

Dealing With Jamming Attacks in Uplink Pairwise NOMA Using Outage Analysis, Smart Relaying, and Redundant Transmissions

VAN-LAN DAO (Member, IEEE), ELISABETH UHLEMANN^{1b} (Senior Member, IEEE),
AND SVETLANA GIRS^{1b} (Member, IEEE),

School of Innovation, Design and Engineering, Mälardalen University, 722 20 Västerås, Sweden

CORRESPONDING AUTHOR: V.-L. DAO (e-mail: van.lan.dao@mdu.se)

This work was supported by the European Union's Horizon 2020 Research and Innovation Programme through the Marie Skłodowska-Curie under Grant 764785 (FORA—Fog Computing for Robotics and Industrial Automation). The work of Elisabeth Uhlemann was supported in part by the Serendipity Project through the Swedish Foundation for Strategic Research and in part by AIDOaRT, the ECSEL Joint Undertaking (JU) under Grant 101007350. The JU receives support from the European Union's Horizon 2020 Research and Innovation Programme and Sweden, Austria, Czech Republic, Finland, France, Italy, Spain.

ABSTRACT This study focuses on optimizing the performance of an uplink pairwise Non-Orthogonal Multiple Access (NOMA) scenario with and without the support of a relay, while subject to jamming attacks. We consider two different relaying protocols, one where the sources and the destination are within range of each other and one where they are not. The relay node can be mobile, e.g., a mobile base station, an unmanned aerial vehicle (UAV) or a stationary node that is chosen as a result of a relay selection procedure. We also benchmark with a NOMA retransmission protocol and an Orthogonal Multiple Access (OMA) scheme without a relay. We analyze, adjust and compare the four protocols for different settings using outage analysis, which is an efficient tool for establishing communication reliability for both individual nodes and the overall wireless network. Closed-form expressions of outage probabilities can be adopted by deep reinforcement learning (RL) algorithms to optimize wireless networks online. Accordingly, we first derive closed-form expressions for the individual outage probability (IOP) of each source node link and the relay link using both pairwise NOMA and OMA. Next, we analyze the IOP for one packet (IOPP) for each source node considering all possible links between the source node to the destination, taking both phases into account for the considered protocols when operating in Nakagami- m fading channels. The overall outage probability for all packets (OOPP) is defined as the maximum IOPP obtained among the source nodes. This metric is useful to optimize the whole wireless network, e.g., to ensure fairness among the source nodes. Then, we propose a method using deep RL where the OOPP is used as a reward function in order to adapt to the dynamic environment associated with jamming attacks. Finally, we discuss valuable guidelines for enhancing the communication reliability of the legitimate system.

INDEX TERMS Dynamic decoding order, imperfect CSI, outage performance, cooperative NOMA, deep reinforcement learning.

I. INTRODUCTION

CONSIDERING the stringent requirements for embedded or cyber-physical systems, Non-Orthogonal Multiple Access (NOMA) has several advantages [1], [2], [3]. It has been shown that the NOMA-based systems can provide more predictable communication than Orthogonal Multiple Access (OMA)-based systems with proper settings for the specific

application [4]. Moreover, performance in terms of outage probability, wireless connectivity, and user fairness can be improved compared to OMA-based systems [3], [5], [6]. In principle, multiple source nodes are served simultaneously using the same time and frequency in uplink NOMA. To separate each source node's signal at the destination, a successive interference cancellation (SIC) unit is used.

However, due to the complexity of the SIC unit and the occurrence of imperfect SIC in practice, having a lower number of active source nodes simultaneously, as in, e.g., pairwise NOMA, is practical [7]. In fact, pairwise NOMA can be deployed on top of existing OMA-based protocols for various applications such as factory automation [4], [8], [9], [10].

Due to the nature of open wireless transmissions, interference from co-located clusters and other wireless networks operating at the same frequency band and within the same area cannot be avoided due to the exponential growth of the number of wireless devices [11], [12]. In the worst case, harmful jamming attacks also aim to interrupt the ongoing transmissions of legitimate communication systems by generating noise signals over relevant wireless channels [13], [14], [15]. For example, a potential reactive jamming attack is reported by the National Institute of Standards and Technology (NIST) in [16]. As a result, taking interference in terms of any jammers and/or interferers into account is needed. Moreover, the jammers may be smart and deploy their own strategies, e.g., power allocation and location to optimize their own systems. Therefore, the legitimate system should also have a proper tool to deal with this situation.

To continue enhancing the communication reliability of pairwise NOMA also in the presence of jamming or strong interference, relaying can be used [17], [18]. In general, a selected relay located between the source node and the destination can help to forward the source node's packet to the destination to increase the probability that the source node's packet is decoded correctly at the destination. In the literature, there are different types of relaying strategies which are applicable in uplink NOMA and in this work we focus on relaying protocols that include two source nodes, one or multiple relays, and a destination. In the first phase, the source nodes transmit their own packets to all relays and to the destination. Thereafter, only one (selected) relay forwards the correctly decoded packet(s) to the destination during the second phase using OMA. Examples of relaying protocols aiming to improve the communication reliability for multiple source nodes communicating with only one destination can be found in [19], [20], [21], [22]. However, achieve the reliability improvement, both the power allocation and the behavior of the relaying protocols, e.g., the selected relay's position, play an important role, especially in the presence of jamming attacks and/or strong interference. But the protocols presented in previous publications, e.g., [19], [20], [21], [22], have not taken jamming attacks into account, and thus an investigation on the effects of jamming attacks and how to deal with them in uplink pairwise NOMA using relaying is needed.

In this work, we consider the scenario as shown in Fig. 1 where a pair of two source nodes are communicating with one destination, aided by a (mobile) relay, e.g., an unmanned aerial vehicle (UAV). The situation when the direct links between the two source nodes and the destination are not

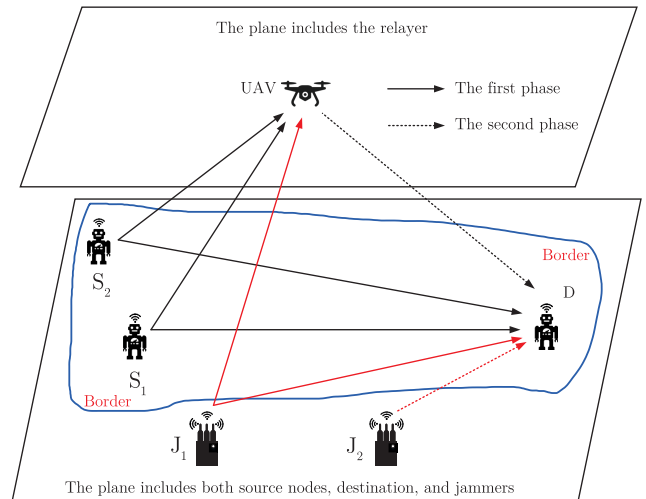


FIGURE 1. System model.

available, we denote conventional NOMA relaying protocol (CNRP). In contrast, when the direct links between both source nodes and the destination exist, the relaying protocol is termed full NOMA relaying protocol (FNRP). We also consider a pure retransmission NOMA protocol (RNP) in which both source nodes re-transmit their packets one more time in uplink pairwise NOMA without help from any relay. To benchmark these three protocols, CNRP, FNRP, RNP, an adaptive power allocation OMA protocol (APAOP) is also considered in which each source node transmits its own packet to the destination in its own phase using adaptive power allocation. We analyze, adjust, and compare the four protocols such that the outage probability is minimized. To evaluate the communication reliability for the whole network, we extend the individual outage probability (IOP) for a single wireless link to the IOP for one packet (IOPP) for each source node taking multiple links into account and then finally the overall outage probability for all packets (OOPP). Accordingly, the main contributions of this work can be summarized as follows:

- Closed-form expressions for the IOP of each source to destination link when using pairwise NOMA or OMA, each source node to the relay link using pairwise NOMA, as well as the relay to destination link using OMA are derived for Nakagami- m fading channels.
- We analyze the IOPP for each source node considering multiple links between the two source nodes and the destination, between both source nodes and the relay, and between the relay and the destination taking both phases into account for the considered protocols.
- The OOPP is defined as the maximum IOPP obtained among the source nodes. This metric is useful to optimize the entire wireless network, e.g., to ensure the fairness condition among the source nodes.
- Based on the OOPP, we propose a method to use deep reinforcement learning (RL) with OOPP as the reward

function to adapt to the dynamic environment associated with jamming attacks.

The rest of this paper is organized as follows: The related works are described in Section II, followed by the system model and a description of the four protocols, with and without a relay in Section III. Next, the considered protocols are introduced in Section IV. After that, Section V presents the calculation of the IOP for one link, the IOPP for each source node, and OOPP in the uplink pairwise NOMA scenario. Next, the deep RL architecture is introduced in Section VI. After that, numerical results are presented in Section VII before providing some general guidelines. Finally, Section VIII concludes the paper.

II. RELATED WORKS

In [23], a relaying scheme for uplink NOMA is proposed with a fixed decoding order (FDO) scheme to enable decoding via linear combining without performing SIC at the destination. Closed-form expressions of the outage probability are derived in Nakagami- m fading channels for both source nodes. However, dynamic decoding order (DDO) is shown as a way to improve the communication reliability [10], [15], [24]. Therefore, a DDO should be adopted. A spectrally-efficient cooperative relaying protocol is presented for use in a scenario with multiple relayers [19]. The chosen relay is decided based on the maximum channel gain criterion among channel gains between relayers and the destination. A distributed energy efficiency maximization strategy is considered using a game theoretic approach [25]. Several amplify and forward (AF) schemes are investigated in [26], [27], [28], in which the maximum signal-to-interference-plus-noise ratio (SINR) is used to select the best relay in [28]. In [20], a UAV acting as a relay in a disaster area is proposed to provide uplink relaying services, and the system throughput is also evaluated. To improve the achievable effective capacity for the machine type communication application, an uplink NOMA and buffer-aided relaying are proposed to assist the finite blocklength transmission [29], [30], [31], while multiple relayers scenario is considered using various Relay Selection (RS) strategies in [32]. A massive machine-type communication scheme is also investigated in [33]. In [34], [35], ergodic capacity and outage performance are presented and discussed for an uplink full-duplex cooperative NOMA system. A multi-hop cooperative NOMA scheme is introduced to enhance system energy efficiency and reliability [36]. In [37], the authors considered cooperative techniques good for establishing pattern division multiple access and evaluated the system by using outage probability and system throughput. An energy-efficient secure short packet transmission of cooperative NOMA is proposed to assist massive machine-type communication application [38]. In [39], a max-min SINR criterion is used to select the best relay in uplink NOMA AF scheme and then the outage probability and throughput are evaluated. A partial decode and amplify NOMA scheme for uplink cooperative

short packet communication is investigated by using the closed-form expression of average block error rate [21]. In [22], the outage probability of the end-to-end uplink and downlink cooperative NOMA is evaluated over Nakagami- m fading channels. In [40], the effects of the relay's position and power allocation on the ergodic rate are analyzed. In [41], an uplink NOMA-based hybrid satellite-terrestrial relay network is proposed and evaluated by using the outage performance. However, the effects of power allocation for uplink NOMA and relay's position are still necessary to be investigated in the presence of jamming attacks. Note that the relay position can be mobile either due to it being a mobile access point, a UAV or the result of a selection procedure among multiple fixed location relayers. In addition, jamming attacks can be smart to change their strategies, e.g., location, transmit power, etc., to defeat legitimate systems. Therefore, to deal with these situations, deep RL is a good approach to enhance the communication reliability of legitimate systems [42], [43].

III. SYSTEM MODEL

In this work, we consider a system consisting of two source nodes \mathcal{S}_s , $s \in \{1, 2\}$ communicating with a destination \mathcal{D} aided by a mobile access point (AP), \mathcal{R} , e.g., UAV, acting as a relay in the presence of jamming attacks, Fig. 1. For the CNRP, FNRP, and RNP, during the first phase, both source nodes transmit their packets to the relay and the destination in uplink pairwise NOMA, while only jammer \mathcal{J}_1 is active to generate jamming signal over all channels to attack all relayers and the destination. In the second phase of the CNRP and FNRP, the relay can help to forward the received packet(s) from the previous slot to the destination in the presence of a jamming attack from \mathcal{J}_2 . For the RNP, both source nodes retransmit their packets in the second phase. For the APAOP, each source node is active at each phase. Note that both jammers are mobile and smart to change their transmit power level and position to defeat the legitimate system, while still optimizing their own systems. We assume that the length of each phase is equal to one another and only one jammer is active at each phase. Note that all legitimate devices are located inside the border and are protected by fences or walls. Consequently, all jammers are only allowed to stay outside of the border. Here, the UAV acting as a relay is located at the altitude h compared to the plane consisting of both source nodes, the destination, and jammers. This is to have higher chances of line of sight (LoS) between the relay and the legitimate devices on the ground, which can enhance communication reliability. Channels between \mathcal{D} and \mathcal{S}_s , between \mathcal{R} and \mathcal{S}_s , between \mathcal{R} and \mathcal{D} , between \mathcal{J}_1 and \mathcal{D} , between \mathcal{R} and \mathcal{J}_1 , and between \mathcal{J}_2 and \mathcal{D} are $\tilde{g}_s^{SD} = \frac{g_s^{SD}}{\sqrt{1+d_s^{SD} \zeta_s^{SD}}}$, $\tilde{g}_s^{SR} = \frac{g_s^{SR}}{\sqrt{1+d_s^{SR} \zeta_s^{SR}}}$, $\tilde{g}^{RD} = \frac{g^{RD}}{\sqrt{1+d^{RD} \zeta^{RD}}}$, $\tilde{g}_{\mathcal{J}_1} = \frac{g_{\mathcal{J}_1}}{\sqrt{1+d_{\mathcal{J}_1} \zeta_{\mathcal{J}_1}}}$, $\tilde{g}_{\mathcal{J}_1 \mathcal{R}} = \frac{g_{\mathcal{J}_1 \mathcal{R}}}{\sqrt{1+d_{\mathcal{J}_1 \mathcal{R}} \zeta_{\mathcal{J}_1 \mathcal{R}}}}$, and $\tilde{g}_{\mathcal{J}_2} = \frac{g_{\mathcal{J}_2}}{\sqrt{1+d_{\mathcal{J}_2} \zeta_{\mathcal{J}_2}}}$ [15], respectively.

The channel coefficients g_s^{SD} , g_s^{SR} , g^{RD} , $g_{\mathcal{J}_1}$, $g_{\mathcal{J}_1\mathcal{R}}$, and $g_{\mathcal{J}_2}$ are assumed to be Nakagami- m fading, modeling a large number of wireless channels by adjusting its parameters, e.g., Rayleigh fading with $m = 1$, Rician fading with parameter K when $m = \frac{(K+1)^2}{2K+1}$ [44]. d_s^{SD} and ζ_s^{SD} , d_s^{SR} and ζ_s^{SR} , d^{RD} and ζ^{RD} , $d_{\mathcal{J}_1}$ and $\zeta_{\mathcal{J}_1}$, $d_{\mathcal{J}_1\mathcal{R}}$ and $\zeta_{\mathcal{J}_1\mathcal{R}}$, and $d_{\mathcal{J}_2}$ and $\zeta_{\mathcal{J}_2}$ are the distances and path-loss exponents between \mathcal{D} and \mathcal{S}_s , between \mathcal{R} and \mathcal{S}_s , between \mathcal{R} and \mathcal{D} , between \mathcal{J}_1 and \mathcal{D} , between \mathcal{R} and \mathcal{J}_1 , and between \mathcal{J}_2 and \mathcal{D} , respectively. We also assume that all devices operate in half-duplex mode with a single antenna. Moreover, the relay, \mathcal{S}_s , \mathcal{D} , \mathcal{J}_1 , and \mathcal{J}_2 are located at $(x^{\mathcal{R}}, y^{\mathcal{R}}, h)$, $(x_s^{\mathcal{S}}, y_s^{\mathcal{S}}, 0)$, $(x_{\mathcal{D}}, y_{\mathcal{D}}, 0)$, $(x_{\mathcal{J}_1}, y_{\mathcal{J}_1}, 0)$, and $(x_{\mathcal{J}_2}, y_{\mathcal{J}_2}, 0)$, respectively. The distances d_s^{SD} , d_s^{SR} , d^{RD} , $d_{\mathcal{J}_1}$, $d_{\mathcal{J}_1\mathcal{R}}$, and $d_{\mathcal{J}_2}$ can be expressed as follows:

$$d_s^{SD} = \sqrt{(x_s^{\mathcal{S}} - x_{\mathcal{D}})^2 + (y_s^{\mathcal{S}} - y_{\mathcal{D}})^2}, \quad (1)$$

$$d_s^{SR} = \sqrt{h^2 + (x^{\mathcal{R}} - x_s^{\mathcal{S}})^2 + (y^{\mathcal{R}} - y_s^{\mathcal{S}})^2}, \quad (2)$$

$$d^{RD} = \sqrt{(x^{\mathcal{R}} - x_{\mathcal{D}})^2 + (y^{\mathcal{R}} - y_{\mathcal{D}})^2}, \quad (3)$$

$$d_{\mathcal{J}_1} = \sqrt{(x_{\mathcal{J}_1} - x_{\mathcal{D}})^2 + (y_{\mathcal{J}_1} - y_{\mathcal{D}})^2}, \quad (4)$$

$$d_{\mathcal{J}_1\mathcal{R}} = \sqrt{h^2 + (x_{\mathcal{J}_1} - x^{\mathcal{R}})^2 + (y_{\mathcal{J}_1} - y^{\mathcal{R}})^2}, \quad (5)$$

$$d_{\mathcal{J}_2} = \sqrt{(x_{\mathcal{J}_2} - x_{\mathcal{D}})^2 + (y_{\mathcal{J}_2} - y_{\mathcal{D}})^2}. \quad (6)$$

In practice, perfect channel state information (CSI) is not available at the receiver(s). Therefore, imperfect CSI is considered in this work. The channel coefficients, g_x^y , $(x, y) \in \{(s, SD), (s, SR), (\mathcal{R}, D)\}$, using linear minimum mean square error are expressed as $g_x^y = \hat{g}_x^y + e_x^y$. Consequently, the channels, \tilde{g}_x^y , can be given as

$$\tilde{g}_x^y = \frac{\hat{g}_x^y + e_x^y}{\sqrt{1 + d_x^y \zeta_x^y}}, \quad (7)$$

where \hat{g}_x^y and $e_x^y \sim CN(0, \sigma_x^y)$ are, respectively, the estimated channel coefficient and channel estimation error and \hat{g}_x^y and e_x^y are uncorrelated. Moreover, all channels follow Nakagami- m fading, therefore channel gains $|\hat{g}_x^y|^2$, $|g_{\mathcal{J}_1}|^2$, $|g_{\mathcal{J}_1\mathcal{R}}|^2$, and $|g_{\mathcal{J}_2}|^2$ can also be characterized by a Gamma distribution with unit mean and shape m_x^y , $m_{\mathcal{J}_1}$, $m_{\mathcal{J}_1\mathcal{R}}$, and $m_{\mathcal{J}_2}$, respectively. In this work, the channel estimation errors are considered to be fixed and independent compared to the average SINR.

For the CNRP and FNRP, the received signals at \mathcal{R} in the first phase and \mathcal{D} in both phases can be represented as follows:

$$y_{\mathcal{R}}^1 = \sum_{s=1}^2 \tilde{g}_s^{SR} \sqrt{P_s} z_s + \frac{g_{\mathcal{J}_1\mathcal{R}}}{\sqrt{1 + d_{\mathcal{J}_1\mathcal{R}}^{\zeta_{\mathcal{J}_1\mathcal{R}}}}} \sqrt{P_{\mathcal{J}_1}} z_{\mathcal{J}_1} + n_{\mathcal{R}}, \quad (8)$$

$$y_{\mathcal{D}}^1 = \sum_{s=1}^2 \tilde{g}_s^{SD} \sqrt{P_s} z_s + \frac{g_{\mathcal{J}_1}}{\sqrt{1 + d_{\mathcal{J}_1}^{\zeta_{\mathcal{J}_1}}}} \sqrt{P_{\mathcal{J}_1}} z_{\mathcal{J}_1} + n_{\mathcal{D}}^1, \quad (9)$$

$$y_{\mathcal{D}}^2 = \sqrt{P_{\mathcal{R}}} z_{\mathcal{R}} + \frac{g_{\mathcal{J}_2}}{\sqrt{1 + d_{\mathcal{J}_2}^{\zeta_{\mathcal{J}_2}}}} \sqrt{P_{\mathcal{J}_2}} z_{\mathcal{J}_2} + n_{\mathcal{D}}^2, \quad (10)$$

where P , $P_{\mathcal{R}}$, $P_{\mathcal{J}_1}$, $P_{\mathcal{J}_2}$, P_s , z_s , $z_{\mathcal{R}}$, $z_{\mathcal{J}_1}$, $z_{\mathcal{J}_2}$, $n_{\mathcal{R}}$, $n_{\mathcal{D}}^1$, and $n_{\mathcal{D}}^2$ are the total transmit power of both source nodes, transmit power of the relay, transmit power of the jammer 1 and 2, power allocation level for each source node \mathcal{S}_s , uplink signal of \mathcal{S}_s , the signal of \mathcal{R} , noise signal of the jamming attacks from jammers 1 and 2, and additive white Gaussian noise at the relay and the destination for both phases modeled as $n_{\mathcal{R}} \sim CN(0, \sigma_{\mathcal{R}0}^2)$, $n_{\mathcal{D}}^1 \sim CN(0, \sigma_{\mathcal{D}1.0}^2)$, and $n_{\mathcal{D}}^2 \sim CN(0, \sigma_{\mathcal{D}2.0}^2)$ respectively. Note that the total transmit power level of both source nodes is equal to P , $P_1 + P_2 = P$.

As mentioned above, DDO can help to enhance the communication reliability of the legitimate system [10], [15], [24]. This is why a DDO scheme is considered for the CNRP, FNRP, and RNP. All relays and the destination use the estimated channel coefficients to decide on the decoding order. Define $h_s^{SD} = \rho_s^{SD} |\hat{g}_s^{SD}|^2$, $\rho_s^{SD} = \frac{P}{W \sigma_{\mathcal{D}1.0}^2 (1 + d_s^{SD} \zeta_s^{SD})}$, $h_s^{SR} = \rho_s^{SR} |\hat{g}_s^{SR}|^2$, $\rho_s^{SR} = \frac{P}{W \sigma_{\mathcal{R}0}^2 (1 + d_s^{SR} \zeta_s^{SR})}$, $h^{\mathcal{RD}} = \rho^{\mathcal{RD}} |\hat{g}^{\mathcal{RD}}|^2$, $\rho^{\mathcal{RD}} = \frac{P_{\mathcal{R}}}{W \sigma_{\mathcal{D}2.0}^2 (1 + d^{\mathcal{RD}} \zeta^{\mathcal{RD}})}$, $h_{\mathcal{J}_1} = \rho_{\mathcal{J}_1} |g_{\mathcal{J}_1}|^2$, $\rho_{\mathcal{J}_1} = \frac{P_{\mathcal{J}_1}}{W \sigma_{\mathcal{D}1.0}^2 (1 + d_{\mathcal{J}_1}^{\zeta_{\mathcal{J}_1}})}$, $h_{\mathcal{J}_1\mathcal{R}} = \rho_{\mathcal{J}_1\mathcal{R}} |g_{\mathcal{J}_1\mathcal{R}}|^2$, $\rho_{\mathcal{J}_1\mathcal{R}} = \frac{P_{\mathcal{J}_1}}{W \sigma_{\mathcal{R}0}^2 (1 + d_{\mathcal{J}_1\mathcal{R}}^{\zeta_{\mathcal{J}_1\mathcal{R}}})}$, $h_{\mathcal{J}_2} = \rho_{\mathcal{J}_2} |g_{\mathcal{J}_2}|^2$, $\rho_{\mathcal{J}_2} = \frac{P_{\mathcal{J}_2}}{W \sigma_{\mathcal{D}2.0}^2 (1 + d_{\mathcal{J}_2}^{\zeta_{\mathcal{J}_2}})}$, where W is the system

bandwidth. In the first phase, when $h_1^{SR} \geq h_2^{SR}$ and $h_1^{SD} \geq h_2^{SD}$, \mathcal{S}_1 's signal is decoded directly by considering both \mathcal{S}_2 's signal and \mathcal{J}_1 's signal as interference and then subtracted by SIC from the received signal $y_{\mathcal{R}}^1$ and $y_{\mathcal{D}}^1$, respectively, before decoding \mathcal{S}_2 's signal treating \mathcal{J}_1 's signal as interference. Accordingly, the received SINRs at the relay \mathcal{R} and the destination to decode z_1 and z_2 can be represented as follows:

$$\gamma_1^{\mathcal{R}} = \frac{\mu_1 h_1^{SR}}{\mu_2 h_2^{SR} + h_{\mathcal{J}_1\mathcal{R}} + a_0^{SR}}, \quad (11)$$

$$\gamma_2^{\mathcal{R}} = \frac{\mu_2 h_2^{SR}}{h_{\mathcal{J}_1\mathcal{R}} + a_0^{SR}}, \quad (12)$$

$$\gamma_1^{\mathcal{D}} = \frac{\mu_1 h_1^{SD}}{\mu_2 h_2^{SD} + h_{\mathcal{J}_1} + a_0^{SD}}, \quad (13)$$

$$\gamma_2^{\mathcal{D}} = \frac{\mu_2 h_2^{SD}}{h_{\mathcal{J}_1} + a_0^{SD}}, \quad (14)$$

where $a_0^{SR} = \mu_1 \sigma_1^{SR2} \rho_1^{SR} + \mu_2 \sigma_2^{SR2} \rho_2^{SR} + 1$, $a_0^{SD} = \mu_1 \sigma_1^{SD2} \rho_1^{SD} + \mu_2 \sigma_2^{SD2} \rho_2^{SD} + 1$, $0 < \mu_s < 1$ is the power allocation factor for \mathcal{S}_s , $\mu_1 + \mu_2 = 1$. In contrast, with $h_1^{SR} < h_2^{SR}$ and $h_1^{SD} < h_2^{SD}$, \mathcal{S}_2 's signal is decoded first before decoding \mathcal{S}_1 's signal, thus the received SINRs at

TABLE 1. The principle of CNRP and FNRP.

Phase	Behavior
1	<ul style="list-style-type: none"> • \mathcal{S}_1 and \mathcal{S}_2 transmit in uplink pairwise NOMA. • CNRP: Only the relay \mathcal{R} is active in receiver mode. The destination is not active, e.g. in deep sleep mode to save power consumption. • FNRP: Both the relay \mathcal{R} and the destination are active in receiver mode.
2	<ul style="list-style-type: none"> • If \mathcal{R} has received at least one correct packet from the previous phase, \mathcal{R} is active using OMA. When receiving one source packet, \mathcal{R} sends that packet to the destination. If both packets have been received correctly, the relay aggregates these two packets into one before sending it to the destination. If no packets have been received successfully, \mathcal{R} stays quiet. • Only the destination is active in receiver mode. \mathcal{R} is not active.

the relay and the destination to decode z_2 and z_1 can be formulated as

$$\gamma_2^{\mathcal{R}'} = \frac{\mu'_2 h_2^{S\mathcal{R}}}{\mu'_1 h_1^{S\mathcal{R}} + h_{\mathcal{J}_1\mathcal{R}} + b_0^{S\mathcal{R}}}, \quad (15)$$

$$\gamma_1^{\mathcal{R}'} = \frac{\mu'_1 h_1^{S\mathcal{R}}}{h_{\mathcal{J}_1\mathcal{R}} + b_0^{S\mathcal{R}}}, \quad (16)$$

$$\gamma_2^{\mathcal{D}'} = \frac{\mu'_2 h_2^{S\mathcal{D}}}{\mu'_1 h_1^{S\mathcal{D}} + h_{\mathcal{J}_1} + b_0^{S\mathcal{D}}}, \quad (17)$$

$$\gamma_1^{\mathcal{D}'} = \frac{\mu'_1 h_1^{S\mathcal{D}}}{h_{\mathcal{J}_1} + b_0^{S\mathcal{D}}}, \quad (18)$$

where $b_0^{S\mathcal{R}} = \mu'_1 \sigma_1^{S\mathcal{R}2} \rho_1^{S\mathcal{R}} + \mu'_2 \sigma_2^{S\mathcal{R}2} \rho_2^{S\mathcal{R}} + 1$, $b_0^{S\mathcal{D}} = \mu'_1 \sigma_1^{S\mathcal{D}2} \rho_1^{S\mathcal{D}} + \mu'_2 \sigma_2^{S\mathcal{D}2} \rho_2^{S\mathcal{D}} + 1$, $0 < \mu'_s < 1$ is the power level for \mathcal{S}_s , $\mu'_1 + \mu'_2 = 1$. Here, (μ'_1, μ'_2) and (μ_1, μ_2) do not need to be the same values.

In the second phase of the CNRP and FNRP, when \mathcal{R} is active, the received SINR at the destination can be represented as

$$\gamma^{\mathcal{R}\mathcal{D}} = \frac{h^{\mathcal{R}\mathcal{D}}}{h_{\mathcal{J}_2} + \sigma^{\mathcal{R}\mathcal{D}2} \rho^{\mathcal{R}\mathcal{D}} + 1}, \quad (19)$$

IV. PROTOCOLS

In this paper, we investigate four protocols: two relaying protocols in uplink pairwise NOMA, a re-transmission scheme using uplink pairwise NOMA, and an APAOP using OMA as follows.

A. CONVENTIONAL NOMA RELAYING PROTOCOL (CNRP)

In the literature, a wide range of relaying protocols has been proposed for various applications [45]. The main principle of many relaying protocols is presented in Table 1 and we name them CNRP when combined with NOMA [36], [46]. The previous studies usually assume that severely bad channels happen to direct links between the source nodes and the destination. This is why any direct links between the source

nodes and the destination are ignored. When CNRP is adopted, several receivers can operate in a deep sleep mode in different phases, saving power consumption for the wireless devices. This protocol can adapt to adjust power allocation factors in the first phase and the relay's position to improve the communication reliability for the legitimate system.

B. FULL NOMA RELAYING PROTOCOL (FNRP)

In practice, various applications such as quarrying, mining in construction sites still have good direct links between the source nodes and the destination and thus direct links cannot be ignored. To take all chances into account, indirect links between the two source nodes and the destination should be considered. We name the protocol including direct links between the source nodes and the destination Full NOMA Relaying Protocol (FNRP). The main idea of FNRP is that all receivers including the destination take a chance to decode the transmitted packet(s) from the first phase in Table 1 when operating in a receiver mode. Note that the destination is not active during the first phase for the CNRP. Therefore, the FNRP is useful when improving the communication reliability compared to the CNRP. Moreover, at the receivers, a maximal-ratio-combining (MRC) scheme can be deployed to improve the communication reliability but it is considered as a future work. In addition, when the destination can decode correctly one or two packets in the first phase, a feedback signal to the relay and the two source nodes is useful. However, we consider the worst case in the presence of jamming attacks that all feedback are dropped. Similar to the CNRP, both power allocation and relay's position are strategies to deal with jammers.

C. RETRANSMISSION NOMA PROTOCOL (RNP)

In the case without the relay, the two source nodes just transmit and retransmit their packets in uplink pairwise NOMA and we call this protocol as RNP. Accordingly, during the second phase, the received signal and SINRs at the destination are similar to (9), (13), (14), (17), and (18) but different power allocation factors. With RNP, power allocation factors for the two source nodes in both phases are useful strategies to improve the communication reliability.

D. ADAPTIVE POWER ALLOCATION OMA PROTOCOL (APAOP)

When there is no relay, we also consider Adaptive Power Allocation OMA Protocol (APAOP), in which \mathcal{S}_1 is active in the first phase with transmit power level of $\mu_1 P$ and \mathcal{S}_2 transmits its own packet in the second phase with transmit power level of $(1 - \mu_1)P$. Accordingly, the received SINRs at the destination in both phase can be represented as

$$\gamma_1^{OMA} = \frac{\mu_1 h_1^{S\mathcal{D}}}{h_{\mathcal{J}_1} + a_0^{OMA}}, \quad (20)$$

$$\gamma_1^{OMA} = \frac{\mu_2 h_2^{S\mathcal{D}}}{h_{\mathcal{J}_1} + b_0^{OMA}}, \quad (21)$$

where $a_0^{OMA} = \mu_1 \sigma_1^{SD^2} \rho_1^{SD} + 1$ and $b_0^{OMA} = \mu_2 \sigma_2^{SD^2} \rho_2^{SD} + 1$. To be fair for comparison among the four protocols, the total transmit power of both source nodes in this protocol are twice compared to the previous protocols. Similar to the RNP, only power allocation is used to cope with jammers for this protocol.

V. OUTAGE PERFORMANCE ANALYSIS

In this section, we first derive the closed-form expressions of the IOP for each source node at all legitimate receivers using pairwise NOMA in the first phase. Next, the closed-form expressions of the IOP at the destination are derived when using OMA. Finally, the IOPP for each source node considering both phases is presented for all four schemes.

A. THE IOP OF EACH SOURCE NODE AT THE RELAYER AND THE DESTINATION IN THE FIRST PHASE

In this subsection, we derive the IOP for only one source node link to the relayer and the destination using uplink pairwise NOMA. First, we analyze and then derive the IOP of each source node at the destination and the relayer. The S_1 's signal cannot be decoded successfully at \mathcal{R} and \mathcal{D} when either of the following three disjoint cases occurs: (i) \mathcal{R} and \mathcal{D} fail to decode S_1 's signal correctly by considering S_2 's signal and \mathcal{J}_1 's signal as interference when $h_1^{SR/SD} \geq h_2^{SR/SD}$; (ii) S_2 's signal cannot be decoded correctly by considering both signals from S_1 and \mathcal{J}_1 as interference when $h_1^{SR/SD} < h_2^{SR/SD}$; (iii) S_2 's signal is decoded successfully and removed by SIC when $h_1^{SR/SD} < h_2^{SR/SD}$, but \mathcal{R} and \mathcal{D} are still unable to decode S_1 's signal. It is the same for the S_2 signal. To be convenient for writing, we just remove the subscript $\mathcal{SR}/\mathcal{SD}$ in (11)–(18). To be convenient for use, we use h_s instead of $h_s^{SR/SD}$. Accordingly, the IOPs of S_1 and S_2 can be expressed as follows [15]:

$$p_1^{SR/SD} = 1 - I_1 - I_2, \quad (22)$$

$$p_2^{SR/SD} = 1 - I_3 - I_4, \quad (23)$$

in which I_1 is calculated as

$$I_1 = \Pr\{(\gamma_1 \geq A_1) \cap (h_1 \geq h_2)\} \\ = \begin{cases} I_{1a} & \mu_2 A_1 \geq \mu_1 \\ I_{1b} & \text{otherwise} \end{cases}, \quad (24)$$

$$I_3 = \Pr\{(\gamma_2' \geq A_2) \cap (h_1 < h_2)\} \\ = \begin{cases} I_{3a} & \mu_1' A_2 \geq \mu_2' \\ I_{3b} & \text{otherwise} \end{cases}, \quad (25)$$

where

$$I_{1a} = \Pr\{h_1 \geq a_1 h_2 + a_2 h_{\mathcal{J}_1} + a_3\}, \quad (26)$$

$$I_{1b} = \Pr\{(h_1 \geq a_1 h_2 + a_2 h_{\mathcal{J}_1} + a_3) \cap (h_1 \geq h_2)\}, \quad (27)$$

$$I_{3a} = \Pr\{h_2 \geq b_1 h_1 + b_2 h_{\mathcal{J}_1} + b_3\}, \quad (28)$$

$$I_{3b} = \Pr\{(h_2 \geq b_1 h_1 + b_2 h_{\mathcal{J}_1} + b_3) \cap (h_2 > h_1)\}, \quad (29)$$

$$I_2 = \Pr\{(\gamma_1' \geq A_1) \cap (\gamma_2' \geq A_2) \cap (h_1 < h_2)\} \\ = \Pr\left\{\left(\frac{\mu_1' h_1}{h_{\mathcal{J}_1} + b_0} \geq A_1\right) \cap (h_1 < h_2)\right\}$$

$$\cap \left(\frac{\mu_2' h_2}{\mu_1' h_1 + h_{\mathcal{J}_1} + b_0} \geq A_2\right)\}, \quad (30)$$

$$I_4 = \Pr\{(\gamma_1 \geq A_1) \cap (\gamma_2 \geq A_2) \cap (h_1 \geq h_2)\}$$

$$= \Pr\left\{\left(\frac{\mu_1 h_1}{\mu_2 h_2 + h_{\mathcal{J}_1} + a_0} \geq A_1\right) \cap (h_1 \geq h_2)\right. \\ \left. \cap \left(\frac{\mu_2 h_2}{h_{\mathcal{J}_1} + a_0} \geq A_2\right)\right\}, \quad (31)$$

where A_1 and A_2 are the SINR thresholds to decode correctly the S_s 's signal at the destination and \mathcal{R} , respectively. Taking all possible cases into account, we can re-write the probabilities I_2 and I_4 as follows:

$$I_2 = \begin{cases} I_{20} & \mu_1' A_2 \geq \mu_2' \\ I_{21} + I_{22} & \text{otherwise} \end{cases}, \quad (32)$$

$$I_4 = \begin{cases} I_{40} & \mu_2 A_1 \geq \mu_1 \\ I_{41} + I_{42} & \text{otherwise} \end{cases}, \quad (33)$$

in which I_{40} , I_{41} , I_{42} , I_{20} , I_{21} , and I_{22} are given as

$$I_{40} = \Pr\left\{\begin{aligned} &(h_1 \geq a_1 h_2 + a_2 h_{\mathcal{J}_1} + a_3) \\ &\cap (h_2 \geq a_4 h_{\mathcal{J}_1} + a_5) \end{aligned}\right\}, \quad (34)$$

$$I_{41} = \Pr\left\{\begin{aligned} &(h_2 \leq a_6 h_{\mathcal{J}_1} + a_7) \\ &\cap (h_2 \geq a_4 h_{\mathcal{J}_1} + a_5) \\ &\cap (h_1 \geq a_1 h_2 + a_2 h_{\mathcal{J}_1} + a_3) \end{aligned}\right\}, \quad (35)$$

$$I_{42} = \Pr\left\{\begin{aligned} &(h_1 \geq h_2) \cap (h_2 \geq a_4 h_{\mathcal{J}_1} + a_5) \\ &\cap (h_2 > a_6 h_{\mathcal{J}_1} + a_7) \end{aligned}\right\}, \quad (36)$$

$$I_{20} = \Pr\left\{\begin{aligned} &(h_2 \geq b_1 h_1 + b_2 h_{\mathcal{J}_1} + b_3) \\ &\cap (h_1 \geq b_4 h_{\mathcal{J}_1} + b_5) \end{aligned}\right\}, \quad (37)$$

$$I_{21} = \Pr\left\{\begin{aligned} &(h_2 \geq b_1 h_1 + b_2 h_{\mathcal{J}_1} + b_3) \\ &\cap (h_1 \geq b_4 h_{\mathcal{J}_1} + b_5) \\ &\cap (h_1 \leq b_6 h_{\mathcal{J}_1} + b_7) \end{aligned}\right\}, \quad (38)$$

$$I_{22} = \Pr\left\{\begin{aligned} &(h_2 > h_1) \cap (h_1 \geq b_4 h_{\mathcal{J}_1} + b_5) \\ &\cap (h_1 > b_6 h_{\mathcal{J}_1} + b_7) \end{aligned}\right\}, \quad (39)$$

where $a_1 = \frac{\mu_2 A_1}{\mu_1}$, $a_2 = \frac{A_1}{\mu_1}$, $a_3 = \frac{A_1 a_0}{\mu_1}$, $a_4 = \frac{A_2}{\mu_2}$, $a_5 = \frac{a_0 A_2}{\mu_2}$, $a_6 = \frac{A_1}{\mu_1 - \mu_2 A_1}$, $a_7 = \frac{A_1 a_0}{\mu_1 - \mu_2 A_1}$, $b_1 = \frac{A_2 \mu_1'}{\mu_2'}$, $b_2 = \frac{A_2}{\mu_2'}$, $b_3 = \frac{b_0 A_2}{\mu_2'}$, $b_4 = \frac{A_1}{\mu_1'}$, $b_5 = \frac{b_0 A_1}{\mu_1'}$, $b_6 = \frac{A_2}{\mu_2' - A_2 \mu_1'}$, $b_7 = \frac{b_0 A_2}{\mu_2' - A_2 \mu_1'}$. I_{1a} , I_{1b} , I_{3a} , I_{3b} , I_{40} , I_{41} , I_{42} , I_{20} , I_{21} , and I_{22} are derived in the Lemma 1 in Appendix A.

B. THE IOPS AT THE DESTINATION WHEN USING OMA

The IOP is a reliability metric evaluating a single wireless link. In this subsection, we derive the IOP for each source node link to the destination using OMA and the relayer link to the destination also adopting OMA. The IOP of the relayer link to the destination is derived first. When \mathcal{R} is active, the outage probability that $z_{\mathcal{R}}$'s signal is not decoded correctly at the destination is given as

$$p^{\mathcal{RD}} = \Pr\{\gamma^{\mathcal{RD}} < A_0\}, \\ = 1 - \Pr\{h^{\mathcal{RD}} \geq b h_{\mathcal{J}_2} + c\}, \quad (40)$$

where $b = A_0$, $c = A_0(\sigma^{\mathcal{RD}^2} \rho^{\mathcal{RD}} + 1)$. A_0 is the SINR threshold to decode successfully the transmitted packet from the relayer at the destination.

$p^{\mathcal{RD}}$ is derived in the Lemma 2 in Appendix-B.

When APAOP is deployed, the IOP for each source node link to the destination is determined as

$$\begin{aligned} p_1^{APAOP} &= \Pr\{\gamma_1^{OMA} < A_1\} \\ &= 1 - \Pr\{h_1^{SD} \geq b_{01}h_{\mathcal{J}_1} + c_{01}\}, \end{aligned} \quad (41)$$

$$\begin{aligned} p_2^{APAOP} &= \Pr\{\gamma_2^{OMA} < A_2\} \\ &= 1 - \Pr\{h_2^{SD} \geq b_{02}h_{\mathcal{J}_2} + c_{02}\} \end{aligned} \quad (42)$$

where $b_{01} = \frac{A_1}{\mu_1}$, $c_{01} = \frac{A_1 a_0^{OMA}}{\mu_1}$, $b_{02} = \frac{A_2}{\mu_2}$, and $c_{02} = \frac{A_2 b_0^{OMA}}{\mu_2}$. p_1^{APAOP} and p_2^{APAOP} are derived in the Lemma 3 in Appendix-C.

C. THE IOP FOR ONE PACKET OF EACH SOURCE NODE FOR THE CNRP

The transmitted packet from each source node can travel to the destination via two links consisting of the source node link to the relay and the relay link to the destination. Therefore, we consider the IOPP of each source node taking multiple links into account. With the CNRP, the source nodes' packets cannot be delivered correctly to the destination when the destination and/or the relay have failed to decode the transmitted packets in the second and/or first phases, respectively. In other words, the IOPP of each source node can be derived based on the IOP for single wireless links in Sections V-A and V-B. Accordingly, the IOPP for each source node is given as follows:

$$p_s^{CNRP} = 1 - (1 - p_s^{SR})(1 - p_s^{\mathcal{RD}}). \quad (43)$$

Here, p_s^{SR} and $p_s^{\mathcal{RD}}$ are the IOPs of the s -th source node link to the relay and the IOP of the relay link to the destination, respectively. And $(1 - p_s^{SR})(1 - p_s^{\mathcal{RD}})$ represents the probability that the source node packet- s is decoded successfully at the destination after two phases.

D. THE IOP FOR ONE PACKET OF EACH SOURCE NODE FOR THE FNRP

Compared to the CNRP, the IOPP for each source node for the FNRP considers two paths including three links: the source node link to the destination, the source node link to the relay, and the relay link to the destination. Based on the principle of the FNRP, the source nodes' packets cannot reach the destination when: (i) the destination is failed to decode the source nodes' packets in the first phase, and (ii) the destination has no correct packets following the path from $S_s \rightarrow \mathcal{R} \rightarrow \mathcal{D}$. We can see that the IOPP for each source node is a joint probability of multiple IOPs for different single wireless links in both phases. Moreover, the IOP of each source node for single wireless links in different phases are independent to each other. To this end, the IOPP for each source node can be expressed as

$$p_s^{FNRP} = p_s^{SD} \left[1 - (1 - p_s^{SR})(1 - p^{\mathcal{RD}}) \right], \quad (44)$$

in which p_s^{SD} , p_s^{SR} , and $p^{\mathcal{RD}}$ are the IOPs for single wireless links and obtained in Sections V-A and V-B.

E. THE IOP FOR ONE PACKET OF EACH SOURCE NODE FOR THE RNP

Both source nodes transmit their packets in uplink pairwise NOMA two times with different power allocation factors. The IOPs for each source node link to the destination in different phases are independent to each other. Therefore, the IOPP for each source node with one retransmission more can be calculated as follows:

$$p_s^{RNP} = \prod_{i=1}^2 p_{s,i}^{SD}, \quad (45)$$

where $p_{s,i}^{SD}$ is the IOP for the S_s 's link to the destination at phase- i with power allocation factor $\mu_{1,i}$. This probability is calculated as in (22) and (23).

F. DEFINITION OF OOPP

In the NOMA-based systems, user fairness is also a strict requirement [6], [47], [48]. In which, all users belong to the same class of priority should experience the same quality of service (QoS), e.g., the same communication reliability level. To this end, we define a new metric, namely OOPP, based on the attained IOPP for each source node. Accordingly, the OOPPs for the CNRP, FNRP, RNP, and APAOP are given, respectively as

$$p_{CNRP} = \max(p_1^{CNRP}, p_2^{CNRP}), \quad (46)$$

$$p_{FNRP} = \max(p_1^{FNRP}, p_2^{FNRP}), \quad (47)$$

$$p_{RNP} = \max(p_1^{RNP}, p_2^{RNP}), \quad (48)$$

$$p_{APAOP} = \max(p_1^{APAOP}, p_2^{APAOP}). \quad (49)$$

VI. DEALING WITH DYNAMIC CONDITIONS USING DEEP REINFORCEMENT LEARNING

The legitimate communication system aims to improve the communication reliability by minimizing the OOPP. Accordingly, we formulate three problems as shown in (50), (51), and (52). For both CNRP and FNRP protocols, the constraints (50b), (50c), and (50d) are related to the position of relays inside the border. The constraints on power allocation factors are presented in (50e) and (50f). However, the constraints for the RNP protocols are power allocation factors in different phases in (51b) and (51c). For the APAOP, only power allocation factor is a constraint in (52b).

$$(P1) : \min_{\mu_1, \mu'_1, (x_i^{\mathcal{R}}, y_i^{\mathcal{R}}, h_i)} p_{CNRP/FNRP} \quad (50a)$$

$$\text{s.t. } x_{min}^{\mathcal{R}} \leq x^{\mathcal{R}} \leq x_{max}^{\mathcal{R}} \quad (50b)$$

$$y_{min}^{\mathcal{R}} \leq y^{\mathcal{R}} \leq y_{max}^{\mathcal{R}} \quad (50c)$$

$$h_{min} \leq h \leq h_{max} \quad (50d)$$

$$0 \leq \mu_1 \leq 1 \quad (50e)$$

$$0 \leq \mu'_1 \leq 1 \quad (50f)$$

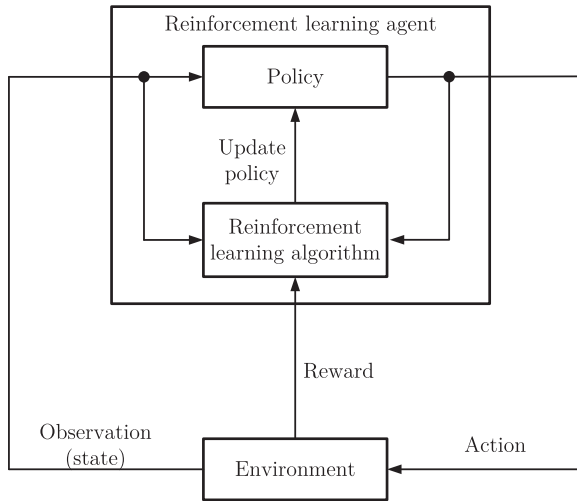


FIGURE 2. Reinforcement learning architecture.

$$(P2) : \min_{\mu_{1,i}, \mu'_{1,i}} p_{RNP} \quad (51a)$$

$$\text{s.t. } 0 < \mu_{1,i} < 1, i \in \{1, 2\} \quad (51b)$$

$$0 < \mu'_{1,i} < 1, i \in \{1, 2\}. \quad (51c)$$

$$(P2): \min_{\mu_1} p_{APAOP} \quad (52a)$$

$$\text{s.t. } 0 < \mu_1 < 1 \quad (52b)$$

In practice, the positions of jamming attacks and their transmit power levels can be changed randomly to defeat the legitimate system. Therefore, the legitimate system should take this into account. In this case, RL is considered as a suitable approach to deal with dynamic conditions as mentioned above [49].

The RL architecture includes an agent and environment, Fig. 2. The main goal of RL is to train an agent to complete a task within an uncertain environment. The RL agent consists of two parts: (i) The policy on how to choose actions based on the states (observations) from the environment. The policy is typically a function approximator with tunable parameters using deep neural networks. (ii) The RL algorithm updates the policy parameters continuously based on actions, states, and reward to find an optimal policy maximizing the cumulative reward received during the task. In this work, a deep deterministic policy gradient (DDPG) agent is adopted to search for an optimal policy maximizing the expected cumulative long-term reward because it supports continuous actions, e.g., power allocation factor μ_1, μ'_1 . A DDPG agent is an actor-critic RL agent, including four function approximators as (a) Actor: The actor takes states as input and returns the corresponding action to maximize the long-term reward. (b) Target actor: This function approximator helps to enhance the stability of the optimization by updating the target actor parameters periodically based on the latest actor parameter values. (c) Critic: The critic takes state and action to return the corresponding expectation of the long-term reward. (d) Target critic: The target critic parameter

is updated periodically based on the latest critic parameter values to enhance the stability of the optimization. In the training phase, the DDPG agent updates both actor and critic properties at each time step and stores past experiences adopting a circular experience buffer. The agent updates both actor and critic by employing a mini-batch of experiences randomly sampled from the buffer. Adopting a stochastic noise model at each training step is to perturb the action selected by the policy. The observation space, the action space, and the reward are defined as follows:

- *State space*: To attack efficiently, the positions of all jammers and their transmit power levels in different slots can be changed. We also consider that the positions of both source nodes can be updated in different slots as well. Therefore, the state space is defined including $(x_{\mathcal{J}_1}, y_{\mathcal{J}_1}, 0), P_{\mathcal{J}_1}, (x_{\mathcal{J}_2}, y_{\mathcal{J}_2}, 0), P_{\mathcal{J}_2}$, and $(x_s^S, y_s^S, 0)$.
- *Action space*: To improve the reliability performance of the legitimate communication system in (50), (51), and (52), the action space includes power allocation factor μ_1, μ'_1 , and (x^R, y^R, h) for both CNRP and FNRP. However, the action space for the RNP only consists of power allocation factors in different phases, μ_1 , and μ'_1 , while only μ_1 is the action for the APAOP. To reduce the complexity of the DDO scheme while enhancing the fairness condition among the two source nodes, we use the DDO-fixed pairwise power allocation (FPPA) scheme, $\mu_1 = \mu'_2, \mu_2 = \mu'_1, \mu_{1,i} = \mu'_{2,i}, \mu_{2,i} = \mu'_{1,i}$ [10], [15]. Accordingly, we only have μ_1 and $\mu_{1,i}$ as actions and their ranges are $0 < \mu_1 < 1$ and $0 < \mu_{1,i} < 1$.
- *Reward*: The legitimate communication system maximizes the communication reliability in terms of minimizing the long term OOPP. Consequently, the reward depends on which protocol is used, $r = -P_{CNRP}/P_{FNRP}/P_{RNP}/P_{APAOP}$ as shown in (46), (47), (48), and (49).

Based on the state space, action space, and reward, the DDPG can be implemented. Each episode is comprised of multiple steps, where the DDPG algorithm follows a sequence: generating an action based on the current state, determining the next state based on the selected action, and then updating the four neural networks of the DDPG agent to facilitate learning.

VII. NUMERICAL RESULTS

In this section, we present numerical results for the IOPP for each source node and the OOPP of the considered system for the four schemes. The following system parameters are used: $W = 1$ Hz, $P = 1$ W, $P_{\mathcal{R}} = 1$ W, $A_1 = A_2 = A_0 = 2^{0.1} - 1$, $h = 20$, $\zeta_s^{SD} = \zeta_s^{SR} = \zeta_s^{RD} = \zeta_{\mathcal{J}_1} = \zeta_{\mathcal{J}_2} = 2$, $m_s^{SD} = m_s^{SR} = m_s^{RD} = m_{\mathcal{J}_1} = m_{\mathcal{J}_2} = 3$, $\sigma_x^2 = 1e - 4$, and $\sigma_{\mathcal{R}0}^2 = \sigma_{\mathcal{D}1}^2 = \sigma_{\mathcal{D}2}^2 = 10^{-10}$ W/Hz [50], [51]. To ensure a fair comparison among all considered protocols, we also consider the total transmit power of the two source nodes for the APAOP is 2W. To check the correctness of the

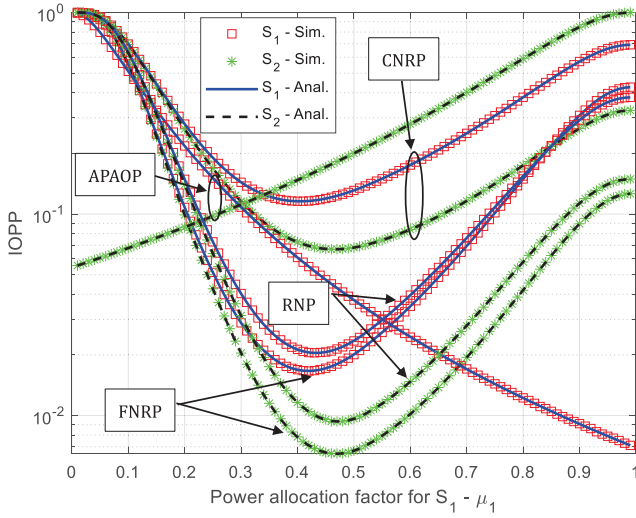


FIGURE 3. The IOPP for each source node taking both phases into account.

analysis in Section V, we also conduct computer simulations using MATLAB. In particular, for each considered IOPP, we first generate 10^7 samples of the channel gains following a Gamma distribution and then check the outage conditions as defined in (22), (23), (40), (41), (42), (43), (44), and (45). The simulation results of the IOPP for each source node are then attained by taking the average of all outage events across 10^7 samples.

To investigate the effect of power allocation on all schemes, we configure the positions and transmit power of legitimate nodes and jammers as follows: $(x^{\mathcal{R}}, y^{\mathcal{R}}, h) = (100, 50, 20)$, $(x_1^{\mathcal{S}}, y_1^{\mathcal{S}}, 0) = (0, 0, 0)$, $(x_2^{\mathcal{S}}, y_2^{\mathcal{S}}, 0) = (25, 100, 0)$, $(x_{\mathcal{D}}, y_{\mathcal{D}}, 0) = (200, 50, 0)$, $P_{\mathcal{J}_1} = 1\text{W}$, $P_{\mathcal{J}_2} = 1\text{W}$, $P_{\mathcal{R}} = 1\text{W}$, $(x_{\mathcal{J}_1}, y_{\mathcal{J}_1}, 0) = (100, -20, 0)$, and $(x_{\mathcal{J}_2}, y_{\mathcal{J}_2}, 0) = (200, -20, 0)$. For the RNP, the power allocation factors for both phases are the same. Fig. 3 illustrates the effect of the power allocation factor on the IOPP for both source nodes using the four protocols mentioned in Section III. We can see that the analytical results and the simulation match very well validating the accuracy of the calculation. Moreover, the power allocation factor affects significantly the IOPP. It can be seen from the figure that the IOPP for both source nodes adopting the FNRP is much smaller than that of using the CNRP. This is because the channels between both source nodes and the destination are good but they are ignored by the CNRP. We also can see that the IOPP for each source node using the APAOP is not fair to each other, e.g., while the IOPP for S_1 's packet decreases dramatically, the IOPP for S_2 's packet increases significantly.

Fig. 4 indicates how power allocation and relay position affect the OOPP of the CNRP when the two source nodes' parameters and both jammers' parameters are the same as investigated in Fig. 3. It is clear that both power allocation and UAV position play a very crucial role to enhance the communication reliability in terms of minimizing the OOPP

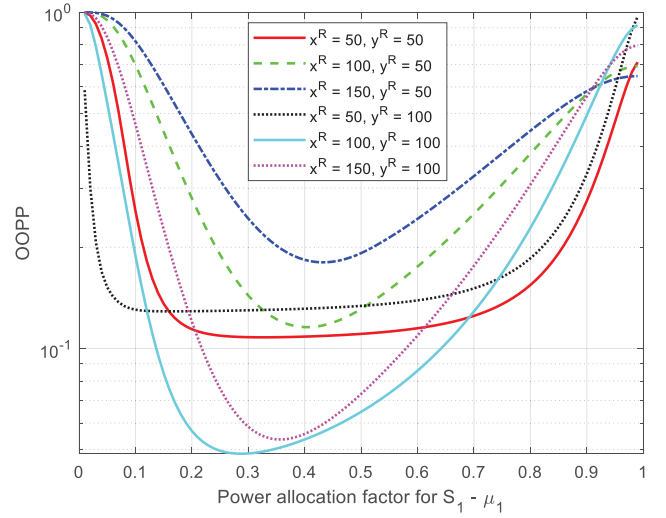


FIGURE 4. The OOPP versus UAV position and power allocation for the CNRP (\mathcal{R} is closer to \mathcal{D} with bigger $x^{\mathcal{R}}$, \mathcal{R} is further away from jammers with bigger $y^{\mathcal{R}}$).

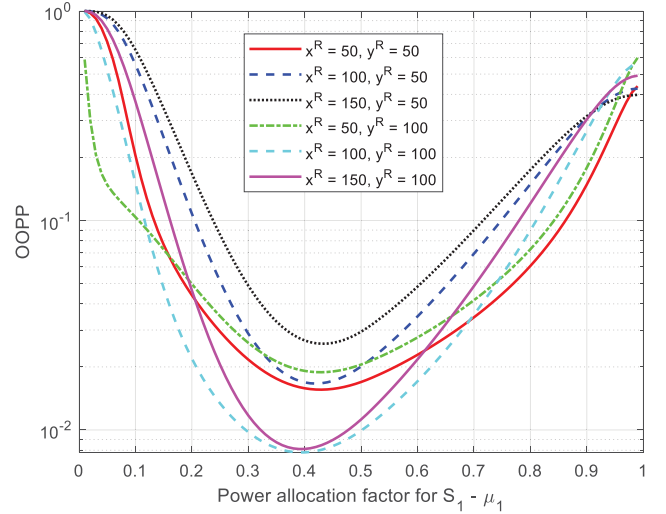


FIGURE 5. The OOPP versus UAV position and power allocation for the FNRP (\mathcal{R} is closer to \mathcal{D} with bigger $x^{\mathcal{R}}$, \mathcal{R} is further away from jammers with bigger $y^{\mathcal{R}}$).

for each source node. We also can see the effects of power allocation and relay position on the OOPP for the FNRP in Fig. 5.

Regarding the RL, three hidden layers including 100 neurons for each layer followed by rectified linear units (ReLU) activation functions are adopted for the actor-network. The activation function for the output layer is the hyperbolic tangent function. For the critic network, both states and actions are considered as inputs. First, all states are fed to a neural network using two hidden layers with 100 neurons for each, and action is fed to another neural network using a hidden layer with 100 neurons. Then, these both neural networks are concatenated before feeding to another neural network with 100 neurons for a hidden layer. The ReLU activation functions are employed for the critic network. Other configurations for the neural networks, such

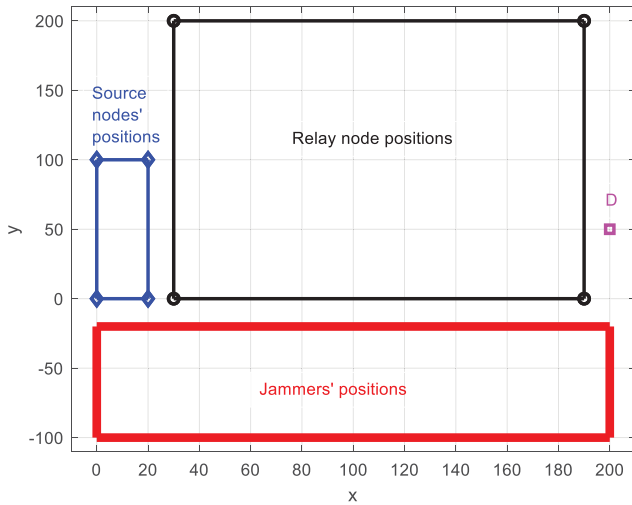


FIGURE 6. Possible positions of the source nodes, relay, and jammers for the deep RL.

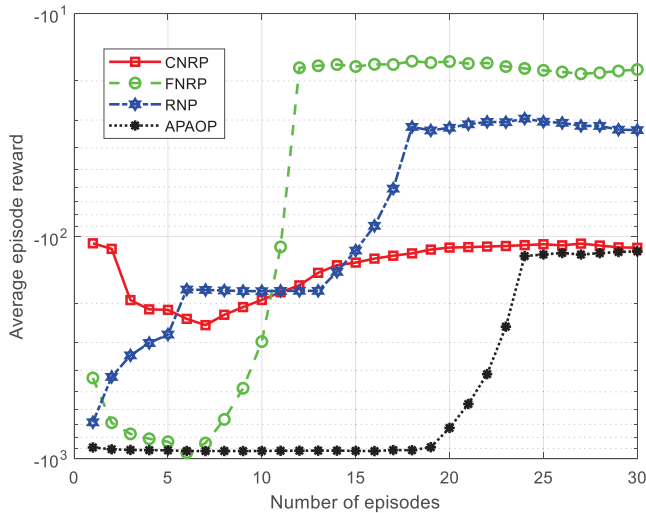


FIGURE 7. Training progress for the four protocols.

as learning rate, gradient threshold, regularization factor, sample time, experience buffer length, minibatch size, and the number of steps per episode are set to $1e-3$, 1 , $1e-4$, 1 , $1e6$, 128 , and $1e3$, respectively. Note that the number of hidden layers, neurons and other hyperparameters of the deep RL architecture have been selected via trial and error to find the best performing deep RL for the considered problem. For the state space of the RL, we configure that $0 \leq x_s^S \leq 20$, $0 \leq y_s^S \leq 100$, $P_{\mathcal{J}_s} \in \{0.1 : 0.1 : 3\}W$, $0 \leq x_{\mathcal{J}_s} \leq 200$, $-100 \leq y_{\mathcal{J}_s} \leq -20$ as shown in Fig. 6. We configure for the action space related to the relay node position as follows: $h = 20$, $30 \leq x^{\mathcal{R}} \leq 190$, and $0 \leq y^{\mathcal{R}} \leq 200$. Both training and inference phases of the RL are implemented using MATLAB on Desktop HP Z2 TWR Base G9, Core i9-12900K 3.20G 30MB 16 cores, 64GB DDR5, NVIDIA GeForce RTX 3080.

Fig. 7 presents the average episode reward in terms of the OOPP versus the number of episodes for the four schemes.

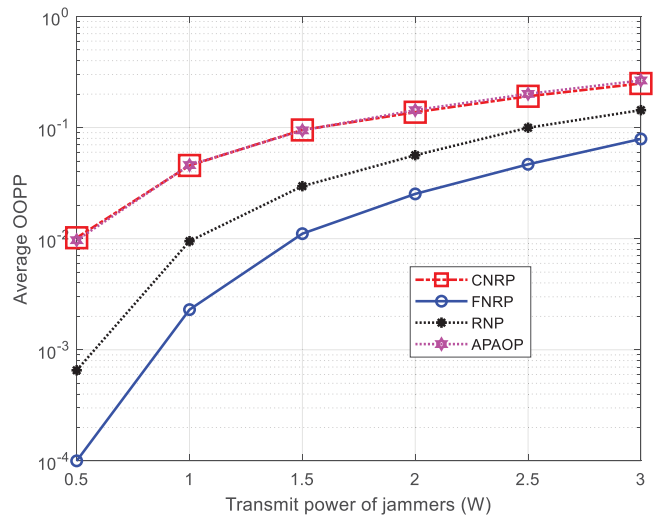


FIGURE 8. Effect of transmit power levels of both jammers on average OOPP in the inference phase when $P_{\mathcal{J}_1} = P_{\mathcal{J}_2}$.

From the figure, we can see that the average episode reward in terms of average OOPP for all schemes converges after 23 episodes. It can be seen from the figure that the FNRP can offer the highest communication reliability, followed by the RNP, while the communication reliability of the APAOP and CNRP schemes is the worst and approximately equal to each other. This phenomenon is also demonstrated in Fig. 3. In the inference phase, we fix the transmit power of both jammers as the same to investigate the average OOPP over 1000 steps. In Fig. 8, a change of the average OOPP following the transmit power of both jammers is provided. In general, the average OOPP grows up significantly when the transmit power of both jammers increases. We also can observe the same trend of the communication reliability offering by each protocol as shown in Fig. 7.

A. CREATING GUIDELINES FOR A SPECIFIC SCENARIO

Based on the obtained results, both FNRP and RNP can offer higher reliability in terms of smaller OOPP compared to the CNRP and APAOP. Finally, the following guidelines can be provided:

- When a mobile AP, e.g., a UAV, is available to act as a relay, the FNRP should be employed to enhance the communication reliability of the legitimate wireless communication system. Even when the reliability requirement in terms of the OOPP is stringent, a mobile relay may be a must.
- When there is no mobile AP acting as a relay, the RNP can be adopted to ensure that the communication reliability of the legitimate communication system is still good as shown in Figs. 7 and 8.

VIII. CONCLUSION

In this paper, we investigate an uplink NOMA scenario with and without the support of a mobile AP acting as

a relay in the presence of jamming attacks. Particularly, we investigate two relaying protocols CNRP and FNRP, a retransmission scheme RNP using uplink pairwise NOMA, and an OMA scheme APAOP. First, we derive the IOP for each source node link to the relay and the destination as well as the relay link to the destination, and the IOPP for each source node considering multiple links in Nakagami-m fading channels during two phases. Subsequently, we define the OOPP as a suitable metric, the maximum value from the obtained IOPP for both source nodes. To address the uncertain environment associated with jamming attacks, we propose a method using RL that adapts to dynamic parameters, including information related to jammers and the positions of source nodes. The results indicate that both power allocation and the relay position play an important role to improve the communication reliability for the relaying protocols. Furthermore, the FNRP and RNP schemes offer superior reliability performance compared to the CNRP and APAOP. Finally, we provide a few guidelines for enhancing the communication reliability of the legitimate communication system. It can be seen that the derived closed-form expressions of the IOP, the IOPP, and the OOPP are useful to analyze and design the considered network also in the presence of jamming. Moreover, it can be concluded that outage analysis is an important tool of high practical relevance and it can be used by deep RL for online prediction of how to minimize outage probabilities.

APPENDIX

A. APPENDIX I

Lemma 1: Given that $h_s \sim G(m_s, \frac{\rho_s}{m_s})$ and $h_{\mathcal{J}_1} \sim G(m_{\mathcal{J}_1}, \frac{\rho_{\mathcal{J}_1}}{m_{\mathcal{J}_1}})$, the closed-form expressions of I_{1a} , I_{1b} , I_{3a} , I_{3b} , I_{40} , I_{41} , I_{42} , I_{20} , I_{21} , and I_{22} can be obtained as follows:

$$I_{1a} = \frac{(m_2 \rho_2^{-1})^{m_2} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-m_1 \rho_1^{-1} a_3}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_2)} \times \sum_{i=0}^{m_1-1} \frac{(m_1 \rho_1^{-1})^i}{i!} \sum_{j=0}^i \binom{i}{j} a_1^j \sum_{k=0}^{i-j} a_3^{i-j-k} a_2^k \times \binom{i-j}{k} \frac{\Gamma(m_2+j)}{B_1^{m_2+j}} \frac{\Gamma(m_{\mathcal{J}_1}+k)}{B_{12}^{m_{\mathcal{J}_1}+k}}, \quad (53)$$

$$I_{1b} = I_{8a} + I_{11a} - I_{10} + \frac{(m_2 \rho_2^{-1})^{m_2} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-B_3 a_7}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_2)} \times \sum_{j=0}^{m_1-1} \frac{(m_1 \rho_1^{-1})^j \Gamma(m_2+j)}{j! B_4^{m_2+j}} \sum_{l=0}^{m_2+j-1} \frac{B_4^l}{l!} \times \sum_{q=0}^l \binom{l}{q} a_6^q a_7^{l-q} \frac{\Gamma(m_{\mathcal{J}_1}+q)}{B_{13}^{m_{\mathcal{J}_1}+q}}, \quad (54)$$

$$I_{3a} = \frac{(m_1 \rho_1^{-1})^{m_1} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-m_2 \rho_2^{-1} b_3}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_1)}$$

$$\times \sum_{i=0}^{m_2-1} \frac{(m_2 \rho_2^{-1})^i}{i!} \sum_{j=0}^i \binom{i}{j} b_1^j \sum_{k=0}^{i-j} b_3^{i-j-k} b_2^k \times \binom{i-j}{k} \frac{\Gamma(m_1+j)}{B_7^{m_1+j}} \frac{\Gamma(m_{\mathcal{J}_1}+k)}{B_{14}^{m_{\mathcal{J}_1}+k}}, \quad (55)$$

$$I_{3b} = I_{9a} + I_{21a} - I_{20} + \frac{(m_1 \rho_1^{-1})^{m_1} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-B_4 b_7}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_1)} \times \sum_{j=0}^{m_2-1} \frac{(m_2 \rho_2^{-1})^j \Gamma(m_1+j)}{j! B_4^{m_1+j}} \sum_{l=0}^{m_1+j-1} \frac{B_4^l}{l!} \times \sum_{q=0}^l \binom{l}{q} b_6^q b_7^{l-q} \frac{\Gamma(m_{\mathcal{J}_1}+q)}{B_{15}^{m_{\mathcal{J}_1}+q}}, \quad (56)$$

$$I_{40} = \frac{(m_2 \rho_2^{-1})^{m_2} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-m_1 \rho_1^{-1} a_3 - B_1 a_5}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_2)} \times \sum_{i=0}^{m_1-1} \frac{(m_1 \rho_1^{-1})^i}{i!} \sum_{j=0}^i \binom{i}{j} a_1^j \Gamma(m_2+j) \sum_{k=0}^{i-j} a_3^{i-j-k} a_2^k \times \binom{i-j}{k} \sum_{l=0}^{m_2+j-1} \frac{B_1^l}{l!} \sum_{q=0}^l \binom{l}{q} a_5^{l-q} a_4^q \frac{\Gamma(m_{\mathcal{J}_1}+k+q)}{B_2^{m_{\mathcal{J}_1}+k+q}}, \quad (57)$$

$$I_{41} = \begin{cases} I_{41a} & a_5 < a_7, x_1 \leq 0 \\ I_{41b} & a_5 \geq a_7, x_1 > 0 \\ 0 & a_5 \geq a_7, x_1 \leq 0 \\ I_{41a} - I_{41b} & a_5 < a_7, x_1 > 0 \end{cases}, \quad (58)$$

in which I_{41a} , and I_{41b} are given as

$$I_{41a} = I_{40} - \frac{(m_2 \rho_2^{-1})^{m_2} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-m_1 \rho_1^{-1} a_3 - B_1 a_7}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_2)} \times \sum_{i=0}^{m_1-1} \frac{(m_1 \rho_1^{-1})^i}{i!} \sum_{j=0}^i \binom{i}{j} a_1^j \Gamma(m_2+j) \sum_{k=0}^{i-j} a_3^{i-j-k} a_2^k \times \binom{i-j}{k} \sum_{l=0}^{m_2+j-1} \frac{B_1^l}{l!} \sum_{q=0}^l \binom{l}{q} a_7^{l-q} a_6^q \frac{\Gamma(m_{\mathcal{J}_1}+k+q)}{B_3^{m_{\mathcal{J}_1}+k+q}}, \quad (59)$$

$$I_{41b} = \frac{(m_2 \rho_2^{-1})^{m_2} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-m_1 \rho_1^{-1} a_3 - B_1 a_5}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_2)} \times \sum_{i=0}^{m_1-1} \frac{(m_1 \rho_1^{-1})^i}{i!} \sum_{j=0}^i \binom{i}{j} a_1^j \Gamma(m_2+j) \sum_{k=0}^{i-j} a_3^{i-j-k} a_2^k \times \binom{i-j}{k} \sum_{l=0}^{m_2+j-1} \frac{B_1^l}{l!} \sum_{q=0}^l \binom{l}{q} a_5^{l-q} a_4^q \frac{\Gamma(m_{\mathcal{J}_1}+k+q, B_2 x_1)}{B_2^{m_{\mathcal{J}_1}+k+q}}$$

$$\frac{(m_2 \rho_2^{-1})^{m_2} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-m_1 \rho_1^{-1} a_3 - B_1 a_7}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_2)} \times \sum_{i=0}^{m_1-1} \frac{(m_1 \rho_1^{-1})^i}{i!} \sum_{j=0}^i \binom{i}{j} \frac{a_1^j \Gamma(m_2 + j)}{B_1^{m_2+j}} \sum_{k=0}^{i-j} a_3^{i-j-k} a_2^k \times \binom{i-j}{k} \sum_{l=0}^{m_2+j-1} \frac{B_1^l}{l!} \sum_{q=0}^l \binom{l}{q} a_7^{l-q} a_6^q \frac{\Gamma(m_{\mathcal{J}_1} + k + q, B_3 x_1)}{B_3^{m_{\mathcal{J}_1} + k + q}}, \quad (60)$$

$$I_{42} = \begin{cases} I_{42a} & a_5 \geq a_7, x_1 \leq 0 \\ I_{42b} & a_5 < a_7, x_1 \leq 0 \\ I_{42a} - Q_1 & a_5 \geq a_7, x_1 > 0 \\ I_{42b} + Q_1 & a_5 < a_7, x_1 > 0 \end{cases}, \quad (61)$$

in which I_{12a} , I_{12b} and Q_1 are given as

$$I_{42a} = \frac{(m_2 \rho_2^{-1})^{m_2} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-B_4 a_5}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_2)} \times \sum_{j=0}^{m_1-1} \frac{(m_1 \rho_1^{-1})^j \Gamma(m_2 + j)}{j! B_4^{m_2+j}} \sum_{l=0}^{m_2+j-1} \frac{B_4^l}{l!} \times \sum_{q=0}^l \binom{l}{q} a_4^q a_5^{l-q} \frac{\Gamma(m_{\mathcal{J}_1} + q)}{B_5^{m_{\mathcal{J}_1} + q}}, \quad (62)$$

$$I_{42b} = \frac{(m_2 \rho_2^{-1})^{m_2} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-B_4 a_7}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_2)} \times \sum_{j=0}^{m_1-1} \frac{(m_1 \rho_1^{-1})^j \Gamma(m_2 + j)}{j! B_4^{m_2+j}} \sum_{l=0}^{m_2+j-1} \frac{B_4^l}{l!} \times \sum_{q=0}^l \binom{l}{q} a_6^q a_7^{l-q} \frac{\Gamma(m_{\mathcal{J}_1} + q)}{B_6^{m_{\mathcal{J}_1} + q}}, \quad (63)$$

$$Q_1 = \frac{(m_2 \rho_2^{-1})^{m_2} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-B_4 a_5}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_2)} \times \sum_{j=0}^{m_1-1} \frac{(m_1 \rho_1^{-1})^j \Gamma(m_2 + j)}{j! B_4^{m_2+j}} \sum_{l=0}^{m_2+j-1} \frac{B_4^l}{l!} \times \sum_{q=0}^l \binom{l}{q} a_4^q a_5^{l-q} \frac{\Gamma(m_{\mathcal{J}_1} + q, B_5 x_1)}{B_5^{m_{\mathcal{J}_1} + q}} - \frac{(m_2 \rho_2^{-1})^{m_2} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-B_4 a_7}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_2)} \times \sum_{j=0}^{m_1-1} \frac{(m_1 \rho_1^{-1})^j \Gamma(m_2 + j)}{j! B_4^{m_2+j}} \sum_{l=0}^{m_2+j-1} \frac{B_4^l}{l!} \times \sum_{q=0}^l \binom{l}{q} a_6^q a_7^{l-q} \frac{\Gamma(m_{\mathcal{J}_1} + q, B_6 x_1)}{B_6^{m_{\mathcal{J}_1} + q}}, \quad (64)$$

$$I_{20} = \frac{(m_1 \rho_1^{-1})^{m_1} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-m_2 \rho_2^{-1} b_3 - B_7 b_5}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_1)} \times \sum_{i=0}^{m_2-1} \frac{(m_2 \rho_2^{-1})^i}{i!} \sum_{j=0}^i \binom{i}{j} \frac{b_1^j \Gamma(m_1 + j)}{B_7^{m_1+j}} \sum_{k=0}^{i-j} b_3^{i-j-k} b_2^k \times \binom{i-j}{k} \sum_{l=0}^{m_1+j-1} \frac{B_7^l}{l!} \sum_{q=0}^l \binom{l}{q} b_5^{l-q} b_4^q \frac{\Gamma(m_{\mathcal{J}_1} + k + q)}{B_8^{m_{\mathcal{J}_1} + k + q}}, \quad (65)$$

$$I_{21} = \begin{cases} I_{21a} & b_5 < b_7, x_2 \leq 0 \\ I_{21b} & b_5 \geq b_7, x_2 > 0 \\ 0 & b_5 \geq b_7, x_2 \leq 0 \\ I_{21a} - I_{21b} & b_5 < b_7, x_2 > 0 \end{cases}, \quad (66)$$

in which I_{21a} and I_{21b} are given as

$$I_{21a} = I_{20} - \frac{(m_1 \rho_1^{-1})^{m_1} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-m_2 \rho_2^{-1} b_3 - B_7 b_7}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_1)} \times \sum_{i=0}^{m_2-1} \frac{(m_2 \rho_2^{-1})^i}{i!} \sum_{j=0}^i \binom{i}{j} \frac{b_1^j \Gamma(m_1 + j)}{B_7^{m_1+j}} \sum_{k=0}^{i-j} b_3^{i-j-k} b_2^k \times \binom{i-j}{k} \sum_{l=0}^{m_1+j-1} \frac{B_7^l}{l!} \sum_{q=0}^l \binom{l}{q} b_7^{l-q} b_6^q \frac{\Gamma(m_{\mathcal{J}_1} + k + q)}{B_9^{m_{\mathcal{J}_1} + k + q}}, \quad (67)$$

$$I_{21b} = \frac{(m_1 \rho_1^{-1})^{m_1} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-m_2 \rho_2^{-1} b_3 - B_7 b_5}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_1)} \times \sum_{i=0}^{m_2-1} \frac{(m_2 \rho_2^{-1})^i}{i!} \sum_{j=0}^i \binom{i}{j} \frac{b_1^j \Gamma(m_1 + j)}{B_7^{m_1+j}} \sum_{k=0}^{i-j} b_3^{i-j-k} b_2^k \times \binom{i-j}{k} \sum_{l=0}^{m_1+j-1} \frac{B_7^l}{l!} \sum_{q=0}^l \binom{l}{q} b_5^{l-q} b_4^q \frac{\Gamma(m_{\mathcal{J}_1} + k + q, B_8 x_2)}{B_8^{m_{\mathcal{J}_1} + k + q}} - \frac{(m_1 \rho_1^{-1})^{m_1} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-m_2 \rho_2^{-1} b_3 - B_7 b_7}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_1)} \times \sum_{i=0}^{m_2-1} \frac{(m_2 \rho_2^{-1})^i}{i!} \sum_{j=0}^i \binom{i}{j} \frac{b_1^j \Gamma(m_1 + j)}{B_7^{m_1+j}} \sum_{k=0}^{i-j} b_3^{i-j-k} b_2^k \times \binom{i-j}{k} \sum_{l=0}^{m_1+j-1} \frac{B_7^l}{l!} \sum_{q=0}^l \binom{l}{q} b_7^{l-q} b_6^q \frac{\Gamma(m_{\mathcal{J}_1} + k + q, B_9 x_2)}{B_9^{m_{\mathcal{J}_1} + k + q}}, \quad (68)$$

$$I_{22} = \begin{cases} I_{22a} & b_5 \geq b_7, x_2 \leq 0 \\ I_{22b} & b_5 < b_7, x_2 \leq 0 \\ I_{22a} - Q_2 & b_5 \geq b_7, x_2 > 0 \\ I_{22b} + Q_2 & b_5 < b_7, x_2 > 0, \end{cases} \quad (69)$$

in which I_{22a} , I_{22b} , and Q_2 are given as

$$\begin{aligned} I_{22a} &= \frac{(m_1 \rho_1^{-1})^{m_1} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-B_4 b_5}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_1)} \\ &\times \sum_{j=0}^{m_2-1} \frac{(m_2 \rho_2^{-1})^j \Gamma(m_1 + j)}{j! B_4^{m_1+j}} \sum_{l=0}^{m_1+j-1} \frac{B_4^l}{l!} \\ &\times \sum_{q=0}^l \binom{l}{q} b_4^q b_5^{l-q} \frac{\Gamma(m_{\mathcal{J}_1} + q)}{B_{10}^{m_{\mathcal{J}_1}+q}}, \quad (70) \\ I_{22b} &= \frac{(m_1 \rho_1^{-1})^{m_1} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-B_4 b_7}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_1)} \\ &\times \sum_{j=0}^{m_2-1} \frac{(m_2 \rho_2^{-1})^j \Gamma(m_1 + j)}{j! B_4^{m_1+j}} \sum_{l=0}^{m_1+j-1} \frac{B_4^l}{l!} \\ &\times \sum_{q=0}^l \binom{l}{q} b_6^q b_7^{l-q} \frac{\Gamma(m_{\mathcal{J}_1} + q)}{B_{11}^{m_{\mathcal{J}_1}+q}}, \quad (71) \\ Q_2 &= \frac{(m_1 \rho_1^{-1})^{m_1} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-B_4 b_5}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_1)} \\ &\times \sum_{j=0}^{m_2-1} \frac{(m_2 \rho_2^{-1})^j \Gamma(m_1 + j)}{j! B_4^{m_1+j}} \sum_{l=0}^{m_1+j-1} \frac{B_4^l}{l!} \\ &\times \sum_{q=0}^l \binom{l}{q} b_4^q b_5^{l-q} \frac{\Gamma(m_{\mathcal{J}_1} + q, B_{10} x_2)}{B_{10}^{m_{\mathcal{J}_1}+q}} \\ &- \frac{(m_1 \rho_1^{-1})^{m_1} (m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-B_4 b_7}}{\Gamma(m_{\mathcal{J}_1}) \Gamma(m_1)} \\ &\times \sum_{j=0}^{m_2-1} \frac{(m_2 \rho_2^{-1})^j \Gamma(m_1 + j)}{j! B_4^{m_1+j}} \sum_{l=0}^{m_1+j-1} \frac{B_4^l}{l!} \\ &\times \sum_{q=0}^l \binom{l}{q} b_6^q b_7^{l-q} \frac{\Gamma(m_{\mathcal{J}_1} + q, B_{11} x_2)}{B_{11}^{m_{\mathcal{J}_1}+q}}, \quad (72) \end{aligned}$$

where $B_1 = m_2 \rho_2^{-1} + m_1 \rho_1^{-1} a_1$, $B_2 = m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1} + m_1 \rho_1^{-1} a_2 + B_1 a_4$, $B_3 = m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1} + m_1 \rho_1^{-1} a_2 + B_1 a_6$, $B_4 = m_2 \rho_2^{-1} + m_1 \rho_1^{-1}$, $B_5 = m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1} + B_4 a_4$, $B_6 = m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1} + B_4 a_6$, $B_7 = m_1 \rho_1^{-1} + m_2 \rho_2^{-1} b_1$, $B_8 = m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1} + m_2 \rho_2^{-1} b_2 + B_7 b_4$, $B_9 = m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1} + m_2 \rho_2^{-1} b_2 + B_7 b_6$, $B_{10} = m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1} + B_4 b_4$, $B_{11} = m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1} + B_4 b_6$, $B_{12} = m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1} + m_1 \rho_1^{-1} a_2$, $B_{13} = m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1} + B_4 a_6$, $B_{14} = m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1} + m_2 \rho_2^{-1} b_2$, $B_{15} = m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1} + B_4 b_6$, $x_1 = \frac{a_5 - a_7}{a_6 - a_4}$, and $x_2 = \frac{b_5 - b_7}{b_6 - b_4}$. $\binom{n}{k} = \frac{n!}{k!(n-k)!}$ is the binomial coefficient. $\Gamma(m, \mu) = \int_{\mu}^{\infty} t^{m-1} e^{-t} dt$ and $\Gamma(m) =$

$\int_0^{\infty} t^{m-1} e^{-t} dt$ are the upper incomplete Gamma function and Gamma function, respectively.

Proof: Applying [15, Lemma 1, Lemma 2, and Lemma 3], this Lemma is obtained. ■

B. APPENDIX II

Lemma 2: Given two random variables $h_{\mathcal{J}_2} \sim G(m_{\mathcal{J}_2}, \frac{\rho_{\mathcal{J}_2}}{m_{\mathcal{J}_2}})$ and $h^{\mathcal{RD}} \sim G(m^{\mathcal{RD}}, \frac{\rho^{\mathcal{RD}}}{m^{\mathcal{RD}}})$ where $m_{\mathcal{J}_2}$ and $m^{\mathcal{RD}}$ are positive integers, the closed-form expression of the probability $p^{\mathcal{RD}}$ can be derived as follows:

$$\begin{aligned} p^{\mathcal{RD}} &= 1 - \frac{(m_{\mathcal{J}_2} \rho_{\mathcal{J}_2}^{-1})^{m_{\mathcal{J}_2}} e^{-m^{\mathcal{RD}} \rho^{\mathcal{RD}} c}}{\Gamma(m_{\mathcal{J}_2})} \\ &\times \sum_{k=0}^{m^{\mathcal{RD}}-1} \frac{(m^{\mathcal{RD}} \rho^{\mathcal{RD}} c)^k}{k!} \\ &\times \sum_{j=0}^k \left\{ \binom{k}{j} \left(\frac{b}{c}\right)^j \frac{\Gamma(m_{\mathcal{J}_2} + j)}{B_{16}^{m_{\mathcal{J}_2}+j}} \right\}, \quad (73) \end{aligned}$$

where $B_{16} = m_{\mathcal{J}_2} \rho_{\mathcal{J}_2}^{-1} + m^{\mathcal{RD}} \rho^{\mathcal{RD}} c$.

Proof: Applying [10, Th. 1], this Lemma is proven. ■

C. APPENDIX III

Lemma 3: Given two random variables $h_{\mathcal{J}_1} \sim G(m_{\mathcal{J}_1}, \frac{\rho_{\mathcal{J}_1}}{m_{\mathcal{J}_1}})$, $h_{\mathcal{J}_2} \sim G(m_{\mathcal{J}_2}, \frac{\rho_{\mathcal{J}_2}}{m_{\mathcal{J}_2}})$ and $h_s^{SD} \sim G(m_s^{SD}, \frac{\rho_s^{SD}}{m_s^{SD}})$ where $m_{\mathcal{J}_1}$, $m_{\mathcal{J}_2}$ and m_s^{SD} are positive integers, the closed-form expressions of the probability p_1^{APAOP} and p_2^{APAOP} can be derived as follows:

$$\begin{aligned} p_1^{APAOP} &= 1 - \frac{(m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1})^{m_{\mathcal{J}_1}} e^{-m_1^{SD} \rho_1^{SD} c_{01}}}{\Gamma(m_{\mathcal{J}_1})} \\ &\times \sum_{k=0}^{m_1^{SD}-1} \frac{(m_1^{SD} \rho_1^{SD} c_{01})^k}{k!} \\ &\times \sum_{j=0}^k \left\{ \binom{k}{j} \left(\frac{b_{01}}{c_{01}}\right)^j \frac{\Gamma(m_{\mathcal{J}_1} + j)}{B_{17}^{m_{\mathcal{J}_1}+j}} \right\}, \quad (74) \end{aligned}$$

$$\begin{aligned} p_2^{APAOP} &= 1 - \frac{(m_{\mathcal{J}_2} \rho_{\mathcal{J}_2}^{-1})^{m_{\mathcal{J}_2}} e^{-m_2^{SD} \rho_2^{SD} c_{02}}}{\Gamma(m_{\mathcal{J}_2})} \\ &\times \sum_{k=0}^{m_2^{SD}-1} \frac{(m_2^{SD} \rho_2^{SD} c_{02})^k}{k!} \\ &\times \sum_{j=0}^k \left\{ \binom{k}{j} \left(\frac{b_{02}}{c_{02}}\right)^j \frac{\Gamma(m_{\mathcal{J}_2} + j)}{B_{18}^{m_{\mathcal{J}_2}+j}} \right\}, \quad (75) \end{aligned}$$

where $B_{17} = m_{\mathcal{J}_1} \rho_{\mathcal{J}_1}^{-1} + m_1^{SD} \rho_1^{SD} c_{01}$ and $B_{18} = m_{\mathcal{J}_2} \rho_{\mathcal{J}_2}^{-1} + m_2^{SD} \rho_2^{SD} c_{02}$.

Proof: Applying [10, Th. 1], this Lemma is proven. ■

REFERENCES

- [1] I. Budhiraja et al., "A systematic review on NOMA variants for 5G and beyond," *IEEE Access*, vol. 9, pp. 85573–85644, 2021.
- [2] Y. Liu, W. Yi, Z. Ding, X. Liu, O. A. Dobre, and N. Al-Dhahir, "Developing NOMA to next generation multiple access: Future vision and research opportunities," *IEEE Wireless Commun.*, vol. 29, no. 6, pp. 120–127, Dec. 2022.
- [3] Y. Liu et al., "Evolution of NOMA toward next generation multiple access (NGMA) for 6G," *IEEE J. Sel. Areas Commun.*, vol. 40, no. 4, pp. 1037–1071, Apr. 2022.
- [4] J. Montalban, E. Iradier, P. Angueira, O. Seijo, and I. Val, "NOMA-based 802.11n for industrial automation," *IEEE Access*, vol. 8, pp. 168546–168557, 2020.
- [5] Z. Ding, Z. Yang, P. Fan, and H. V. Poor, "On the performance of non-orthogonal multiple access in 5G systems with randomly deployed users," *IEEE Signal Process. Lett.*, vol. 21, no. 12, pp. 1501–1505, Dec. 2014.
- [6] S. Timotheou and I. Krikidis, "Fairness for non-orthogonal multiple access in 5G systems," *IEEE Signal Process. Lett.*, vol. 22, no. 10, pp. 1647–1651, Oct. 2015.
- [7] L. Dai, B. Wang, Y. Yuan, S. Han, I. Chih-Lin, and Z. Wang, "Non-orthogonal multiple access for 5G: Solutions, challenges, opportunities, and future research trends," *IEEE Commun. Mag.*, vol. 53, no. 9, pp. 74–81, Sep. 2015.
- [8] E. Iradier et al., "Analysis of NOMA-based retransmission schemes for factory automation applications," *IEEE Access*, vol. 9, pp. 29541–29554, 2021.
- [9] Y. Wang, M. Zheng, and H. Yu, "NODR: An NOMA-based retransmission scheme for URLLC in industrial wireless networks," *IEEE Sensors J.*, vol. 22, no. 20, pp. 20073–20084, Oct. 2022.
- [10] V.-L. Dao, L.-N. Hoang, S. Girs, and E. Uhlemann, "Outage performance of pairwise NOMA allowing a dynamic decoding order and optimal pairs of power levels," *IEEE Open J. Commun. Soc.*, vol. 1, pp. 1886–1906, 2020.
- [11] "Cisco Annual Internet Report (2018–2023)," Cisco Systems, Inc., San Jose, CA, USA, White Paper, Mar. 2020. [Online]. Available: <https://www.cisco.com/c/en/us/solutions/collateral/executive-perspectives/annual-internet-report/white-paper-c11-741490.html>
- [12] S. Yan, X. Cao, Z. Liu, and X. Liu, "Interference management in 6G space and terrestrial integrated networks: Challenges and approaches," *Intell. Conver. Netw.*, vol. 1, no. 3, pp. 271–280, 2020.
- [13] H. Pirayesh and H. Zeng, "Jamming attacks and anti-jamming strategies in wireless networks: A comprehensive survey," *IEEE Commun. Surveys Tuts.*, vol. 24, no. 2, pp. 767–809, 2nd Quart., 2022.
- [14] P. Angueira et al., "A survey of physical layer techniques for secure wireless communications in industry," *IEEE Commun. Surveys Tuts.*, vol. 24, no. 2, pp. 810–838, 2nd Quart., 2022.
- [15] V.-L. Dao, L.-N. Hoang, S. Girs, and E. Uhlemann, "Defeating jamming using outage performance aware joint power allocation and access point placement in uplink pairwise NOMA," *IEEE Open J. Commun. Soc.*, vol. 2, pp. 1957–1979, 2021.
- [16] (Nat. Inst. Stand. Technol., Gaithersburg, MD, USA). *Reactive Jamming Attack*. Accessed: Aug. 21, 2023. [Online]. Available: https://nvd.nist.gov/vuln/search/results?form_type=Basic&results_type=overview&query=wireless+jamming+attack&search_type=all
- [17] S. Diggavi, N. Al-Dhahir, A. Stamoulis, and A. Calderbank, "Great expectations: The value of spatial diversity in wireless networks," *Proc. IEEE*, vol. 92, no. 2, pp. 219–270, Feb. 2004.
- [18] Z. Ding, H. Dai, and H. V. Poor, "Relay selection for cooperative NOMA," *IEEE Wireless Commun. Lett.*, vol. 5, no. 4, pp. 416–419, Aug. 2016.
- [19] J. S. Yeom, Y.-B. Kim, and B. C. Jung, "Spectrally efficient uplink cooperative NOMA with joint decoding for relay-assisted IoT networks," *IEEE Internet Things J.*, vol. 10, no. 1, pp. 210–223, Jan. 2023.
- [20] B. Hu, L. Wang, S. Chen, J. Cui, and L. Chen, "An uplink throughput optimization scheme for UAV-enabled urban emergency communications," *IEEE Internet Things J.*, vol. 9, no. 6, pp. 4291–4302, Mar. 2022.
- [21] T.-H. Vu, T.-T. Nguyen, Q.-V. Pham, D. B. Da Costa, and S. Kim, "A novel partial decode-and-amplify NOMA-inspired relaying protocol for uplink short-packet communications," *IEEE Wireless Commun. Lett.*, vol. 12, no. 7, pp. 1244–1248, Jul. 2023.
- [22] S. Ghosh, A. Al-Dweik, and M.-S. Alouini, "On the performance of end-to-end cooperative NOMA-based IoT networks with wireless energy harvesting," *IEEE Internet Things J.*, vol. 10, no. 18, pp. 16253–16270, Sep. 2023.
- [23] H. Liu, N. I. Miridakis, T. A. Tsiftsis, K. J. Kim, and K. S. Kwak, "Coordinated uplink transmission for cooperative NOMA systems," in *Proc. IEEE Glob. Commun. Conf. (GLOBECOM)*, Abu Dhabi, UAE, 2018, pp. 1–6.
- [24] Y. Gao, B. Xia, Y. Liu, Y. Yao, K. Xiao, and G. Lu, "Analysis of the dynamic ordered decoding for uplink NOMA systems with imperfect CSI," *IEEE Trans. Veh. Technol.*, vol. 67, no. 7, pp. 6647–6651, Jul. 2018.
- [25] M. W. Baidas, "A game-theoretic approach to distributed energy-efficiency maximization in uplink NOMA relay Ad-Hoc networks," in *Proc. Int. Symp. Wireless Pers. Multimedia Commun. (WPMC)*, 2020, pp. 1–6.
- [26] V.-T. Truong, M.-T. Vo, Y. Lee, and D.-B. Ha, "Amplify-and-forward relay transmission in uplink non-orthogonal multiple access networks," in *Proc. NAFOSTED Conf. Inf. Comput. Sci. (NICS)*, Hanoi, Vietnam, 2019, pp. 1–6.
- [27] X. Xie, Y. Bi, and X. Nie, "Performance analysis of uplink cooperative NOMA system with an AF relay," in *Proc. IEEE Int. Conf. Commun. Technol. (ICCT)*, Nanning, China, 2020, pp. 178–181.
- [28] V.-T. Truong, D.-B. Ha, Y. Lee, and A.-N. Nguyen, "On performance of cooperative transmission in uplink non-orthogonal multiple access wireless sensor networks," in *Proc. Int. Conf. Recent Adv. Signal Process., Telecommun. Comput. (SigTelCom)*, Hanoi, Vietnam, 2020, pp. 56–60.
- [29] S. Lv, X. Xu, S. Han, and P. Zhang, "Buffer-aided relaying in NOMA-based MTC networks with finite blocklength and statistical QoS constraints," *IEEE Trans. Wireless Commun.*, vol. 21, no. 7, pp. 5406–5419, Jul. 2022.
- [30] P. Xu, J. Quan, Z. Yang, G. Chen, and Z. Ding, "Performance analysis of buffer-aided hybrid NOMA/OMA in cooperative uplink system," *IEEE Access*, vol. 7, pp. 168759–168773, 2019.
- [31] J. Quan, P. Xu, Y. Wang, and Z. Yang, "Hybrid NOMA/OMA with buffer-aided relaying for cooperative uplink system," in *Proc. 15th EAI Int. Conf. Heterog. Netw. Qual., Rel., Secur. Robust.*, 2019, pp. 169–183.
- [32] N. Nomikos, E. T. Michailidis, P. Trakadas, D. Vouyioukas, T. Zahariadis, and I. Krikidis, "flex-NOMA: Exploiting buffer-aided relay selection for massive connectivity in the 5G uplink," *IEEE Access*, vol. 7, pp. 88743–88755, 2019.
- [33] L. Feng, J. Chai, F. Zhou, and W. Li, "Energy-efficient joint optimization of channel assignment, power allocation, and relay selection based on hypergraph for uplink mMTC networks," *IEEE Trans. Green Commun. Netw.*, vol. 5, no. 1, pp. 203–215, Mar. 2021.
- [34] X. Xie, J. Liu, J. Huang, and S. Zhao, "Ergodic capacity and outage performance analysis of uplink full-duplex cooperative NOMA system," *IEEE Access*, vol. 8, pp. 164786–164794, 2020.
- [35] X. Yu, J. Cai, M. Xie, and T. Teng, "Ergodic rate analysis of massive spatial modulation MIMO system with NOMA in cooperative relay networks," *IEEE Syst. J.*, vol. 16, no. 3, pp. 4513–4524, Sep. 2022.
- [36] D. Ran, Z. Yang, L. Feng, W. Li, Q. Ou, and M. Cheriet, "Optimizing global channel matching for multi-hop uplink NOMA-assisted cellular IoT with cooperative relay," in *Proc. Int. Wireless Commun. Mobile Comput. (IWCMC)*, Limassol, Cyprus, 2020, pp. 1806–1811.
- [37] X. Zhang, S. Kang, and X. Fu, "Pattern division multiple access featuring amplify-and-forward relaying in an uplink network," *IEEE Access*, vol. 8, pp. 85656–85663, 2020.
- [38] S. Lv, X. Xu, S. Han, X. Tao, and P. Zhang, "Energy-efficient secure short-packet transmission in NOMA-assisted mMTC networks with relaying," *IEEE Trans. Veh. Technol.*, vol. 71, no. 2, pp. 1699–1712, 2022.
- [39] V.-L. Nguyen, V.-T. Truong, D.-B. Ha, T.-L. Vo, and Y. Lee, "Performance analysis of relay selection on cooperative uplink NOMA network with wireless power transfer," in *Proc. 6th Int. Conf. Ind. Netw. Intell. Syst.*, 2020, pp. 32–44.
- [40] S. Abdel-Razek, "Ergodic capacity analysis of uplink cooperative NOMA network based on statistical channel state information," *Wireless Pers. Commun.*, vol. 125, no. 1, pp. 737–753, Feb. 2022.

- [41] L. Han, W.-P. Zhu, and M. Lin, "Uplink outage performance of NOMA-based hybrid satellite-terrestrial relay networks over generalized inhomogeneous fading channels," *IEEE Trans. Commun.*, vol. 70, no. 4, pp. 2417–2434, Apr. 2022.
- [42] X. Lu et al., "Reinforcement learning-based physical cross-layer security and privacy in 6G," *IEEE Commun. Surveys Tuts.*, vol. 25, no. 1, pp. 425–466, 1st Quart., 2023.
- [43] X. Xu, Y. Liu, X. Mu, Q. Chen, H. Jiang, and Z. Ding, "Artificial intelligence enabled NOMA toward next generation multiple access," *IEEE Wireless Commun.*, vol. 30, no. 1, pp. 86–94, Feb. 2023.
- [44] A. Goldsmith, *Wireless Communications*. Stanford, CA, USA: Cambridge Univ. Press, 2005.
- [45] D. G. Reina, M. Askalani, S. L. Toral, F. Barrero, E. Asimakopoulou, and N. Bessis, "A survey on multihop Ad Hoc networks for disaster response scenarios," *Int. J. Distrib. Sens. Netw.*, vol. 11, Jan. 2015, Art. no. 647037, doi: [10.1155/2015/647037](https://doi.org/10.1155/2015/647037).
- [46] V.-H. Dang et al., "Throughput optimization for NOMA energy harvesting cognitive radio with multi-UAV-assisted relaying under security constraints," *IEEE Trans. Cogn. Commun. Netw.*, vol. 9, no. 1, pp. 82–98, Feb. 2023.
- [47] G. Celik and H. Celebi, "User fairness for RSS-based positioning in uplink cooperative NOMA," *IET Commun.*, vol. 13, no. 20, pp. 3434–3442, Dec. 2019.
- [48] V. Dao, H. Tran, S. Girs, and E. Uhlemann, "Reliability and fairness for UAV communication based on non-orthogonal multiple access," in *Proc. IEEE Int. Conf. Commun. Workshops*, Shanghai, China, 2019, pp. 1–6.
- [49] R. S. Sutton and A. G. Barto, *Reinforcement Learning: An Introduction*. Cambridge, MA, USA: MIT Press, 2018.
- [50] H. Tabassum, M. S. Ali, E. Hossain, M. J. Hossain, and D. I. Kim, "Uplink vs. downlink NOMA in cellular networks: Challenges and research directions," in *Proc. IEEE Veh. Technol. Conf.*, Sydney, NSW, Australia, 2017, pp. 1–7.
- [51] X. Li, J. Li, Y. Liu, Z. Ding, and A. Nallanathan, "Residual transceiver hardware impairments on cooperative NOMA networks," *IEEE Trans. Wireless Commun.*, vol. 19, no. 1, pp. 680–695, Jan. 2020.



VAN-LAN DAO (Member, IEEE) received the B.S. degree in electrical and electronic engineering and the M.S. degree in automation and control engineering from Le Quy Don Technical University, Vietnam, in 2011 and 2016, respectively, and the Ph.D. degree in computer science and engineering from Malardalen University, Sweden, in 2023, where he is a Researcher. He was a Visiting Researcher with the University of Electro-Communications, Japan, for two months in 2017. He also worked with industrial partner, Robot

Technology—Danish Technological Institute, Denmark, for five months in 2021. His research interests include embedded system design for Internet of Things, wireless sensor networks, and wireless body area networks, hardware security, digital VLSI/ASIC design and FPGA-based system hardware design, non-orthogonal multiple access, timely and reliable communication for wireless networks, physical layer security, metaheuristic optimization, game theory, and machine learning.



ELISABETH UHLEMANN (Senior Member, IEEE) received the Ph.D. degree in communications theory from Chalmers University of Technology, Sweden, in 2004. He is a Full Professor in Data Communications with Malardalen University, Sweden. She has held visiting positions with the University of South Australia in 2005, the Technical University of Berlin in 2007, and the University of Canterbury, New Zealand, in 2011. She has also worked as a Consultant with Volvo Technology from 2005 to 2009 dealing with connected vehicles; as a Consultant with Ikanos Communications, CA, USA, in 2005 with VDSL protocols and with Free2move, Sweden, from 2009 to 2010 with wireless audio. She serves as a Senior Editor for *IEEE Vehicular Technology Magazine* on Connected and Automated Vehicles and was the Former Vice Chair of the Swedish VT/COM/IT chapter. She has served as a Research Grant Reviewer in three different countries and as a member of different Ph.D. examination committees over 25 times in five different countries.



SVETLANA GIRS (Member, IEEE) received the B.Sc. and M.Sc. degrees in telecommunications from Saint-Petersburg State Polytechnic University, Russia, in 2009 and 2011, respectively, and the Ph.D. degree in computer science and engineering from Malardalen University in February 2016. She is a Senior Lecturer with the School of Innovation, Design and Engineering, Malardalen University, Västerås, Sweden, where she is currently leading the Networked and Embedded Systems Division. She has been a

Visiting Researcher with the University of Canterbury, Christchurch, New Zealand, for three months in 2014. Her research interests include real-time communication, cooperative relay networks, and reliable wireless communication for industrial systems. She is also the Co-Chair of the Subcommittee on Industrial Communication Systems within the IEEE IES Technical Committee on Factory Automation.