduced in Section II, and the weighted sum computation rate maximization problem is formulated in Section III. In Section IV, three algorithms are proposed to solve the formulated problem. Numerical simulation results are provided to validate the proposed algorithms in Section V. Finally, this paper is concluded in Section VI. The code and dataset are available at https://github.com/CPNGroup/WPTMEC.

II. SYSTEM MODEL

A. NETWORK MODEL

In this paper, we consider a multi-user wireless-powered MEC network. The network consists of a HAP integrated with a computing unit and a set of single-antenna devices, denoted as $\mathcal{N} = \{1, \dots, N\}$. The HAP with N antennas provides stable wireless energy supply and enables task offloading for the N devices, where each device can convert the received RF signals into power for subsequent data offloading and processing. The system time is segmented into consecutive time frames, each with a uniform duration of T, deliberately chosen to be smaller than the channel coherence time. For instance, this could be on the order of several seconds in a static IoT environment [31]. At the beginning of each time frame, the devices are categorized into three types based on the diversity in data importance, computational capabilities, and channel qualities. The three classes are source devices (SDs), auxiliary devices (ADs), and independent devices (IDs), as illustrated in Fig. 1. Let $o = \{o_1, o_2, \dots, o_m\}, p = \{p_1, p_2, \dots, p_m\}, \text{ and } q =$ $\{q_1, q_2, \ldots, q_{N-2m}\}$ denote the sets of SDs, ADs, and IDs, respectively, where 2m < N and $\boldsymbol{o} \cup \boldsymbol{p} \cup \boldsymbol{q} = \mathcal{N}$. Each pair $\{o_i, p_i\}, \forall i \in \{1, 2, \dots, m\}$, forms a collaborative cluster, where m is the number of collaborative clusters. As shown in Fig. 2, the data of each source device (SD) is partitioned into three parts for processing at the local device, HAP, and corresponding auxiliary device (AD), respectively. Specifically, for the portion of an SD's data that is designated for processing at the corresponding AD, it will first be

location are
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ing mode sele- g) and time ptimized.
U frequencies, de selection
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TABLE 1. Overview of the literature related to wireless-powered MEC networks.

transmitted to the HAP and then forwarded to the AD by the HAP. This avoids communication issues due to distance limitations between the SD and AD. The data of each AD is processed locally. The data of each independent device (ID) is partitioned into two parts for processing at the local device and HAP, respectively. The main symbols and notations are summarized in Table 2.

B. HARVEST-AND-OFFLOADING PROTOCOL

To avoid co-channel interference and facilitate collaboration among users as well as between users and HAP, we propose a harvest-and-offload protocol utilizing UC and OFDMA, as depicted in Fig. 3. This protocol achieves collaboration among users as well as between users and HAP by planning the time allocation for different categories of IoT devices and HAP in each time frame. It is assumed that the same frequency band is utilized for both RF energy signals and data transmission. Moreover, based on OFDMA, the system bandwidth is divided into N sub-channels. The HAP randomly assigns each device to a sub-channel to avoid inter-device interference. Consistent with [25], we adopt a quasistatic scenario where the wireless channel remains static within each time block but varies across blocks.

As shown in Fig. 3(a), the time block of the HAP is divided into three stages. The first stage, with a duration of $\alpha_1 T$ where $\alpha_1 \in (0, 1]$, is allocated for WPT. The second stage, lasting for $\alpha_2 T$ where $\alpha_2 \in [0, 1)$, involves the HAP receiving offloaded data from all devices, including data



FIGURE 1. The wireless-powered MEC network with multi-user cooperation.



FIGURE 2. The data allocation for three types of devices. (The SD and AD form a collaborative cluster.)

designated for processing at the HAP and ADs. During the remaining time of $T - (\alpha_1 + \alpha_2)T$, the HAP processes data received from ADs and IDs, which is intended for processing at the HAP.

As illustrated in Fig. 3(b), IDs can locally process a portion of their data throughout the entire time block. Concurrently, during the initial $\alpha_1 T$ of each time frame, it is dedicated to energy harvesting (EH). Subsequently, in the following $\alpha_2 T$, IDs offload the data prepared for processing at the HAP through uplinks. It is noteworthy that the time for data offloading from each ID must be less than or equal to $\alpha_2 T$.

Fig. 3(c) illustrates the time allocation and interaction process between two devices in a cooperative cluster. In the initial $\alpha_1 T$ of each time frame, both the SD o_i and AD p_i undergo EH. In the subsequent $\alpha_2 T$, SD o_i offloads the data prepared for processing at both the HAP and AD p_i to the

TABLE 2. Summary of system notations.

Notation	Description
\mathcal{N}	The set of IoT devices
0	The set of SDs
m	The number of collaborative clusters
p	The set of ADs
q	The set of IDs
T	The length of a time frame
α_1	The fraction of time that the HAP broadcasts RF
	energy for the IoT devices to harvest
α_2	The fraction of time that the HAP receives data
	offloaded by the IoT devices
η	The energy harvesting efficiency
p_n	The transmit power of the nth IoT device
P	The transmit power of the HAP
h_n	The wireless channel gain between the HAP and the
	nth device
E_n^h	The amount of energy that the nth device harvests
t_n^{loc}	The local computation time of the n th device
l_n^{loc}	The amount of data allocated for local processing
	by the <i>n</i> th IoT device
ϕ_n	The number of cycles needed to process one bit of
	task data from the nth IoT device
f_n	The processor's computing speed (cycles per second)
	of the nth IoT device
k_n	The computation energy efficiency coefficient
E_n^{loc}	The energy consumption of the n th device due to
	the local computing
l_n^{ap}	The amount of data allocated for processing on the
	HAP by the nth IoT device
$l_{o_i}^{p_i}$	The amount of data allocated for processing on the
	AD p_i by the SD o_i
t_n^{off}	The data transfer time for the nth IoT device
	during computing offloading
E_n^{off}	The energy consumption of the n th device due to
	computing offloading
l_{th}	The minimum data processing requirement of each
	device within a time frame
f_{ap}	The processor's computing speed (cycles per
	second) of the HAP
l_o	The profile of data distribution scheme for SDs
l_p	The profile of data distribution scheme for ADs
1	The profile of data distribution scheme for IDs

HAP via the uplink. Notably, the actual data offloading time from SD o_i to the HAP must be less than or equal to $\alpha_2 T$. However, it is only after $\alpha_1 T + \alpha_2 T$ that the HAP begins transmitting data to ADs, which is designated for processing by ADs and originates from SDs.

The time $t_{ap}^{p_i}$ for HAP to send data to AD p_i through the downlink depends on their transmission rate and the amount of data that SD o_i allocates for processing by AD p_i . Since AD p_i can only start assisting processing after receiving data from SD o_i , the assisting processing time of AD p_i cannot exceed $T - (\alpha_1 + \alpha_2)T - t_{ap}^{p_i}$. Before the assisting processing, AD p_i needs to process all its data locally. Additionally,



FIGURE 3. The harvest-and-offloading protocol with UC and OFDMA.

throughout the entire time block, SD o_i can process a portion of its data locally. It is worth noting that the delay for results collection is ignored, as the size of the results is negligible compared to the input data size, as stated in [32], [33].

The proposed protocol utilizes OFDMA to reduce the possibility of co-channel interference among devices. It employs Time Division Multiplexing (TDD) on subchannels, enabling energy harvesting and data transmission to share the same frequency band. This approach enhances both spectral and energy efficiency. Simultaneously, the protocol promotes efficient collaboration among various system components. It achieves this by strategically allocating time among different devices, ensuring sufficient time for energy harvesting, data transmission, and processing within each timeframe. Nevertheless, the practical implementation of this protocol faces potential challenges. These include communication quality fluctuations caused by dynamic environmental factors, the influence of collaboration delays on real-time performance, and variations in energy harvesting efficiency due to environmental changes and device mobility. Due to space limitations, a detailed discussion of these challenges will not be pursued.

C. WIRELESS ENERGY TRANSMISSION

At the beginning of each time frame, the HAP broadcasts wireless power to each user for a duration of $\alpha_1 T$. Therefore, the harvested energy by device *n* can then be expressed as

$$E_n^h = \begin{cases} \eta P h_n \alpha_1 T, \quad \eta P h_n \alpha_1 T \leq E_n^{max}, \\ E_n^{max}, \quad \eta P h_n \alpha_1 T > E_n^{max}, \\ \end{cases}, \quad (1)$$

where constant $\eta \in (0, 1)$ denotes the EH efficiency factor and *P* represents the transmit power at the HAP. h_n represents the wireless channel gain between the *n*th IoT device and the HAP, assuming it is the same for both uplink and downlink due to channel reciprocity [23], [34]. E_n^{max} denotes the maximum battery capacity of the *n*th IoT device.

It should be noted that while some existing works have discussed the non-linear EH model and argued that it is closer to real-world scenarios [35], numerous studies (e.g., [36], [37]) have reported that EH circuits typically operate within the linear region when the distance between





the HAP and IoT devices is not very short. This is due to the significant influence of path loss in WPT systems on the amount of harvested energy. Additionally, the linear model is frequently employed as they are tractable and accurate for moderate harvested power [16]. Thus, similar to previous studies [16], [23], [36], [37], we also adopt the linear EH model when $E_n^h \leq E_n^{max}$ as shown in (1) in this paper.

D. DATA PROCESSING AND CONSTRAINTS1) SOURCE DEVICES

The data of SD o_i is partitioned into three components: local processing, processing by the HAP, and processing by the corresponding AD p_i . The data processed by AD p_i needs to be offloaded to the HAP first and then transmitted from the HAP to AD p_i for processing. Regarding energy consumption, the SD o_i 's energy expenditure can be divided into two main parts: local processing and computing offloading. These two aspects are analyzed as follows.

- Local processing: An IoT device can harvest energy and process its data simultaneously [38], [39]. Let f_{o_i} denote the processor's computational speed (cycles per second) of SD o_i , and let $l_{o_i}^{loc}$ denote the amount of data processed locally. Therefore, the time required for local processing is given by $t_{o_i}^{loc} = \frac{l_{o_i}^{loc}\phi_{o_i}}{f_{o_i}}$, where ϕ_{o_i} represents the number of cycles needed to process one bit of task data from SD o_i . Moreover, the local processing must be completed before the time frame ends, which can be expressed as $t_{o_i}^{loc} \leq T$. The energy consumed by local computation is given by $E_{o_i}^{loc} = k_{o_i} f_{o_i}^2 l_{o_i}^{loc} \phi_{o_i}$, where k_{o_i} represents the computational energy efficiency factor [28].
- Computing offloading: Let $l_{o_i}^{ap}$ denote the amount of SD o_i 's data processed at the HAP, and $l_{o_i}^{p_i}$ denote the amount of SD o_i 's data processed at the AD p_i . Therefore, the time required for SD o_i to offload its data is given by $t_{o_i}^{off} = \frac{l_{o_i}^{ap} + l_{o_i}^{o_i}}{R_{o_i}^{ap}}$, where $R_{o_i}^{ap} = B \log_2(1 + \frac{p_{o_i}h_{o_i}}{N_0})$ denotes the data transmission rate between SD o_i and the HAP. Here, *B* denotes the sub-channel bandwidth, N_0 denotes the receiver noise power, and p_{o_i} represents

the transmit power at SD o_i . The offloading time of SD o_i needs to satisfy $t_{o_i}^{off} \leq \alpha_2 T$, since the HAP receives the offloaded data during $\alpha_2 T$, as shown in Fig. 3. Additionally, the energy consumption due to offloading the data is given by $E_{o_i}^{off} = p_{o_i} t_{o_i}^{off} = \frac{p_{o_i} (d_{o_i}^{op} + b_{i_i}^{o_i})}{R_{o_i}^{op}}$.

According to the energy causality property, an IoT device cannot consume more energy than it harvests within a given time frame. Hence, the energy consumption of SD o_i must adhere to the condition: $E_{o_i}^{loc} + E_{o_i}^{off} \leq E_{o_i}^{h}$. Furthermore, the minimum data processing constraint dictates that each IoT device's data needs to be processed at least l_{th} bits per time frame. Therefore, the total amount of data from SD o_i that needs to be processed must satisfy the following inequality:

$$l_{o_i}^{loc} + l_{o_i}^{ap} + l_{o_i}^{p_i} \ge l_{th}.$$
 (2)

2) AUXILIARY DEVICES

The data of AD p_i is processed exclusively on the local device. However, its energy consumption can be divided into two main components: local processing energy consumption and auxiliary processing energy consumption. This is due to the need to process a portion of SD o_i 's data. These two aspects are analyzed as follows.

• Auxiliary processing: The time required for AD p_i to complete auxiliary processing can be calculated as $t_{p_i}^{o_i} = \frac{l_{o_i}^{o_i}\phi_{o_i}}{f_{p_i}}$, where f_{p_i} represents the computational speed of AD p_i 's processor. Additionally, auxiliary processing can only start after AD p_i receives the data from the HAP. The transmission time is denoted as $t_{ap}^{p_i} = \frac{l_{o_i}^{p_i}}{R_{ap}^{p_i}}$, where $R_{ap}^{p_i} = Blog_2(1 + \frac{Ph_{p_i}}{N_0})$ denotes the data transmission rate between AD p_i and the HAP. Therefore, as shown in Fig. 3, the following inequality must be satisfied:

$$(\alpha_1 + \alpha_2)T + t_{ap}^{p_i} + t_{p_i}^{o_i} \le T.$$
 (3)

Moreover, the energy consumption resulting from auxiliary processing can be calculated as $E_{p_i}^{o_i} = k_{p_i} f_{p_i}^{2} l_{o_i}^{p_i} \phi_{o_i}$. • *Local processing:* Let $l_{p_i}^{loc}$ denote the amount of AD p_i 's data processed locally. The time required for AD p_i to complete local processing can be calculated as $t_{p_i}^{loc} = \frac{l_{p_i}^{loc} \phi_{p_i}}{f_{p_i}}$. Additionally, the minimum data processing constraint requires that the total amount of data from AD p_i that needs to be processed satisfies the following inequality:

$$l_{p_i}^{loc} \ge l_{th}.\tag{4}$$

And the energy consumption resulting from local processing can be calculated as $E_{p_i}^{loc} = k_{p_i} f_{p_i}^3 l_{p_i}^{loc} = k_{p_i} f_{p_i}^2 l_{p_i}^{loc} \phi_{p_i}$.

According to the energy causality property, the energy consumption of AD p_i must adhere to the condition: $E_{p_i}^{loc} + E_{p_i}^{o_i} \leq E_{p_i}^h$. And the total computation time of AD p_i

includes both auxiliary processing and local processing, so the following inequality must hold:

$$t_{p_i}^{loc} + t_{p_i}^{o_i} \le T.$$
 (5)

3) INDEPENDENT DEVICES

The IDs do not form collaborative clusters with other devices. Consequently, the data of ID q_k is divided into two parts: one part is processed locally, and the other is processed on the HAP. In terms of energy consumption, the energy expenditure of ID q_k can be divided into two main parts: local processing and computing offloading. Next, these two aspects are analyzed as follows.

- Local processing: Let $l_{q_k}^{loc}$ represent the amount of data allocated for local processing by ID q_k . Consequently, the local processing delay can be calculated using the formula $t_{q_k}^{loc} = \frac{l_{q_k}^{loc}\phi_{q_k}}{f_{q_k}}$, where f_{q_k} represents the computational speed of ID q_k 's processor. The local processing energy consumption is given by $E_{q_k}^{loc} = k_{q_k} f_{q_k}^2 l_{q_k}^{loc} \phi_{q_k}$. Similarly, the local processing delay needs to satisfy the inequality: $t_{q_k}^{loc} \leq T$.
- Computing offloading: Let $l_{q_k}^{q_k}$ represent the amount of data from ID q_k that is processed at the HAP. The offloading time can be calculated using the formula $t_{q_k}^{off} = \frac{l_{q_k}^{op}}{R_{q_k}^{op}}$, where $R_{q_k}^{ap} = Blog_2(1 + \frac{p_{q_k}h_{q_k}}{N_0})$ represents the data transmission rate between ID q_k and the HAP. Similarly, the offloading time of ID q_k needs to satisfy the condition $t_{q_k}^{off} \leq \alpha_2 T$. Additionally, the energy consumption due to data offloading can be obtained as $E_{q_k}^{off} = p_{q_k} t_{q_k}^{off}$.

Similarly, the total energy consumption of ID q_k must satisfy the energy causality constraint represented by $E_{q_k}^{loc} + E_{q_k}^{off} \leq E_{q_k}^h$. Moreover, there exists a constraint on the minimum amount of data to be processed, expressed as $l_{q_k}^{loc} + l_{q_k}^{ap} \geq l_{th}$.

4) HYBRID ACCESS POINT

The HAP initially transmits RF signals within $\alpha_1 T$ and subsequently receives data offloaded by IoT devices within $\alpha_2 T$. Processing commences only after data reception. Consequently, the processing time for the HAP is $(1 - \alpha_1 - \alpha_2)T$, and the workload which needs to be processed by the HAP is given by $\sum_{i=1}^{m} l_{o_i}^{ap} \phi_{o_i} + \sum_{k=1}^{N-2m} l_{q_k}^{ap} \phi_{q_k}$. Thus, the following inequality must be satisfied:

$$\frac{\sum_{i=1}^{m} l_{o_i}^{ap} \phi_{o_i} + \sum_{k=1}^{N-2m} l_{q_k}^{ap} \phi_{q_k}}{f_{ap}} \le (1 - \alpha_1 - \alpha_2)T, \quad (6)$$

where f_{ap} represents the processor's computing speed (cycles per second) of the HAP.

III. PROBLEM FORMULATION AND ANALYSIS

Our study aims to maximize the WSCR of the IIoT system by jointly optimizing the collaboration strategy $\psi = \{m, o, p, q\}$ among IoT devices and the HAP, the system time allocation $\alpha = \{\alpha_1, \alpha_2\}$, and the data allocation $l = \{l_o, l_p, l_q\}$, where

 $l_o = [l_{o_1}^{loc}, l_{o_1}^{ap}, l_{o_1}^{p_1}, \dots, l_{o_m}^{loc}, l_{o_m}^{ap}, l_{o_m}^{p_m}]^T$, $l_p = [l_{p_1}^{loc}, \dots, l_{p_m}^{loc}]^T$, and $l_q = [l_{q_1}^{loc}, l_{q_1}^{ap}, \dots, l_{q_{N-2m}}^{loc}]^T$. Additionally, the optimization is conducted considering energy causality, the minimum data processing requirements, and time constraints. Mathematically, the WSCR can be formulated as

$$\begin{split} \mathcal{P}_{1} : \max_{\psi,\alpha,l} & \sum_{i=1}^{m} \left\{ w_{o_{i}} \left(t_{o_{i}}^{loc} + t_{o_{i}}^{ap} + t_{o_{i}}^{p_{i}} \right) + w_{p_{i}} t_{p_{i}}^{loc} \right\} \\ & + \sum_{k=1}^{N-2m} w_{q_{k}} \left(t_{q_{k}}^{loc} + t_{q_{k}}^{ap} \right) \\ \text{s.t.} & (7a) : |o| = |p| = m \in \left\{ 0, 1, 2, \dots, \left\lfloor \frac{N}{2} \right\rfloor \right\} \\ & (7b) : o \cup p \cup q = \mathcal{N} \\ & (7c) : \frac{t_{o_{i}}^{loc} \phi_{o_{i}}}{f_{o_{i}}} \leq T, \forall i \in \{1, 2, \dots, m\} \\ & (7d) : \frac{t_{o_{i}}^{ap} + t_{o_{i}}^{p_{i}}}{R_{o_{i}}^{ap}} \leq \alpha_{2}T, \forall i \in \{1, 2, \dots, m\} \\ & (7e) : \frac{\phi_{o_{i}} t_{o_{i}}^{p_{i}}}{f_{p_{i}}} \leq (1 - \alpha_{1} - \alpha_{2})T - \frac{t_{o_{i}}^{p_{i}}}{R_{op}^{p_{i}}}, \forall i \\ & (7f) : \frac{\phi_{p_{i}} t_{p_{i}}^{p_{i}}}{f_{p_{i}}} \leq T - \frac{\phi_{o_{i}} t_{o_{i}}^{p_{i}}}{f_{p_{i}}}, \forall i \in \{1, 2, \dots, m\} \\ & (7g) : \frac{t_{q_{k}}^{lac} \phi_{q_{k}}}{f_{q_{k}}} \leq \alpha_{2}T, \forall k \in \{1, 2, \dots, N-2m\} \\ & (7h) : \frac{t_{q_{k}}}{R_{q_{k}}} \leq \alpha_{2}T, \forall k \in \{1, 2, \dots, N-2m\} \\ & (7h) : \frac{\sum_{i=1}^{m} \phi_{o_{i}} t_{o_{i}}^{ap} + \sum_{k=1}^{N-2m} \phi_{q_{k}} t_{q_{k}}^{ap}}{f_{q_{k}}}} \leq (1 - \alpha_{1} - \alpha_{2})T \\ & (7i) : \frac{\sum_{i=1}^{m} \phi_{o_{i}} t_{o_{i}}^{ap} + \sum_{k=1}^{N-2m} \phi_{q_{k}} t_{q_{k}}^{ap}}{f_{q_{k}}}} \leq (1 - \alpha_{1} - \alpha_{2})T \\ & (7i) : \frac{R_{o_{i}}^{p_{i}}}{R_{op}^{p_{i}}} \leq (1 - \alpha_{1} - \alpha_{2})T, \forall i \in \{1, 2, \dots, m\} \\ & (7h) : k_{o_{i}} f_{o_{i}}^{2} t_{o_{i}}^{loc} \phi_{o_{i}} + \frac{P_{o_{i}} (f_{o_{i}}^{p_{i}} + t_{o_{i}}^{p_{i}})}{R_{o_{i}}^{o_{i}}} \leq E_{o_{i}}^{h}, \forall i \\ & (7h) : k_{o_{i}} f_{o_{i}}^{2} t_{o_{i}}^{loc} \phi_{o_{i}} + \frac{P_{o_{i}} (f_{o_{i}}^{p_{i}} + t_{o_{i}}^{p_{i}})}{R_{o_{i}}^{o_{i}}}} \leq E_{o_{i}}^{h}, \forall i \\ & (7h) : k_{q_{k}} t_{q_{k}}^{2} t_{q_{k}}^{loc} \phi_{q_{k}} + \frac{P_{q_{k}} t_{q_{k}}^{q_{k}}}{R_{q_{k}}^{q_{k}}}} \leq E_{q_{k}}^{h}, \forall k \\ & (7o) : t_{o_{i}}^{loc} + t_{o_{i}}^{p_{i}} + t_{o_{i}}^{p_{i}}} \geq t_{h}, \forall i \in \{1, 2, \dots, m\} \\ & (7p) : t_{o_{i}}^{loc} \in t_{i}, \forall i, \forall i \in \{1, 2, \dots, m\} \\ & (7p) : t_{o_{i}}^{loc} \in t_{i}, \psi_{i}, \psi_{i} \in \{1, 2, \dots, m\} \\ & (7p) : t_{o_{i}}^{loc} + t_{q_{k}}^{p_{k}} \geq t_{h}, \forall k \in \{1, 2, \dots, N-2m\} \\ & (7p) : t_{o_{i}}^{loc} \in t_{i}, \psi_{i}, \psi_{i}, \psi_{i} \in \{1, k, m, m\} \\ & (7p) : t_{$$

where w_{o_i} , w_{p_i} , w_{q_k} are unique weights assigned to SD o_i , AD p_i , and ID q_k , respectively. Constraints (7*a*) and (7*b*) pertain to the formation of collaborative clusters, requiring an SD to collaborate with a unique AD. Constraints (7*c*) and (7*d*) concern the time overhead of SDs. Constraints (7*e*) and (7*f*) relate to the time overhead of ADs. Constraints (7*g*) and (7*h*) involve the time overhead of IDs. Constraints (7*i*) to (7*k*) govern the time allocation of the HAP. Constraints (7*l*) to (7*n*) pertain to the energy consumption causality for SDs, ADs, and IDs, respectively. Constraints (7*o*) to (7*q*) specify the minimum data processing requirements for SDs, ADs, and IDs, respectively. Constraint (7r) pertains to the non-negativity for the data allocation.

According to our assumption, the energy consumed by the *n*th device for continuous data processing within a time frame should be greater than the energy collected.¹ Mathematically, this can be expressed as $k_n f_n^3 T > E_n^h$. Based on this assumption, we give the following lemma.

Lemma 1: At the end of each time frame, all devices exhaust the energy collected during that time frame.

Proof: Assume that under the optimal solution of problem \mathcal{P}_1 , there is a device that has not exhausted its energy. Based on the above assumption, we can increase the amount of data processed locally by that device to exhaust the remaining energy without violating the constraints in \mathcal{P}_1 . As a result, the WSCR is clearly enhanced, which contradicts the initial assumptions. This proves Lemma 1. Hence, we can disregard the remaining energy from the last time frame.

Several lemmas are given next to simplify the constraints of the problem.

Lemma 2: Constraints (7c) and (7g) must hold.

Proof: Given $k_n f_n^3 T > E_n^h$, it follows that the actual local processing time of SD o_i cannot exceed T, i.e., $\frac{l_{o_i}^{loc}\phi_{o_i}}{f_{o_i}} < T$, thus satisfying constraint (7*c*). Similarly, it can be demonstrated that constraint (7*g*) must also be met. \blacksquare *Lemma 3:* The constraint (7*j*) must hold.

Proof: Due to constraint (7e), $\frac{I_{o_i}^{i}}{R_{qp}^{P_i}} \leq (1-\alpha_1-\alpha_2)T - \frac{\phi_{o_i}I_{o_i}^{P_i}}{f_{p_i}}$. Additionally, based on constraint (7r), it can be derived that $\frac{I_{o_i}^{P_i}}{R_{qp}^{P_i}} \leq (1-\alpha_1-\alpha_2)T$. As a result, constraint (7j) can be removed.

Lemma 4: The constraint (7k) must hold.

Proof: From constraints (7*i*) and (7*r*), it follows that $0 \leq \frac{\sum_{i=1}^{m} \phi_{o_i} l_{o_i}^{a_p} + \sum_{k=1}^{N-2m} \phi_{q_k} l_{q_k}^{a_p}}{f_{a_p}} \leq (1 - \alpha_1 - \alpha_2)T$, which implies $\alpha_1 + \alpha_2 \leq 1$. Moreover, from constraints (7*h*), (7*l*) and (7*o*), it is evident that $\alpha_1 > 0$ and $\alpha_2 \geq 0$. Consequently, we can infer that $\alpha_1 \in (0, 1]$ and $\alpha_2 \in [0, 1)$. This establishes that constraint (7*k*) holds.

Additionally, according to (1), it is apparent that constraints (7l) - (7n) display nonlinearity. To simplify problem resolution, we reformulate these three constraints as:

$$\begin{aligned} k_{o_i} f_{o_i}^2 l_{o_i}^{loc} \phi_{o_i} &+ \frac{p_{o_i} (l_{o_i}^{ap} + l_{o_i}^{p_i})}{R_{o_i}^{ap}} \leq \eta P h_{o_i} \alpha_1 T, \forall i, \\ k_{o_i} f_{o_i}^2 l_{o_i}^{loc} \phi_{o_i} &+ \frac{p_{o_i} (l_{o_i}^{ap} + l_{o_i}^{p_i})}{R_{o_i}^{ap}} \leq E_{o_i}^{max}, \forall i, \\ k_{p_i} f_{p_i}^2 l_{p_i}^{loc} \phi_{p_i} &+ k_{p_i} f_{p_i}^2 l_{o_i}^{p_i} \phi_{o_i} \leq \eta P h_{p_i} \alpha_1 T, \forall i, \\ k_{p_i} f_{p_i}^2 l_{p_i}^{loc} \phi_{p_i} &+ k_{p_i} f_{p_i}^2 l_{o_i}^{p_i} \phi_{o_i} \leq E_{p_i}^{max}, \forall i, \end{aligned}$$

¹In wireless-powered MEC systems, wireless devices (WDs) are typically energy-constrained. This means a WD can exhaust all harvested energy within a given time frame by operating at maximum computing speed. Thus, in accordance with [22], [23], [25], [38], we adopt this assumption. Exploring the potential remaining energy of each IoT device after a time frame will be a future focus of our research.

$$k_{q_k} f_{q_k}^2 l_{q_k}^{loc} \phi_{q_k} + \frac{p_{q_k} l_{q_k}^{ap}}{R_{q_k}^{ap}} \le \eta P h_{q_k} \alpha_1 T, \forall k,$$

$$k_{q_k} f_{q_k}^2 l_{q_k}^{loc} \phi_{q_k} + \frac{p_{q_k} l_{q_k}^{ap}}{R_{q_k}^{ap}} \le E_{q_k}^{max}, \forall k.$$
(8)

Building upon the lemmas and constraint transformations mentioned earlier, the original problem can be reduced to:

$$\mathcal{P}_{2}: \max_{\psi,\alpha,l} \sum_{i=1}^{m} \left\{ w_{o_{i}} \left(l_{o_{i}}^{loc} + l_{o_{i}}^{ap} + l_{o_{i}}^{p_{i}} \right) + w_{p_{i}} l_{p_{i}}^{loc} \right\} \\ + \sum_{k=1}^{N-2m} w_{q_{k}} \left(l_{q_{k}}^{loc} + l_{q_{k}}^{ap} \right) \\ \text{s.t.} \quad (9a): (7a), (7b), (7d) - (7f), \\ (7h), (7i), (7o) - (7r), (8) \qquad (9)$$

Lemma 5: The problem \mathcal{P}_2 is NP-hard.

Proof: To establish the NP-hardness of \mathcal{P}_2 , we present a specific case of the problem and demonstrate its NPhardness. Assume that given the number of collaborative clusters m > 0, the construction of collaborative clusters is specified as $\{(i, i + m) \mid i \in \{1, 2, \dots, m\}\}$, where (i, i + m)represents the indices of the corresponding IoT devices. Consequently, we can obtain $o = \{i \mid i \in \{1, 2, \dots, m\}\}$, p = $\{i + m \mid i \in \{1, 2, \dots, m\}\}$, and $q = \{2m + 1, 2m + 2, \dots, N\}$. Based on this, problem \mathcal{P}_2 is reduced to \mathcal{P}_3 . It can be observed that the objective function and constraints of \mathcal{P}_3 are both linear. Additionally, the optimization variable *m* is an integer. As a result, \mathcal{P}_3 constitutes a mixed-integer linear programming (MILP) problem known as NP-hard. Therefore, this shows that problem \mathcal{P}_2 is also NP-hard.

$$\mathcal{P}_{3}: \max_{m,\boldsymbol{\alpha},\boldsymbol{l}} \sum_{i=1}^{m} \left\{ w_{o_{i}} \left(l_{o_{i}}^{loc} + l_{o_{i}}^{ap} + l_{o_{i}}^{p_{i}} \right) + w_{p_{i}} l_{p_{i}}^{loc} \right\} \\ + \sum_{k=1}^{N-2m} w_{q_{k}} \left(l_{q_{k}}^{loc} + l_{q_{k}}^{ap} \right) \\ \text{s.t.} \quad (10a): (7d) - (7f), (7h), (7i), (7o) - (7r), (8) \\ (10b): m \in \left\{ 1, 2, \dots, \left\lfloor \frac{N}{2} \right\rfloor \right\}.$$
(10)

IV. SOLUTION APPROACH

In this section, efficient algorithms are proposed to solve problem \mathcal{P}_2 . As shown in Fig. 4, problem (\mathcal{P}_2) can be decomposed into two sub-problems, namely, time and data allocation (\mathcal{P}_4) and collaboration strategy. Thus, we can optimize the time and data allocation { α , l} for a given collaboration strategy ψ^0 first. Then, we can search among all possible collaboration strategies to find an optimal or satisfying sub-optimal collaboration strategy ψ^* . However, due to the exponentially large search space, it is computationally prohibitive when N is large. Therefore, a low-complexity algorithm is needed. Next, we provide detailed solutions for these two subproblems.



FIGURE 4. The two-level optimization structure for solving \mathcal{P}_2 .

A. TIME AND DATA ALLOCATION

When the collaborative strategy ψ is fixed, \mathcal{P}_2 transforms into the following problem:

$$\mathcal{P}_{4}: \max_{\boldsymbol{\alpha}, \boldsymbol{l}} \sum_{i=1}^{m} \left\{ w_{o_{i}} \left(l_{o_{i}}^{loc} + l_{o_{i}}^{ap} + l_{o_{i}}^{p_{i}} \right) + w_{p_{i}} l_{p_{i}}^{loc} \right\} \\ + \sum_{k=1}^{N-2m} w_{q_{k}} \left(l_{q_{k}}^{loc} + l_{q_{k}}^{ap} \right) \\ \text{s.t.} (11a): (7d) - (7f), (7h), (7i), (7o) - (7r), (8)$$

$$(11)$$

Problem \mathcal{P}_4 is a convex optimization problem because both the objective function and constraints are linear in terms of the optimization variables. In order to facilitate the solution of the problem, \mathcal{P}_4 is first transformed into the standard matrix form:

$$\mathcal{P}_{5}: \min_{\boldsymbol{x}_{1}} \boldsymbol{c}_{1}^{T}\boldsymbol{x}_{1}$$

s.t. (12*a*) : $\boldsymbol{A}_{1}\boldsymbol{x}_{1} \leq \boldsymbol{b}$
(12*b*) : $\boldsymbol{x}_{1} \geq \boldsymbol{0}.$ (12)

The vector x_1 represents the optimization variables, and it is defined as $x_1 = [\alpha_1, \alpha_2, l_{o_1}^{loc}, l_{o_1}^{ap}, l_{o_1}^{p_1}, l_{p_1}^{loc}, \dots,$ $l_{o_m}^{loc}, l_{o_m}^{ap}, l_{o_m}^{p}, l_{p_m}^{loc}, l_{q_1}^{loc}, l_{q_1}^{ap}, \dots, l_{q_{N-2m}}^{loc}, l_{q_{N-2m}}^{ap}]^T$. The weight corresponding coefficients to the optimization variables in the objective function of \mathcal{P}_{4} are represented by $c_{1},$ given as **C**1 _ $[0, 0, -w_{o_1}, -w_{o_1}, -w_{o_1}, -w_{p_1}, \dots, -w_{o_m}, -w_{o_m}, -w_{o_m}, -w_{o_m}, -w_{p_m}, -w_{q_1}, -w_{q_1}, \dots, -w_{q_{N-2m}}, -w_{q_{N-2m}}]^T$. The matrix A_1 represents the coefficients in the constraints, which has a total of 2N + 2 columns and 4N + m + 1 rows. The vector **b** represents the constants on the right-hand side of the inequality in the constraints, and it has a total of 4N + m + 1elements.

Introducing slack variables allows for the transformation of problem \mathcal{P}_5 into the following form:

$$\mathcal{P}_{6}: \min_{x_{1},x_{2}} c_{1}^{T} x_{1} + c_{2}^{T} x_{2}$$

s.t. (13*a*): $A_{1}x_{1} + x_{2} = b$,
(13*b*): $x_{1}, x_{2} \geq 0$, (13)

 x_2 is a vector of slack variables with 4N + m + 1 elements, while $c_2 = 0$ and its length is also 4N + m + 1. Let $c_3 = [c_1; c_2], x_3 = [x_1; x_2], A_2 = [A_1, I]$, the problem \mathcal{P}_6 can be equated to:

$$\mathcal{P}_7: \min_{\mathbf{x}_3} \mathbf{c}_3^T \mathbf{x}_3$$

s.t. (14*a*) : $\mathbf{A}_2 \mathbf{x}_3 = \mathbf{b}$
(14*b*) : $\mathbf{x}_3 \succeq \mathbf{0}$. (14)

To eliminate the inequality constraints in the problem \mathcal{P}_7 , a barrier function is introduced for transformation:

$$\mathcal{P}_8: \min_{\boldsymbol{x}_3} \boldsymbol{c}_3^T \boldsymbol{x}_3 - \mu \sum_{i=1}^{6N+m+3} log(\boldsymbol{x}_3^i)$$

s.t. (15*a*): $\boldsymbol{A}_2 \boldsymbol{x}_3 = \boldsymbol{b},$ (15)

where x_3^i represents the *i*th element in x_3 and $\mu > 0$. Additionally, the solution obtained by solving problem \mathcal{P}_8 gets closer to the optimal solution of problem \mathcal{P}_7 as the value of μ decreases.

It can be observed that problem \mathcal{P}_8 remains a convex optimization problem, allowing the utilization of the Lagrangian duality method to obtain an optimal solution when the value of μ is given. The Lagrangian function for \mathcal{P}_8 is given by:

$$L_{\mu}(\mathbf{x}_{3}, \mathbf{v}) = \mathbf{c}_{3}^{T} \mathbf{x}_{3} - \mu \sum_{i=1}^{6N+m+3} \log\left(x_{3}^{i}\right) + \mathbf{v}^{T} (\mathbf{A}_{2} \mathbf{x}_{3} - \mathbf{b}), \quad (16)$$

where vector $\boldsymbol{\nu}$ comprises Lagrange multipliers. By applying the Karush-Kuhn-Tucker (KKT) conditions and setting $\lambda_i = \frac{\mu}{x_3^T}$, $\boldsymbol{\lambda} = [\lambda_1, \dots, \lambda_{6N+m+3}]^T$, $\boldsymbol{X} = diag(\mathbf{x}_3)$, $\boldsymbol{\Lambda} = diag(\boldsymbol{\lambda})$, and $\mathbf{1} = \{1, \dots, 1\}^T$, we can obtain the following results:

$$\boldsymbol{A}_2 \boldsymbol{x}_3 = \boldsymbol{b},\tag{17}$$

$$X\Lambda \mathbf{1} = \mu \mathbf{1},\tag{18}$$

and

$$\boldsymbol{A}_{2}^{T}\boldsymbol{v}+\boldsymbol{c}_{3}=\boldsymbol{\lambda}. \tag{19}$$

Based on the above transformations, the solution to problem \mathcal{P}_8 can be determined by solving the following equation:

$$F_{\mu}(\boldsymbol{\lambda}, \boldsymbol{\nu}, \boldsymbol{x}_3) = \begin{pmatrix} A_2 \boldsymbol{x}_3 - \boldsymbol{b} \\ \boldsymbol{X} \boldsymbol{\Lambda} \boldsymbol{1} - \mu \boldsymbol{1} \\ A_2^T \boldsymbol{\nu} + \boldsymbol{c}_3 - \boldsymbol{\lambda} \end{pmatrix} = \boldsymbol{0}.$$
 (20)

The Newton iterative method can be employed to solve the problem. Specifically, in the (k+1)th iteration, the first-order Taylor expansion of function $F_{\mu}(\lambda, \nu, x_3)$ is given by:

$$F_{\mu}(\boldsymbol{\lambda}, \boldsymbol{\nu}, \boldsymbol{x}_{3}) \approx F_{\mu}\left(\boldsymbol{\lambda}^{k}, \boldsymbol{\nu}^{k}, \boldsymbol{x}_{3}^{k}\right) + \nabla F_{\mu}\left(\boldsymbol{\lambda}^{k}, \boldsymbol{\nu}^{k}, \boldsymbol{x}_{3}^{k}\right) \begin{pmatrix} \boldsymbol{\lambda} - \boldsymbol{\lambda}^{k} \\ \boldsymbol{\nu} - \boldsymbol{\nu}^{k} \\ \boldsymbol{x}_{3} - \boldsymbol{x}_{3}^{k} \end{pmatrix}.$$
 (21)

Here, λ^k , v^k , and x_3^k represent the solutions obtained in the *k*th iteration. Therefore, solving the equation $F_{\mu}(\lambda, v, x_3) = \mathbf{0}$ is equivalent to solving the following equation:

$$F_{\mu}\left(\boldsymbol{\lambda}^{k},\boldsymbol{\nu}^{k},\boldsymbol{x}_{3}^{k}\right)+\nabla F_{\mu}\left(\boldsymbol{\lambda}^{k},\boldsymbol{\nu}^{k},\boldsymbol{x}_{3}^{k}\right)\begin{pmatrix}\boldsymbol{\lambda}-\boldsymbol{\lambda}^{k}\\\boldsymbol{\nu}-\boldsymbol{\nu}^{k}\\\boldsymbol{x}_{3}-\boldsymbol{x}_{3}^{k}\end{pmatrix}=\boldsymbol{0},\ (22)$$

where

$$\nabla F_{\mu}(\boldsymbol{\lambda}, \boldsymbol{\nu}, \boldsymbol{x}_{3}) = \begin{pmatrix} \mathbf{0} & \mathbf{0} & \mathbf{A}_{2} \\ \boldsymbol{X} & \mathbf{0} & \mathbf{A} \\ -\mathbf{I} & \mathbf{A}_{2}^{T} & \mathbf{0} \end{pmatrix}.$$
 (23)

Hence, determining the optimal solution involves solving the following equations:

$$\begin{pmatrix} \mathbf{0} & \mathbf{0} & A_2 \\ \mathbf{X}^k & \mathbf{0} & \mathbf{\Lambda}^k \\ -\mathbf{I} & A_2^T & \mathbf{0} \end{pmatrix} \begin{pmatrix} \Delta \boldsymbol{\lambda} \\ \Delta \boldsymbol{\nu} \\ \Delta \boldsymbol{x}_3 \end{pmatrix} = \begin{pmatrix} -\boldsymbol{\gamma}_1^k \\ -\boldsymbol{\gamma}_2^k \\ -\boldsymbol{\gamma}_3^K \end{pmatrix}, \quad (24)$$

where

$$\begin{pmatrix} \boldsymbol{\gamma}_1^k \\ \boldsymbol{\gamma}_2^k \\ \boldsymbol{\gamma}_3^k \end{pmatrix} = F_{\mu} \Big(\boldsymbol{\lambda}^k, \boldsymbol{v}^k, \boldsymbol{x}_3^k \Big), \qquad (25)$$

and

$$\begin{pmatrix} \Delta \boldsymbol{\lambda} \\ \Delta \boldsymbol{\nu} \\ \Delta \boldsymbol{x}_3 \end{pmatrix} = \begin{pmatrix} \boldsymbol{\lambda} - \boldsymbol{\lambda}^k \\ \boldsymbol{\nu} - \boldsymbol{\nu}^k \\ \boldsymbol{x}_3 - \boldsymbol{x}_3^k \end{pmatrix}.$$
 (26)

The solution obtained by the k + 1th iteration process is:

$$\begin{pmatrix} \boldsymbol{\lambda}^{k+1} \\ \boldsymbol{\nu}^{k+1} \\ \boldsymbol{x}^{k+1} \end{pmatrix} = \begin{pmatrix} \boldsymbol{\lambda}^{k} \\ \boldsymbol{\nu}^{k} \\ \boldsymbol{x}^{k} \end{pmatrix} + \beta \begin{pmatrix} \Delta \boldsymbol{\lambda} \\ \Delta \boldsymbol{\nu} \\ \Delta \boldsymbol{x}_{3} \end{pmatrix}, \quad (27)$$

where β is the iteration step. Thus, we can derive an approximate optimal solution to the problem \mathcal{P}_8 for a given value of μ . Next, we can decrease the value of μ and utilize the solution obtained with the previous μ value as the initial solution for the problem \mathcal{P}_8 with the updated μ value. Afterwards, we can apply the same approach to solve the updated problem \mathcal{P}_8 . The iterative process continues until the accuracy requirement is satisfied or the maximum number of iterations is reached. This approach requires fewer iterations than solving problem \mathcal{P}_8 with the minimum μ value, accelerating the solution process for problem \mathcal{P}_5 . In order to further accelerate the value of μ while iterating to solve problem \mathcal{P}_8 . This approach is outlined in Algorithm 1.

Remark 1 (Details on the Iteration Process of Algorithm 1): In our designed algorithm, the key factors influencing the iterative process include the convergence judgment criterion, the iteration step size β , and the attenuation coefficient φ . The convergence criterion determines if the algorithm has achieved a solution satisfying the accuracy requirement. In this paper, we define it as the largest component of the absolute error vector resulting from substituting the current solution into equation (25) and comparing it to the target result **0**. To limit excessive



FIGURE 5. Algorithm 1: Interior-Point Method Based Algorithm for Solving Problem $\mathcal{P}_{\text{S}}.$

iterations, we establish a maximum iteration count *L*. The iteration step size β regulates the variable updates' magnitude in each iteration. Typically, the step size choice involves a line search technique, such as the Goldstein Condition [40]. Moreover, the attenuation coefficient φ controls the decay of the barrier function in each iteration and is typically chosen within the range of (0, 1). Empirically, common choices for the attenuation coefficient include values such as 0.1, 0.2, or smaller. Experimenting with various values and observing their effects on the algorithm's performance is advisable.

Remark 2 (Computational Complexity Analysis of Algorithm 1): In each iteration, the computational complexity of computing the Jacobian matrix for $F_{\mu}(\lambda^{k}, \nu^{k}, \mathbf{x}_{3}^{k})$ is $\mathcal{O}(N^{3})$. Moreover, solving equation (24) requires $\mathcal{O}(N^{3})$, while calculating the updated $\boldsymbol{\gamma}_{1}^{k}, \boldsymbol{\gamma}_{2}^{k}, \boldsymbol{\gamma}_{3}^{k}$ and its paradigm needs $\mathcal{O}(N^{2})$. Therefore, the computational complexity of Algorithm 1 is approximately $\mathcal{O}(LN^{3})$.

The solution for problem \mathcal{P}_5 is identical to that for problem \mathcal{P}_4 . Therefore, given the collaborative strategy, we can derive the optimal solution for time and data allocation. Subsequently, our objective is to identify the most favorable collaborative strategy. To achieve this, we can employ an exhaustive search method to enumerate all feasible solutions $\{\boldsymbol{\psi}_k | k \in \mathcal{K}\}$ for collaborative strategies satisfying constraints (7*a*) and (7*b*). Next, we can solve problem \mathcal{P}_4 for each of these collaborative strategies. The optimal solution $\boldsymbol{\psi}^*$ can

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be determined by employing the following equation:

z

$$\boldsymbol{\psi}^* = \arg \max_{\boldsymbol{\psi}_k} f(\boldsymbol{\psi}_k), \tag{28}$$

where $f(\boldsymbol{\psi}_{k}) = \max_{\boldsymbol{\alpha}, l} \sum_{i=1}^{m_{k}} \{w_{o_{k,i}}(l_{o_{k,i}}^{loc} + l_{o_{k,i}}^{ap} + l_{o_{k,i}}^{p_{k,i}}) + w_{p_{k,i}}l_{p_{k,i}}^{loc}\} + \sum_{j=1}^{N-2m_{k}} w_{q_{k,j}}(l_{q_{k,j}}^{loc} + l_{q_{k,j}}^{ap}).$ Remark 3 (Computational Complexity Analysis of the

Remark 3 (Computational Complexity Analysis of the Exhaustive Search Scheme): Initially, we need to determine the number of collaborative clusters m, with a total of $\lfloor \frac{N}{2} \rfloor + 1$ possible scenarios. Furthermore, for a given number of collaborative clusters m > 0, there are $C_N^{2m}(2m-1)!!$ possible scenarios for constructing collaborative clusters. In this case, C_N^{2m} represents the number of scenarios in which 2m devices are selected from a total of N IoT devices to form collaborative clusters, and (2m-1)!! denotes the number of scenarios in which collaborative clusters are constructed using 2m IoT devices. In summary, there are a total of $1 + \sum_{m=1}^{m=\lfloor \frac{N}{2} \rfloor} C_N^{2m}(2m-1)!!$ possible scenarios for constructing collaborative clusters. For each instance of collaborative cluster construction, the maximum achievable objective value needs to be determined based on Algorithm 1. Thus, the overall computational complexity is $\mathcal{O}(LN^3(1 + \sum_{m=1}^{m=\lfloor \frac{N}{2} \rfloor} C_N^{2m}(2m-1)!!))$.

It is evident that the exhaustive search method becomes computationally prohibitive when N is large. Therefore, we must explore a low-complexity algorithm to identify the most favorable collaborative strategy.

B. COLLABORATION STRATEGY

This subsection presents a low-complexity, priority-based algorithm designed to find a suboptimal solution for collaboration strategy, offering a trade-off between computational complexity and system performance.

In this study, the diversity among IoT devices is considered in three main aspects: (1) the importance weight w_n assigned to the *n*th device's data, (2) the channel quality h_n between the *n*th device and the HAP, and (3) the local computational capacity f_n of the *n*th device. Based on these three variables, a priority function can be constructed to guide the formation of collaborative clusters.

$$O_n = w_n \frac{k_n f_n^2 \phi_n}{\frac{p_n}{R_n^{app}}} = \frac{w_n B \log_2 \left(1 + \frac{p_n h_n}{N_0}\right) k_n f_n^2 \phi_n}{p_n}.$$
 (29)

The term $k_n f_n^2 \phi_n$ represents the energy consumption per unit of data volume processed by the *n*th device, and $\frac{p_n}{R_n^{ap}}$ represents the energy consumption per unit of data volume offloaded by the *n*th device. Therefore, a higher value of O_n indicates a higher priority for the *n*th device to serve as SD. During the construction of a new collaborative cluster, the device with the highest priority among the current IDs is selected as SD, and the device with the lowest priority among the current IDs is selected as AD. Therefore, an iterative algorithm can be employed to determine the optimal number m^* of collaborative clusters.



FIGURE 6. Algorithm 2: Priority-Based Algorithm for Solving Problem \mathcal{P}_2 .

Initially, setting *m* to 0 indicates that all devices are IDs, and then Algorithm 1 is used to calculate the achievable objective value based on the current collaboration strategy. In each iteration, *m* is incremented by 1, and a new collaborative cluster is constructed by selecting the device with the highest priority among the current IDs as SD and the device with the lowest priority as AD. Algorithm 1 is used again to calculate the achievable objective value based on the current collaborative strategy. After iterating $\lfloor \frac{N}{2} \rfloor + 1$ times, we can obtain the achievable objective value for each value of *m*. Hence, the final output would be the collaboration strategy corresponding to the maximum achievable objective value. The summarized solution is outlined in Algorithm 2.

Remark 4 (Computational Complexity Analysis of the Priority-Based Iterative Algorithm): The proposed algorithm establishes a priority function to narrow down the search space to $\lfloor \frac{N}{2} \rfloor$. Additionally, in each iteration, Algorithm 1 is employed, incurring a computational cost of $\mathcal{O}(LN^3)$. As a result, the overall computational complexity amounts to $\mathcal{O}(LN^4)$.

Nevertheless, the iterative search process for determining the optimal number of collaborative clusters continues to result in extended execution latency, mainly attributable to the computational complexity of Algorithm 1. Section Vcompares the algorithm's execution time, illustrating that the iterative search process for the optimal number of collaborative clusters substantially impacts the total execution latency. Thus, we aim to ascertain the optimal number of collaborative clusters with diminished computational complexity. To accomplish this, we put forth a scheme based on deep learning (DL) that encompasses the following primary steps:

- 1) Step 1: Over a specified period, collect samples of the time-varying system states, represented as $\{w_t, h_t \mid t \in \mathcal{I}\}$. In this context, $w_t = [w_1^t, \ldots, w_N^t]^T$, $h_t = [h_1^t, \ldots, h_N^t]^T$, and \mathcal{I} denotes a set of time indices, and $|\mathcal{I}|$ indicates the total number of gathered samples.
- 2) Step 2: Apply Algorithm 2 to ascertain the optimal number of collaborative clusters for each system state sample. These determined numbers of clusters serve as labels for the system state samples, forming a training dataset denoted as $\{w_t, h_t, m_t^* \mid t \in \mathcal{I}\}$.
- Step 3: Build a Deep Neural Network (DNN) with an input layer dimension of 2N, which aligns with the size of {w_t, h_t}, and an output layer dimension of 1, corresponding to m_t^{*}.
- Step 4: Utilize the stochastic gradient descent method to train the DNN with the training dataset {w_t, h_t, m_t^{*} | t ∈ I}. The training process employs the crossentropy loss function. Before commencing training, it is imperative to standardize and scale the data to adhere to a standard normal distribution with a mean of 0 and a variance of 1. This preprocessing step facilitates a swifter training process.
- 5) *Step 5*: Save the trained model and scaling parameters used to normalize the training data. Next, deploy the model to the HAP and import the corresponding parameters.
- 6) Step 6: In the decision-making process, the HAP inputs the acquired system state indicators, denoted as $\{w_{new}, h_{new}\}$, into the DNN model following an identical normalization procedure. Next, employ Algorithm 3 to ascertain the suitable decision. The particulars of Algorithm 3 are expounded upon above.

By leveraging the trained DNN, the iterative process of finding the optimal number of collaborative clusters is skipped. As a result, the deep learning-based scheme achieves a computational complexity of $\mathcal{O}(LN^3)$.

V. NUMERICAL RESULTS

In this section, some simulation results are provided to verify our theoretical results and evaluate the performance of our proposed algorithms. The network scenario shown in Fig. 1 is simulated. In all simulations, we use the parameters of Powercast TX91501-3W with P = 3 Watts for the energy transmitter at the HAP, and those of P2110 Powerharvester for the energy receiver at each IoT device [38]. The energy harvesting efficiency is set as $\eta = 0.51$, and the average channel gain, denoted as \bar{h}_n , adheres to the free-space path



FIGURE 7. Algorithm 3: DL-Based Algorithm for Solving Problem \mathcal{P}_2 .

loss model outlined below:

$$\bar{h}_n = A_d \left(\frac{3 \times 10^8}{4\pi f_c d_n}\right)^{d_e}, n \in \mathcal{N},$$
(30)

with $A_d = 4.11$ denoting the antenna gain, $f_c = 915$ MHz denoting the transmission frequency, and $d_e = 2.8$ denoting the path loss exponent. The time-varying wireless channel gain of the *n*th IoT device at time frame *t* is generated using a Rayleigh fading channel model, expressed as $h_n^t = \bar{h}_n \alpha_n^t$. Here, α_n^t represents the independent random channel fading factor, which follows an exponential distribution with a unit mean.

Without loss of generality, the time block length is set as T = 1 s, the number of IoT devices is set as N = 20, and the minimum data processing requirement within a time frame of each IoT device is set as $l_{th} = 4000$ bit. For local computing, the computation energy efficiency coefficient and the number of CPU cycles needed to process each bit of raw data are



FIGURE 8. The average WSCR of the system versus the average distance from the loT devices to the HAP under different minimum data processing requirements.

set as $k_n = 10^{-26}$ and $\phi_n = 100$, $n \in \mathcal{N}$, respectively [22]. For computing offloading, the bandwidth is set as B = 2MHz and the noise power is set as $N_0 = 10^{-10}$ W. All above parameters will not change unless otherwise stated. Moreover, all the simulations are performed on a Pytorch 1.10.2 platform with an Intel Core i7-13650HX 4.9 GHz CPU and 16 GB of memory.

First, we assess the performance of our two proposed algorithms in comparison to the exhaustive search method. In this method, all $1 + \sum_{m=1}^{m=\lfloor \frac{N}{2} \rfloor} C_N^{2m} (2m-1)!!$ potential scenarios for constructing collaborative clusters are explored to achieve the highest WSCR.

Fig. 8 shows the achievable average WSCR of the system versus the average distance \overline{d} from the IoT devices to the HAP under different minimum data processing requirements within a time frame, i.e., $l_{th} = 4000$ (the upper curve) and $l_{th} = 8000$ (the lower curve). It is evident that as \overline{d} increases, the achievable average WSCR decreases. This is attributed to the fact that greater distances result in poorer channel conditions. Furthermore, an increase in l_{th} leads to a reduction in the achievable average WSCR for all three algorithms. Notably, the performance of our proposed DL-based algorithm closely aligns with that of the priority-based iterative algorithm and the exhaustive search approach. This suggests that our proposed algorithms have the capacity to converge towards the globally optimal solution.

In Fig. 9, we analyze the performance of the proposed priority-based iterative algorithm by plotting the evolution of the achievable average WSCR of the system at each iteration. Here, we set $l_{th} = 4000$ bit, N = 30, $\overline{d} = 4$ m, while keeping the remaining parameters consistent with Fig. 8. It is notable that the iterative search curve does not exhibit a monotonic or concave pattern characterized by an initial increase followed by a decrease. Additionally, the curve reaches its closest value to the exhaustive search result when the number of collaborating clusters is 12. This behavior stems from the influence of the number of collaborative clusters on the optimal time and data allocation scheme,



FIGURE 9. The priority-based iterative algorithm curve versus iteration number.



FIGURE 10. The average operation time versus the number of IoT devices.

resulting in irregular fluctuations in the average WSCR with respect to the number of collaborative clusters. As a consequence, the proposed priority-based iterative method entails an exploration of all conceivable scenarios for the number of collaborating clusters, leading to a substantial increase in execution time.

In Fig. 10, the average operation time of the three algorithms is plotted against the number of IoT devices N, with the distance from the devices to the HAP being randomly generated. It is noteworthy that for relatively small numbers of IoT devices, e.g., N = 5, the average operation time required for the priority-based iterative algorithm closely matches that of the DL-based algorithm. However, as N increases, the operation time required for the priority-based iterative algorithm experiences a significant rise. This may be attributed to the fact that an increased number of IoT devices introduces more iterations in the process of finding the optimal number of collaborating clusters. In contrast, the required operation time for the DL-based algorithm does not exhibit a substantial increase with the growth of N. Furthermore, the average operation time of



FIGURE 11. The average WSCR of the system versus the number of IoT devices.

the two proposed algorithms is much lower than that of the exhaustive search algorithm. Based on the observations in Fig. 10, it is more advisable to employ the priority-based iterative algorithm in networks with relatively small sizes, while the DL-based algorithm is preferred for large-scale networks. This conclusion aligns with our theoretical analysis of the computational complexities of the two algorithms.

Subsequently, we contrast the proposed two algorithms with the following three additional benchmark schemes. Additionally, we include the result of the exhaustive search method (EX) for comparison.

- Local Computing only (LC): In this scheme, all IoT devices' data is processed locally.
- Non-Collaborative Partial Offloading (NC): This scheme does not involve collaboration among IoT devices. The entire system follows the time and data allocation strategy derived from Algorithm 1.
- 3) **Stochastic collaboration (SC):** In this scheme, the collaboration strategy of IoT devices is randomly selected.

Fig. 11 illustrates the achievable average WSCR of the system versus the number of IoT devices N. It is observed that as N increases, the average WSCR decreases for all schemes except the LC strategy. This can be attributed to the gradual decrease in the average computational resources allocated by the HAP to the IoT devices with the increasing number of IoT devices. Additionally, the performance of the SC scheme consistently lags behind that of the NC scheme, underscoring the importance of a meticulously designed collaboration strategy among users. As the number of IoT devices N increases, the gap between the proposed schemes (priority-based iterative algorithm (PI) and DL-based algorithm (DL)) and the NC scheme widens. This underscores the imperative of optimizing collaboration among users in scenarios with a large number of IoT devices, while the benefits of such collaboration are less pronounced in scenarios with a small number of IoT devices. Finally, both of our proposed schemes demonstrate performance that is highly comparable to that of the exhaustive search



FIGURE 12. The average WSCR of the system versus the processor's computing speed of the HAP.



FIGURE 13. The average WSCR of the system versus the path loss exponent de.

algorithm, indicating their capability to achieve nearly optimal performance.

Fig. 12 illustrates the average WSCR of the system versus the processor's computing speed of the HAP f_{hap} . It is evident that as f_{hap} increases, the average WSCR increases for all schemes except the LC strategy. When $f_{hap} > 1$ GHz, the NC scheme surpasses the SC scheme in performance. Additionally, with the increase of f_{hap} , the performance of the NC scheme gradually converges towards that of the proposed schemes. This emphasizes that collaboration among devices becomes less crucial when the computational resources of the HAP are ample. Conversely, when $f_{ap} < 1$ GHz, the SC scheme outperforms the NC scheme.

Fig. 13 displays the average WSCR of the system with respect to the path loss exponent d_e . It is evident that with an increase in d_e , the average WSCR of all schemes experiences a significant decrease. When $d_e < 2.7$, the SC scheme outperforms the NC scheme. However, for $d_e \ge 2.8$, the NC scheme surpasses the SC scheme and gradually approaches the performance level of the proposed schemes. This may be attributed to the fact that, as the channel quality deteriorates, the energy harvested by the devices decreases, while the



FIGURE 14. The optimal number of collaborating clusters *m* obtained by the priority-based iterative algorithm versus various parameters.

energy required for offloading increases. Consequently, the performance enhancement resulting from user collaboration gradually diminishes. Therefore, collaboration among users becomes unnecessary when the channel quality is poor.

Lastly, we examine the influence of various parameters on device collaboration. In particular, we plot the variation in the optimal number of collaborating clusters obtained through the priority-based iterative algorithm under different parameters.

Fig. 14 presents the optimal number of collaborating clusters *m* obtained through the priority-based iterative algorithm versus various parameters. It is evident that with an increase in the energy harvesting efficiency η or the number of IoT devices N, the optimal number of collaborative clusters also increases. Conversely, as the processor's computing speed of the HAP f_{hap} or the minimum data processing requirement within a time frame for each device increases, the optimal number of collaborative clusters decreases. Consequently, in scenarios where energy is the primary limiting factor, collaboration among devices is restrained, while in scenarios where the computational capacity of the HAP is the main limiting factor, collaboration among users is encouraged. Furthermore, an increase in the minimum data processing requirement within a time frame for each device hinders collaboration among users.

Based on the simulation results above, it is evident that the DL-based algorithm achieves near-optimal performance within a short runtime. Hence, the proposed scheme exhibits significant feasibility and advantages for real-time decisionmaking in real-world fading environments. Nevertheless, the scheme possesses potential limitations, such as inadequate adaptation to IoT devices joining or exiting the network. To mitigate this issue, in practical deployment, DNN models corresponding to different numbers of IoT devices can be pre-trained and switched based on the varying number of IoT devices. However, this approach necessitates the upfront collection of extensive data for different IoT device counts, incurring high costs. In contrast, an alternative scheme employs the priority-based iterative algorithm for initial decision-making, saving samples and labels for subsequent use. Upon accumulating sufficient training data for a specific number of IoT devices, the corresponding DNN model is trained. Subsequently, when encountering the specified number of IoT devices, the trained DNN is utilized. Over time, DNN adoption will become the norm.

VI. CONCLUSION

This paper investigates the design for maximizing the weighted sum computation rate in a wireless-powered MEC network with multi-user cooperation. A mathematical optimization problem is formulated by jointly optimizing collaboration, time allocation, and data distribution among multiple IoT devices and the HAP while adhering to constraints on energy causality and devices' computing rate requirements. To efficiently address this mixed-integer programming problem, an interior-point method based algorithm and a priority-based iterative algorithm are initially proposed. Then, a deep learning-based approach is introduced to accelerate the algorithm's operation further. Simulation results demonstrate that the proposed algorithms achieve comparable performance with the exhaustive search method. Furthermore, the deep learning-based approach significantly reduces the algorithm's execution time compared to the priority-based iterative algorithm and the exhaustive search method.

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