

## RESEARCH ARTICLE

# Efficient Region-Based Skyline Computation for a Group of Users

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**ABSTRACT** Recently, with the advancement of technology, ad-hoc meetings or impromptu gatherings are becoming more and more common. The meetings/gatherings which involve at least two people will require a specific physical point location that is useful or interesting to them, called *point of interest (PoI)*. These people might be residing at different locations; each with their own preferences which most likely to be different. Undoubtedly, given  $n$  people in a group, there will be  $n$  users' preferences. Finding a suitable *PoI* that meets these  $n$  users' preferences is not a straightforward task. Existing solutions that utilise skyline processing in discovering the best, most preferred objects in satisfying the preferences of a group of users within a predetermined area have shown acceptable results. However, these solutions have to be executed repeatedly for each query of a group of users since they do not exploit the possibilities that an area that has been visited by a group of users might be the area of interest of another group of users in the future. Inherently, they require rescanning the objects and recomputing the skylines of a previously visited region which is undoubtedly unwise and costly. This paper proposes the *Region-based Skyline for a Group of Users (RSGU)* and *Extended Region-based Skyline for a Group of Users (ERSGU)* frameworks which attempt to resolve the limitations of existing solutions. In this work, skyline objects are *PoIs* that are recommended to a group of users that are derived by analysing both the locations of the users, i.e. spatial attributes, as well as the spatial and non-spatial attributes of objects that are within a predetermined region of the group of users. Here, each region is partitioned into smaller units called fragments in such a way that overlapping areas between the currently and previously visited regions can be easily determined; while the results of computing the skylines of each fragment, known as fragment skylines, are saved to be utilised by the subsequent requests. Meanwhile, *ERSGU* has an additional feature in which the skylines derived for a group of users are not only based on the evaluation of the spatial and non-spatial attributes of the objects, but also the closeness of the objects to the desirable facilities or other interesting objects in the region. Undeniably, a *PoI* that is nearby to other attractions is appealing and worth the journey. Several experiments have been conducted and the results show that our proposed frameworks outperform the previous work with respect to CPU time.

**INDEX TERMS** Multi-criteria decision making, skyline queries, group of users, spatial and non-spatial attributes.

## I. INTRODUCTION

Query processing is defined as the process of answering a query (request) to a database or an information system, which usually involves the following three main activities:

(i) analysing and interpreting the query, (ii) searching through the space of stored data, and (iii) retrieving the results satisfying the query. The traditional query processing operates either by retrieving objects<sup>1</sup> from a collection of objects

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<sup>1</sup>Without loss of generality, the term object is used throughout this paper to be in line with other research works in similar area. The terms data, data item, record, and tuple can also be used in this context.

that strictly satisfy each condition specified in the query or returning an empty result if otherwise. The recent developments in query processing attempt to relax these stringent requirements, by retrieving the best, most preferred objects from the collection according to the conditions specified in the query, also known as *user-defined preferences*. These preference queries employ preference evaluation techniques, have achieved significant success, as they are widely used in applications related to multi-criteria decision support. During the past two decades, several preference evaluation techniques have been introduced, among them are: top- $k$  [1], [2], [3], [4], [5], [6], [7], [8], [9], [10], [11], skyline [12], [13], [14], [15], [16], [17], [18], [19], [20], [21], [22], [23], [24], [25], [26], [27], [28], [29], [30], [31], [32], [33], [34], [35],  $k$ -dominance [36], [26], [37], [38], [39], [40], [41] top- $k$  dominating [3], [42], [5], [43], [44], [45], [46], [47], [48], and  $k$ -frequency [49].

The skyline operator introduced by [12] which is used to filter a set of interesting objects from a potentially large multi-dimensional set of objects by keeping only those objects that are not worse than any other, has been greatly explored in several studies in an attempt to accurately and efficiently solve problems of real-world applications that are related to decision support and decision making. It attempts to derive the best, most preferred set of objects known as *skylines objects* (or *skylines* in short) according to a set of evaluation criteria.<sup>2</sup> The process of computing skylines becomes more challenging when conflicting criteria are involved while the number of criteria to be considered is huge. A classic example is selecting a hotel for a holiday whereby hotels that are close to the beach are known to be expensive. While other criteria like facilities, rating, and service, are equally important, distance and price are examples of conflicting criteria.

Since the introduction of the skyline operator by [12], an abundance of skyline algorithms have been proposed for data processing in order to retrieve useful insights. These variants of skyline algorithms are introduced to deal with different characteristics of data, such as uncertain data [58], incomplete data [50], [51], [52], [53], encrypted data [54], and streaming data [55]; while others are based on the platform being considered like distributed database [51], cloud computing [38], road networks [56], and others [57], [58], [59]. Nonetheless, these skyline algorithms focus mainly on the optimisation problem of skyline computation for a given single user query. Since the location of the user is insignificant in these studies, hence, it is sufficient to derive skyline objects by considering only the attributes of these objects, also known as *non-spatial* attributes, as the evaluation criteria.

However, in real world scenario, due to the advancement of technology, ad-hoc meetings or impromptu gatherings are

<sup>2</sup>In this paper, the evaluation criteria used in determining the skyline objects are the attributes (the terms dimension and attribute are used interchangeably) of the objects.

becoming more and more common [60], [61], [62], [63], [64], [65], [66], [67]. Intuitively, the meetings/gatherings will involve a group of people (at least two people) and they will have to decide on a specific physical point location that is useful or interesting to them, called *point of interest*<sup>3</sup> (henceforth referred to as *PoI*). Some examples of *PoI* are restaurants, hotels, cafes, etc. These people might be residing at different locations; each with their own preferences which are most likely to be different. Undoubtedly, given  $n$  people in a group, there will be  $n$  users' queries.<sup>4</sup> The existing skyline algorithms are unsuited for such a scenario for two main reasons: (i) they cater a user's preferences at a time (single user query) and (ii) they deal only with the non-spatial attributes of the objects. While, deriving the skyline objects of a group of users that are located at different locations, both the spatial and non-spatial attributes of the objects need to be considered. Thus, identifying objects that best meet the preferences of a group of users is crucial and challenging.

The following scenarios typify the samples of situations considered in this paper.

*Scenario I:* Assuming several users whom are not close to each other would like to meet; hence a group is said to be formed and these users will have to decide on a specific *PoI* within a predetermined area<sup>5</sup> that is useful or interesting to them. There are many *PoIs* that they can choose. However, several criteria need to be considered before they decide on a specific *PoI* to visit. These include the location of the *PoI*, i.e. how far it is from the location of each user (spatial attribute), the opening hour, food, ticket price, rating, facilities provided, etc (non-spatial attributes). A *PoI* which is near to the users might not be the *PoI* that meets all the users' preferences. While a *PoI* which provides facilities that meet most of the users' preferences might be located far away from these users. Therefore, recommending a *PoI* to visit to these users is not a straightforward task as many criteria need to be considered. These include the location of each user (spatial attribute), the locations of the objects (spatial attribute), and the features of the objects (non-spatial attributes). It becomes more complicated when the number of users in the group is large, the objects to be analysed in the space (area) are high multi-dimensional while their number is huge. Meanwhile, an area that has been visited by a group of users might be the area of interest of another group of users in the future. Thus, it is essential to have a method that could find an object(s) that dominates other objects that best suits the preferences of a group of users with respect to both the spatial (location) and non-spatial attributes of the objects that are within a predetermined area. Also, it is hypothesised that utilising

<sup>3</sup>The term object of interest is also used that reflects the *PoI* that is saved as object in a data set.

<sup>4</sup>In this paper, the  $n$  users' queries are assumed to be distinct (reflecting different preferences) due to the fact that the skyline objects for a user might be different from another user as the area covered by each user's query might be different.

<sup>5</sup>Without loss of generality, the terms area, region, and space are used interchangeably throughout this paper.

the previous skyline computation results in identifying the skyline objects of the subsequent group of users can greatly reduce the skyline computation time.

*Scenario II:* Similar to the *Scenario I* described above, the group of users might be looking for a *PoI* to visit which is near to other useful or interesting facilities/objects (other types of *PoI*). Undeniably, a *PoI* that is nearby to other attractions is appealing and worth the journey. Therefore, besides considering the spatial attribute, i.e. the location of the objects, and the non-spatial attributes of the objects like opening hour, food, ticket price, rating, etc; another criterion needs to be established which is the distance between the *PoI* to other useful or interesting facilities/objects like mosque, cinema, hospital, etc. Thus, it is essential to have a method that could find an object(s) that dominates other objects that best suits the preferences of a group of users with respect to both the spatial (location) and non-spatial attributes of the objects as well as how close the object is to other useful or interesting facilities/objects that are within a predetermined area.

This paper takes the challenge to solve the problem associated to identifying skyline objects for a group of users whom intend to have a meeting/gathering. Two different solutions are proposed, each handling a different scenario as described in *Scenario I* and *Scenario II*. Generally, each solution will make use of the spatial attribute of the users, as well as the spatial and non-spatial attributes of the objects. In general, the main contributions of this work are briefly described as follows:

- We have formally introduced the problem of computing skylines of a group of users and justify the significance of addressing the problem.
- We have proposed an efficient solution, named *Region-based Skyline for a Group of Users (RSGU)* framework that is designed for processing the skyline queries of a given group of users by considering both the locations of the users in the group, as well as the spatial and non-spatial attributes of the objects that are within a predetermined region of the group of users; with two main aims that are (i) avoiding the process of rescanning the set of objects within a predetermined region that is known to have been previously visited by a group of users and (ii) avoiding the recomputation of skylines of a set of objects within a predetermined region that has been analysed in earlier computations of previously visited group of users.
- We have proposed the *Extended Region-based Skyline for a Group of Users (ERSGU)* framework, an enhancement of the *RSGU* framework, which has similar aims as *RSGU* with an additional feature in which the closeness of the objects to the other desirable facilities or interesting objects in the region are taken into consideration in the derivation of skyline objects for a group of users.
- We have conducted extensive experiments to prove *RSGU*'s and *ERSGU*'s capabilities in deriving the skyline objects for a group of users.

The rest of the paper is structured as follows. In Section II, the previous works that are related to computing skylines for a single user as well as for a group of users are presented. In Section III, the necessary definitions and notations, which are used throughout the paper, are set out. Section IV and Section V elaborate our proposed frameworks, *RSGU* and *ERSGU* respectively, that are purposely designed for handling the computation of skyline objects for a group of users. A running database example is also given to clarify the phases of the proposed frameworks. The experimental results are demonstrated in Section VI. Conclusion and further research direction are depicted in the final section, Section VII.

## II. RELATED WORKS

Since the introduction of the skyline operator by [12] many variants of skyline algorithms have been proposed. Although the ultimate goal of these algorithms is to derive the best, most preferred objects from a multi-dimensional set of objects, each of them tackled a slightly different issue. We categorised these skyline algorithms into two main categories, namely: skyline algorithms for a single user and skyline algorithms for a group of users.

*Skyline algorithms for a single user* – Generally, these skyline algorithms attempt to optimise the process of filtering the best, most preferred objects from a potentially large multi-dimensional set of objects. These algorithms aim at reducing the processing time by reducing the search space as small as possible. Thus, ensuring that only the set of objects that may potentially be the skylines is analysed. In this category, users' queries are assumed to have the same objective function; hence users are assumed to have the same preferences.

Among the earlier and most cited skyline algorithms in the literature are *Block Nested Loop (BNL)* [12], *Divide-and-Conquer (D&C)* [12], *Linear Elimination Sort for Skyline (LESS)* [68], *Branch and Bound Skyline (BBS)* [69], *SkyCube* [12], and *Sort and Limit Skyline algorithm (SaLSa)* [70]. Recently, several skyline algorithms have been proposed that attempt not only to resolve the optimisation problem but also issues related to the uncertainty of data; which is defined as the degree to which data are inaccurate, imprecise, untrusted, unknown or incomplete. These include among others *ISkyline* [22], *sorting-based bucket skyline* [71], *Incoskyline* [72], *Jincoskyline* [73], and *OIS* [74] that handle the issues of incompleteness of data. The incompleteness of data leads to the loss of *transitivity property* of skyline technique. It also leads to *cyclic dominance* between the objects as some objects are incomparable to each other that results in no object can be considered as skyline. Meanwhile, *probabilistic skyline model* [75],  *$\tau$ -Skyline* [76], *SkyQUD* [77], [78], [79], [80] and *SQUiD* [71] focus on the challenges in computing skyline queries for uncertain database. Here, the exact values of the objects are not known at the point of processing. Consequently, one cannot derive the exact skyline but can only compute the probability of an object being a skyline member. On the other hand, the works by [81] and [82]

attempt to solve the issues related to uncertain data in a data stream. Processing such data is challenging due to the objects in the stream arrive online and data streams are potentially unbounded in size. Besides, the work by [52] focuses on dynamic database. Nonetheless, these skyline algorithms are specifically designed to cater only a single user query, i.e. only a single user’s preferences is considered in the skyline computation.

*Skyline algorithms for a group of users* – These skyline algorithms compute the skyline objects of a group of users from a potentially large multi-dimensional set of objects. As we assume that the objects are static, hence we further elaborate only those works that are similar to our intention. To the best of our knowledge the only works that contribute to skyline queries for a group of users are the works done by [60], [61], and [63]. In processing spatial skyline query for a group of users, two algorithms are proposed by [60], namely:  $B^2S^2$  and  $VS^2$ . Both algorithms assumed that the user points are static. The  $B^2S^2$  algorithm utilises the  $R$ -tree while the  $VS^2$  algorithm utilises the Voronoi diagram. Then, [61] proposed the  $VCS^2$  algorithm which enhanced the work by [60].  $VCS^2$  algorithm aims at processing skyline query by taking into consideration the movements of the users. However,  $VCS^2$  only calculates the last location of the users and does not consider the changes of locations to prevent recalculation of the skylines. In [63], the authors proposed the  $VR$  algorithm, that combined two data structures as used in [61],  $R$ -tree and Voronoi, in order to find spatial skylines for a group of user points. In their work, both the user points and objects are considered static. While the spatial and non-spatial attributes of the objects are analysed to find the skylines. Meanwhile, our previous solution,  $SGMU$  [83], is designed with the main aim to continuously derive skylines for a group of mobile users.

Although [60], [61], [63], [83] considered the spatial attributes of the group of users in determining the skylines, but there is no attempt made to avoid rescanning of objects of previously visited regions and simultaneously avoid repeating the process of pairwise comparisons among the objects.

### III. DEFINITIONS AND NOTATIONS

In this section, we present the necessary definitions and introduce the notations that are used throughout this paper. First, we give the definitions that are related to  $RSGU$ . This is then followed with definitions of  $RSGU$  that are modified/extended to suit with the  $ERSGU$ ’s solution. Examples are provided where necessary to further clarify the definitions. A formal definition of the problem addressed by each solution is then put forward at the end of each section.

#### A. PRELIMINARIES OF RSGU

To clarify the concepts and steps proposed in this work, the following sample of data is used. Table 1(a) and Table 1(b) present the spatial attribute (*Location*) of the users of group  $a$ ,  $G_a$ , and group  $b$ ,  $G_b$ , respectively. Here, the request submitted by  $G_a$  is assumed at time  $t_a$ , while the request submitted

TABLE 1. The spatial attribute of the users.

ID	Location	ID	Location
$u_1$	(8, 8)	$u_1$	(5, 8)
$u_2$	(14, 16)	$u_2$	(10, 10)
$u_3$	(2, 5)	$u_3$	(18, 10)

(a) Group  $a$ ,  $G_a$                       (b) Group  $b$ ,  $G_b$

TABLE 2. The spatial and non-spatial attributes of the objects.

Restaurant	Location	Rate	Price	Restaurant	Location	Rate	Price
$o_1$	(2, 3)	3	70	$o_{24}$	(4, 13.3)	3	75
$o_2$	(3, 4)	4	65	$o_{25}$	(7, 13)	1	90
$o_3$	(3, 1)	5	80	$o_{26}$	(16, 15)	2	86
$o_4$	(7, 1.7)	2	75	$o_{27}$	(20, 14)	5	80
$o_5$	(6, 5)	3	65	$o_{28}$	(23, 20)	3	60
$o_6$	(7, 7)	5	70	$o_{29}$	(21, 21)	5	62
$o_7$	(9, 8)	1	80	$o_{30}$	(17, 23)	4	95
$o_8$	(8, 9.7)	2	85	$o_{31}$	(14, 20)	2	65
$o_9$	(7, 11)	4	73	$o_{32}$	(13, 18)	2	55
$o_{10}$	(10, 5)	3	50	$o_{33}$	(10, 19)	3	70
$o_{11}$	(10.7, 6)	1	65	$o_{34}$	(1, 16)	4	62
$o_{12}$	(15, 2)	2	80	$o_{35}$	(3, 22)	4	81
$o_{13}$	(17, 1)	5	105	$o_{36}$	(7, 20)	3	90
$o_{14}$	(22, 4.7)	4	90	$o_{37}$	(24, 15)	2	66
$o_{15}$	(17, 5.7)	3	85	$o_{38}$	(-3, -1)	1	57
$o_{16}$	(20, 7)	4	90	$o_{39}$	(-1, 7)	1	61
$o_{17}$	(23, 9)	1	55	$o_{40}$	(10.3, 13)	4	71
$o_{18}$	(16, 8)	2	54	$o_{41}$	(-4, 4)	3	98
$o_{19}$	(14, 10)	4	80	$o_{42}$	(8, -2)	2	58
$o_{20}$	(11, 9.7)	5	56	$o_{43}$	(8, 18)	2	85
$o_{21}$	(4, 10)	3	67	$o_{44}$	(-2, 10)	4	70
$o_{22}$	(2, 12)	5	100	$o_{45}$	(3, -1)	5	80
$o_{23}$	(3, 13)	4	74				

by  $G_b$  is at time  $t_b$  where  $t_a < t_b$ . Table 2 presents the spatial (*Location*) and non-spatial (*Rate*, *Price*) attributes of the objects. For the non-spatial attributes, it is assumed that higher rate and lower price are preferable.

Given a data set  $D = \langle R, U, O \rangle$ , where  $U = \{u_1, u_2, \dots, u_n\}$  is a list of  $n$  users,  $O = \{o_1, o_2, \dots, o_m\}$  is a list of  $m$  objects, and  $R = \langle A_S, A_N \rangle$  where  $A_S$  representing a spatial attribute while  $A_N = \{d_1, d_2, \dots, d_l\}$  is a set of non-spatial attributes. The following definitions defined the properties of a user and an object as used in this work.

*Definition 1 (Property of a User)*: Each user,  $u_i \in U$ , is associated with a spatial attribute which represents the location of the user at a time,  $t$ . This is denoted as  $u_i(x_i, y_i)$ . For instance,  $u_1(8, 8)$  of Table 1(a) denotes the location of user  $u_1$  at time  $t_a$ .

*Definition 2 (Properties of an Object)*: Each object  $o_j \in O$  has two main elements denoted by  $o_j = (s_j, ns_j)$  where  $s_j$  is the value of spatial attribute (location),  $A_S$ , and  $ns_j = \{o_j.d_1, o_j.d_2, \dots, o_j.d_l\}$  is a set of values of non-spatial attributes,  $A_N$ , associated to  $o_j$ . The location of an object  $o_j \in O$  is denoted as  $o_j(x_j, y_j)$ . As each object  $o_j \in O$  is assumed to be static, thus the location of the object is fixed regardless the changes in time. Hence,  $o_j = (s_j, ns_j)$  can be written as  $o_j = ((x_j, y_j), \{o_j.d_1, o_j.d_2, \dots, o_j.d_l\})$ . For instance, the object  $o_1$  of Table 2 can be written as  $o_1 = ((2, 3), \{3, 70\})$ .

The following definitions defined the notion of dominance in this work.

**Definition 3 (Dominance<sup>6</sup>):** Given two objects  $o_i = (s_i, ns_i)$  and  $o_j = (s_j, ns_j) \in O$  where  $i \neq j$ ,  $o_i$  is said to dominate  $o_j$  (denoted by  $o_i < o_j$ ) if and only if both of the following conditions hold: (1)  $o_i$  non-spatially dominates  $o_j$  ( $o_i <_{ns} o_j$ ) and (2)  $o_i$  spatially dominates  $o_j$  ( $o_i <_s o_j$ ).

**Definition 4 (Non-spatial Dominance):** Given two objects  $o_i = (s_i, ns_i)$  and  $o_j = (s_j, ns_j) \in O$  where  $i \neq j$ ,  $o_i$  is said to non-spatially dominate  $o_j$  (denoted by  $o_i <_{ns} o_j$ ) if and only if  $o_i$  is no worse than (in this definition, greater value is preferable)  $o_j$  in all the non-spatial attributes,  $A_N$ . This is formally written as follows:  $o_i <_{ns} o_j$  if and only if  $\forall d_k \in A_N, o_i.d_k \geq o_j.d_k \wedge \exists d_l \in A_N, o_i.d_l > o_j.d_l$ . For instance, given  $o_6 = ((7, 7), \{5, 70\})$  and  $o_{12} = ((15, 2), \{2, 80\})$ ,  $o_6 <_{ns} o_{12}$  since  $o_6$  is better than  $o_{12}$  in both the dimensions *Rate* and *Price*; with the assumption that higher rate and lower price are preferable.

**Definition 5 (Spatial Dominance):** Given two objects  $o_i = (s_i, ns_i)$  and  $o_j = (s_j, ns_j) \in O$  where  $i \neq j$ ,  $o_i$  is said to spatially dominate  $o_j$  (denoted by  $o_i <_s o_j$ ) if and only if for every user  $u_k \in U$ , the distance between  $o_i$  and  $u_k$ ,  $dist(o_i, u_k)$ , is no worse than the distance between  $o_j$  and  $u_k$ ,  $dist(o_j, u_k)$ . This is formally written as follows:  $o_i <_s o_j$  if and only if  $\forall u_k \in U, dist(o_i, u_k) \leq dist(o_j, u_k) \wedge \exists u_l \in U, dist(o_i, u_l) < dist(o_j, u_l)$ . For instance, the distances between  $o_1 = ((2, 3), \{3, 70\})$  and  $u_1, u_2$ , and  $u_3$  of group  $G_a$  are 7.81, 17.69, and 2, respectively; while the distances between  $o_2 = ((3, 4), \{4, 65\})$  and  $u_1, u_2$ , and  $u_3$  of group  $G_a$  are 6.4, 16.27, and 1.41, respectively. Thus,  $o_2 <_s o_1$ .

**Definition 6 (Dominance in a Space):** Given a bounded space,  $S$  (region, MBR, fragment, area, polygon, etc.), and two objects  $o_i = (s_i, ns_i)$  and  $o_j = (s_j, ns_j) \in O$  where  $i \neq j$  in  $S$ ,  $o_i$  is said to dominate  $o_j$  (denoted by  $o_i < o_j$ ) in  $S$  if and only if (1)  $o_i$  non-spatially dominates  $o_j$  ( $o_i <_{ns} o_j$ ) in  $S$  and (2)  $o_i$  spatially dominates  $o_j$  ( $o_i <_s o_j$ ) in  $S$ .

**Definition 7 (Skylines of a Space):** An object  $o_i \in O$  in a space  $S$  is a skyline of  $S$  if there are no other objects  $o_j \in O$  where  $i \neq j$  in the space  $S$  that dominates  $o_i$ . In this paper,  $Sky_{G_p}$  is used to denote the skyline set for the group  $G_p$  of a given space  $S$ .

Based on the above definitions, the problem that is tackled by *RSGU* is formulated as follows:

Given a group of users,  $G_p = \{u_1, u_2, \dots, u_p\}$ , where  $G_p \subset U$ , and the candidate skylines of  $G_p$  in region  $R_p$  denoted as  $CS_{G_p}$ . Find the skylines of a group of users  $G_q = \{u_1, u_2, \dots, u_q\}$  in region  $R_q$ , i.e.  $CS_{G_q}$ , where  $G_q \subset U$ ,  $G_q \neq G_p$ , and  $R_q \cap R_p \neq \emptyset$  by utilising  $CS_{G_p}$  that has been derived for  $G_p$ . This is depicted in Fig. 1 where the area covered to compute the skylines for  $G_q$  that falls in the region  $S_{G_q}$  can be reduced to the area defined by  $S_{G_q} - S_{G_p}$ , while the results of skyline computation that have been performed earlier over the area  $S_{G_p} \cap S_{G_q}$  for  $G_p$  can be avoided by simply utilising the obtained results derived for  $G_p$ , i.e.  $CS_{G_p}$ .

<sup>6</sup>Without loss of generality, the definition is applicable for a given bounded space,  $S$ , i.e.  $O$  is a set of objects in the space  $S$ . Similar note applies for Definition 4 and Definition 5.

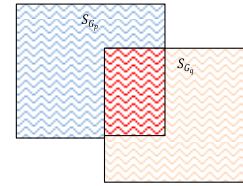


FIGURE 1. The reduction area in deriving skylines for a group of users.

TABLE 3. The spatial attribute of the interesting objects.

ID	Location
$IN_{1-1}$	(10, 7)
$IN_{1-2}$	(5, 12)
$IN_{1-3}$	(3, 15)
$IN_{2-1}$	(2.5, 9)
$IN_{2-2}$	(8, 4)
$IN_{2-3}$	(11.5, 3)
$IN_{3-1}$	(16, 10)
$IN_{3-2}$	(4, 11)
$IN_{3-3}$	(9, 6.5)

### B. PRELIMINARIES OF ERSQU

Since *ERSQU* is an extended framework of *RSGU*, thus the concepts, terms, and notations introduced and clarified in Part A above are applied here. Also, the sample of data given in Table 1 and Table 2 will be referred to as example to clarify the steps of *ERSQU*. Therefore, the definitions of *property of a user*, *properties of an object*, *non-spatial dominance*, *spatial dominance*, *non-spatial dominance of the fragment  $F_k$* , and *candidate skylines of the fragment  $F_k$*  are omitted here; readers can easily refer to their definitions in Part A. Nevertheless, three new definitions are included that are *Definition 10 Property of an Interesting Object*, *Definition 11 Closest Property of a  $t$  Type of Interesting Objects*, and *Definition 13  $p$ -Closest Dominance*. These are further elaborated below.

Given a set of  $p$  distinct types of interesting objects,  $IN = \{IN_1, IN_2, \dots, IN_p\}$ , where  $IN_t = \{IN_{t-1}, IN_{t-2}, \dots, IN_{t-n}\}$  is a list of  $n$  interesting objects of type  $t$ . The notation  $IN_{t-l}$  is used to denote the  $l$ -th interesting object of type  $t$ .

**Definition 10 (Property of an Interesting Object):** Each interesting object,  $IN_{t-l} \in IN_t$ , is associated with a spatial attribute which represents the location of the interesting object. This is denoted as  $IN_{t-l}(x_l, y_l)$ .

Table 3 presents samples of interesting objects that are used throughout this paper. There are three types of interesting objects with each type having three objects. The  $IN_{2-3}$  in Table 3 represents the third object of type 2 with location (11.5, 3).

**Definition 11 (Closest Property of a  $t$  Type of Interesting Objects):** Given a set of  $t$  type of interesting objects,  $IN_t = \{IN_{t-1}, IN_{t-2}, \dots, IN_{t-n}\}$ ,  $IN_{t-i}$  is said to be the closest interesting object to an object  $o_j \in O$  if and only if  $\min(dist(o_j, IN_{t-1}), dist(o_j, IN_{t-2}), \dots, dist(o_j, IN_{t-n})) = dist(o_j, IN_{t-i})$ . Therefore, each object  $o_j \in O$  is associated with  $p$  interesting objects that are closest to it, with their distances captured.

The *Definition 3 Dominance*, *Definition 6 Dominance in a Space*, and *Definition 7 Skylines of a Space* defined in Part A are extended to incorporate the additional condition on interesting objects. These are reflected in *Definition 12 Dominance*, *Definition 14 Dominance in a Space*, and *Definition 15 Skylines of a Space*, respectively.

*Definition 12 (Dominance)*: Given two objects  $o_i = (s_i, ns_i)$  and  $o_j = (s_j, ns_j) \in O$  where  $i \neq j$ ,  $o_i$  is said to dominate  $o_j$  (denoted by  $o_i < o_j$ ) if and only if the following conditions hold: (1)  $o_i$  non-spatially dominates  $o_j$  ( $o_i <_{ns} o_j$ ), (2)  $o_i$  spatially dominates  $o_j$  ( $o_i <_s o_j$ ), and (3)  $o_i$  dominates  $o_j$  with  $p$ -closest interesting objects ( $o_i <_{s-IN} o_j$ ).

The definition of *non-spatially dominate* (condition (1)) is as given in *Definition 5*; while the definition of *spatially dominate* (condition (2)) is as given in *Definition 6* of Part A. Meanwhile, *Definition 13 p-Closest Dominance* is introduced to cater the condition (3) defined in *Definition 12 Dominance*.

*Definition 13 (p-Closest Dominance)*: Given two objects  $o_i = (s_i, ns_i)$  and  $o_j = (s_j, ns_j) \in O$  where  $i \neq j$ ,  $o_i$  is said to dominate  $o_j$  with  $p$  interesting objects (denoted by  $o_i <_{s-IN} o_j$ ) if and only if  $o_i$  is no worse than  $o_j$  in all the  $p$  types of interesting objects. This is formally written as follows:  $o_i <_{s-IN} o_j$  if and only if  $\forall t \in p, o_i.t \leq o_j.t \wedge \exists s \in p, o_i.s < o_j.s$ .

*Definition 14 (Dominance in a Space)*: Given a bounded space,  $S$  (region, MBR, fragment, area, polygon, etc.), and two objects  $o_i = (s_i, ns_i)$  and  $o_j = (s_j, ns_j) \in O$  where  $i \neq j$  in  $S$ ,  $o_i$  is said to dominate  $o_j$  (denoted by  $o_i < o_j$ ) in  $S$  if and only if (1)  $o_i$  non-spatially dominates  $o_j$  ( $o_i <_{ns} o_j$ ) in  $S$ , (2)  $o_i$  spatially dominates  $o_j$  ( $o_i <_s o_j$ ) in  $S$ , and  $o_i$  dominates  $o_j$  with  $p$ -closest interesting objects ( $o_i <_{s-IN} o_j$ ) in  $S$ .

*Definition 15 (Skylines of a Space)*: An object  $o_i \in O$  in a space  $S$  is a skyline of  $S$  if there are no other objects  $o_j \in O$  where  $i \neq j$  in the space  $S$  that dominates  $o_i$ . In this paper,  $Sky_{G_p}$  is used to denote the skyline set for the group  $G_p$  of a given space  $S$ .

Based on the above definitions, the problem that is tackled by *ERSGU* is formulated as follows:

Given a group of users,  $G_p = \{u_1, u_2, \dots, u_p\}$ , where  $G_p \subset U$ , and the candidate skylines of  $G_p$  in region  $R_p$  denoted as  $CS_{G_p}$ . Find the skylines of a group of users  $G_q = \{u_1, u_2, \dots, u_q\}$  in region  $R_q$ , i.e.  $CS_{G_q}$ , where  $G_q \subset U$ ,  $G_q \neq G_p$ , and  $R_q \cap R_p \neq \emptyset$  by utilising  $CS_{G_p}$  that has been derived for  $G_p$ . Note that the  $CS_{G_p}$  has been derived based on the objects' spatial and non-spatial attributes as well as their  $p$ -closest interesting objects. The area covered to compute the skylines for  $G_q$  that falls in the region  $S_{G_q}$  can be reduced to the area defined by  $S_{G_q} - S_{G_p}$ , while the results of skyline computation that have been performed earlier over the area  $S_{G_p} \cap S_{G_q}$  for  $G_p$  can be avoided by simply utilising the obtained results derived for  $G_p$ , i.e.  $CS_{G_p}$ .

#### IV. THE RSGU FRAMEWORK

This section presents the *Region-based Skyline for a Group of Users (RSGU)* framework that is mainly proposed for processing the skyline queries of a given group of users;

with two main aims that are (i) avoiding the process of rescanning the set of objects within a predetermined region that is known to have been previously visited by a group of users and (ii) avoiding the recomputation of skylines of a set of objects within a predetermined region that has been analysed in earlier computations of previously visited group of users. Consequently, the number of pairwise comparisons and skyline computation time can be greatly reduced. To achieve these aims, each region is partitioned into smaller units called fragments in such a way that overlapping areas between the currently and previously visited regions can be easily determined; hence rescanning the set of objects within these overlapping areas can be avoided. The results of computing the skylines of each fragment, known as fragment skylines, are saved to be utilised by the subsequent requests. This avoids the need to recompute the skylines of subsequent requests that fall within the same region. In this work, skylines objects are point of interests (PoIs) that are recommended to a group of users that are derived by analysing both the locations of the users, i.e. spatial attributes, as well as the spatial and non-spatial attributes of the set of objects that is within a predetermined region of the group of users.

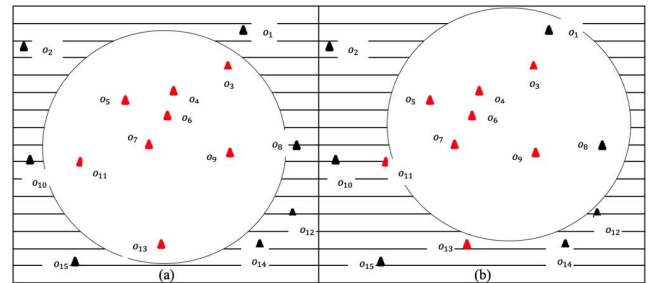
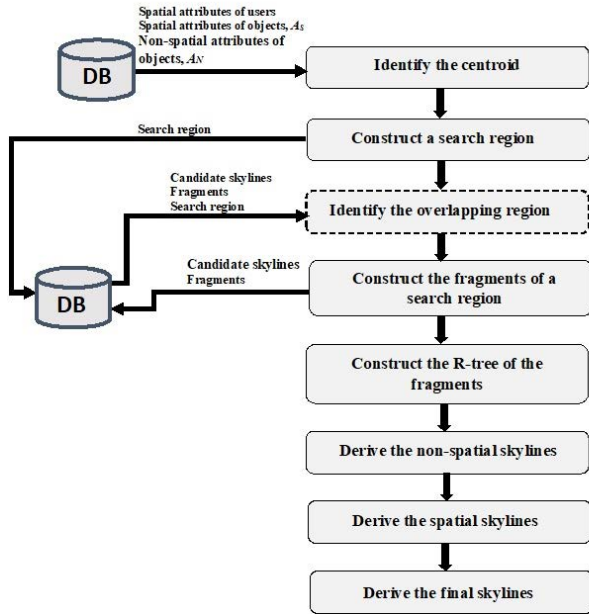


FIGURE 2. Example of (a) previously analysed region (b) current region.

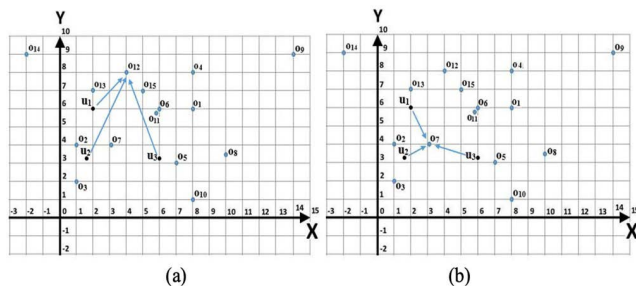
Fig. 2 simulates a sample of situation considered in this paper. There are 15 distinct objects representing point of interests (PoIs). Fig. 2(a) presents the region that is derived based on the locations of a group of users while Fig. 2(b) presents the region that is derived based on the locations of a different group of users. For simplicity, these users are not depicted in the figure. The derivation of these regions is explained in the following sections. Here, the region presented in Fig. 2(b) is considered as the current visited region while Fig. 2(a) represents the previously visited region. As shown in Fig. 2(a) the set of objects that falls within the derived region is  $\{o_3, o_4, o_5, o_6, o_7, o_9, o_{11}, o_{13}\}$ . These objects are then compared to recognise the final skylines to be recommended to the group of users. Assume that the object  $o_7$  is the skyline object. Meanwhile, based on the region depicted in Fig. 2(b), the set of objects that needs to be analysed in order to derive the skyline objects is  $\{o_1, o_3, o_4, o_5, o_6, o_7, o_8, o_9\}$ . From these figures, it is obvious that both regions cover similar area and contain several common objects, namely:  $\{o_3, o_4, o_5, o_6, o_7, \text{ and } o_9\}$ . Attempting to rescan the objects of previously visited region and recompute the skylines of the regions (i.e. repeating the process of pairwise

comparisons among objects) are undoubtedly unwise and costly. In this example, to derive the skylines of  $\{o_1, o_3, o_4, o_5, o_6, o_7, o_8, o_9\}$  (Fig. 2(b)) would require pairwise comparisons to be performed between these objects; while the objects  $o_3, o_4, o_5, o_6, o_7,$  and  $o_9$  have been compared while they are being analysed in identifying the skyline objects of the region presented in Fig. 2(a).



**FIGURE 3.** The Region-based Skyline for a Group of Users (RSGU) framework.

The *RSGU* framework is presented in Fig. 3. It consists of eight main steps that are: (1) Identify the centroid, (2) Construct a search region, (3) Identify the overlapping region, (4) Construct the fragments of a search region, (5) Construct the *R*-tree of the fragments, (6) Derive the non-spatial skylines, (7) Derive the spatial skylines, and (8) Derive the final skylines. Step (3) is conducted only when past computed skyline results of the fragments are available. Each of these steps is elaborated in the following sections.



**FIGURE 4.** (a) The direction of movements of a group of users towards a point which is not the centre point (b) The direction of movements of a group of users towards a point which has the tendency to be the centre point.

**A. IDENTIFY THE CENTROID**

When a group of users,  $G_p = \{u_1, u_2, \dots, u_p\}$ , decided to meet, there must be a point to guide the direction of their movements. Fig. 4 shows some examples of direction of movements of a group of users towards a targeted point. In Fig. 4(a), the users  $u_1, u_2,$  and  $u_3$  are moving towards the object  $o_{12}$  while in Fig. 4(b), these users are moving towards the object  $o_7$ .

In this work, it is assumed that the group of users will move towards a point that has the tendency to be a center based on the users' locations. This point is called centroid and is denoted by  $C_{G_p}(x_{G_p}, y_{G_p})$ . The centroid of a given group of users,  $C_{G_p}$ , is determined using the following formula [84], [85]:

$$C_{G_p}(x_{G_p} = \frac{\sum_{i=1}^n x_i}{n}, y_{G_p} = \frac{\sum_{i=1}^n y_i}{n}) \tag{1}$$

where  $x_i$  is the  $x$  coordinate of user  $u_i$  location,  $y_i$  is the  $y$  coordinate of user  $u_i$  location,  $x_{G_p}$  is the average of the  $x$  coordinates of all users in the group  $G_p$ , and  $y_{G_p}$  is the average of the  $y$  coordinates of all users in the group  $G_p$ . Based on the example given in Table 2, the centroid of  $G_a$  is  $C_{G_a}(8, 9.6)$ .

**B. CONSTRUCT A SEARCH REGION**

The aim of constructing a search region is to limit the searching space to those spaces in which potential candidate skyline objects are derived. Hence, given a group of users, the searching space should include the regions of interest of all users in the group. This is achieved by: (1) identifying the search region for each user,  $S_{u_i}$  and (ii) identifying the search region given a group of users,  $S_{G_p}$ .

*Identify the search region for each user,  $S_{u_i}$*  – Since the centroid of a given group of users, say  $C_{G_p}$ , which is identified in the previous step does not necessarily contain an object, therefore the nearest object,  $o_n$ , to the centroid  $C_{G_p}$  will have to be determined. The nearest object is an object with the shortest Euclidean distance from the centroid, i.e.  $\{o_n | o_n \in O \wedge \forall o_i \in O - \{o_n\}: Ed(C_{G_p}, o_n) < Ed(C_{G_p}, o_i)\}$  where  $Ed$  is the Euclidean distance function. Based on the example given in Table 1(a), the nearest object to the centroid of  $G_a$ , i.e.  $C_{G_a}(8, 9.6)$ , is  $o_8(8, 9.7)$ . The search region for a user,  $u_i$ , denoted as  $S_{u_i}$ , is the area bounded by a rectangle also known as the minimum bounding rectangle,  $MBR_{u_i}$ . The notation  $S_{u_i}$  is used to denote the search region of  $u_i$  while  $MBR_{u_i}$  is used in forming the  $S_{u_i}$ . The distance between a user,  $u_i$ , and the nearest object,  $o_n$ , denoted by  $R_{u_i, o_n}$ , is calculated by the following equation:

$$R_{u_i, o_n} = \sqrt{(x_{o_n} - x_i)^2 + (y_{o_n} - y_i)^2} \tag{2}$$

where  $x_i$  is the  $x$  coordinate of user  $u_i$  location,  $y_i$  is the  $y$  coordinate of user  $u_i$  location,  $x_{o_n}$  is the  $x$  coordinate of object  $o_n$  location, and  $y_{o_n}$  is the  $y$  coordinate of object  $o_n$  location. A *MBR* is formed based on four vertices as explained in the following: the vertex at the bottom left of the *MBR* is denoted by  $bl = (x_{bl}, y_{bl})$ ; the vertex at the bottom right of the *MBR* is denoted by  $br = (x_{br}, y_{br})$ ; the vertex at the top left of the

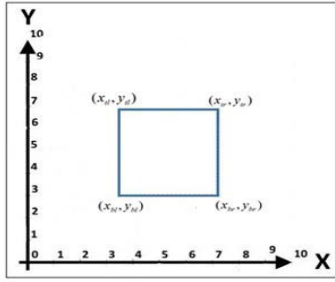


FIGURE 5. Minimum bounding rectangle (MBR).

MBR is denoted by  $tl = (x_{tl}, y_{tl})$ ; and the vertex at the top right of the MBR is denoted by  $tr = (x_{tr}, y_{tr})$ . Fig. 5 depicts these notations. These vertices are calculated as follows:

$$\begin{aligned} bl &= (x_i - R_{u_i o_n}, y_i - R_{u_i o_n}) \\ br &= (x_i + R_{u_i o_n}, y_i - R_{u_i o_n}) \\ tl &= (x_i - R_{u_i o_n}, y_i + R_{u_i o_n}) \\ tr &= (x_i + R_{u_i o_n}, y_i + R_{u_i o_n}) \end{aligned}$$

Identify the search region given a group of users,  $S_{G_p}$  – This step is simply achieved by performing union on the search region of each user in the group, i.e.  $S_{G_p} = \bigcup_{i=1}^p S_{u_i}$ . An example of a search region  $S_{G_a} = \bigcup_{i=1}^3 S_{u_i}$  can be seen in Fig. 6. The search region constructed in this step is saved to be utilised later in identifying the overlapping region of subsequent requests.

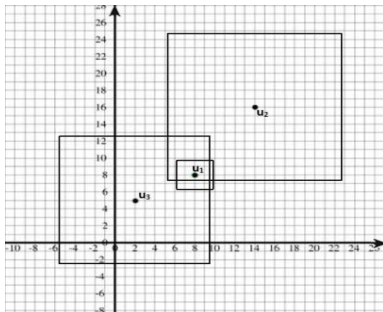


FIGURE 6. The search region for a group of users.

### C. CONSTRUCT THE FRAGMENTS OF A SEARCH REGION

This step partitions the search region of a group of users,  $S_{G_p}$ , into  $m$  fragments (subspaces). Here, the vertices of the MBR associated to each  $S_{u_i}$  are analysed and sorted according to the  $x$ - and  $y$ -axes. The search region (space) is vertically fragmented based on the  $x$ -coordinates, while it is horizontally fragmented based on the  $y$ -coordinates. The MRBs formed within the  $S_{G_p}$  are the fragments of the region.

Objects that fall within each fragment are then identified. Given an object,  $o_j(x_j, y_j)$ , and a fragment,  $F_k$ , with  $bl(x_l, y_b)$ ,  $br(x_r, y_b)$ ,  $tl(x_l, y_t)$ , and  $tr(x_r, y_t)$ , the following cases are identified:

- (a) If  $x_l < x_j < x_r$  and  $y_b < y_j < y_t$ , then the object  $o_j(x_j, y_j)$  is said to fall within the boundary of fragment,  $F_k$ .
- (b) If  $x_j = x_l$  or  $x_j = x_r$  or  $y_j = y_b$  or  $y_j = y_t$ , then the object  $o_j(x_j, y_j)$  is said to intersect with the boundary of fragment,  $F_k$ .
- (c) Objects that do not meet the above two cases are objects that are outside the boundary of fragment,  $F_k$ .

Further, utilising the non-spatial dominance testing given in Definition 8, an extension to the Definition 4, over the set of objects that satisfies the cases (a) or (b) above, denoted by  $O_{F_k}$ , the non-spatial candidate skylines of a fragment are determined,  $CS_{nsF_k}$ , as defined by Definition 9.

**Definition 8 (Non-Spatial Dominance of the Fragment  $F_k$ ):** Given two objects  $o_i = (s_i, ns_i) \in O_{F_k}$  and  $o_j = (s_j, ns_j) \in O_{F_k}$  where  $i \neq j$ ,  $o_i$  is said to non-spatially dominate  $o_j$  (denoted by  $o_i <_{ns} o_j$ ) if and only if  $o_i$  is no worse than (in this definition, greater value is preferable)  $o_j$  in all the non-spatial attributes,  $A_N$ . This is formally written as follows:  $o_i <_{ns} o_j$  if and only if  $\forall d_k \in A_N, o_i.d_k \geq o_j.d_k \wedge \exists d_l \in A_N, o_i.d_l > o_j.d_l$ .

**Definition 9 (Candidate Skylines of the Fragment  $F_k$ ):** An object  $o_i \in O_{F_k}$  in a space  $F_k$  is a non-spatial candidate skyline of  $F_k$  if there are no other objects  $o_j \in O_{F_k}$  where  $i \neq j$  in the space  $F_k$  that non-spatially dominates  $o_i$ .

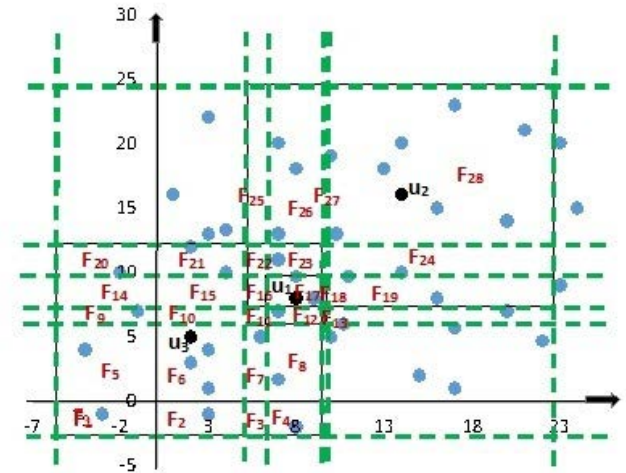


FIGURE 7. The fragments derived based on the  $S_{G_a}$  given in Fig. 6.

This will avoid rescanning the objects of the region and repeating the process of pairwise comparisons among the objects during the skyline computation of subsequent skyline queries. Fig. 7 presents the fragments constructed based on the  $S_{G_a}$  given in Fig. 6. The  $x$ -coordinates =  $\{-5.6, 0, 5.3, 6.3, 9.6, 9.7, 22.7\}$  and  $y$ -coordinates =  $\{-2.6, 0, 6.3, 7.3, 9.7, 12.6, 24.7\}$ . Altogether there are 28 fragments; some samples are given in Table 4.

### D. CONSTRUCT THE R-TREE OF THE FRAGMENTS

The aim of constructing the  $R$ -tree of the fragments is to reduce the searching process time. An  $R$ -tree is a classified data structure and it is used for dynamic classification of a set



TABLE 4. Sample of fragments and their associated candidate skylines.

$x$ Coordinate ( $x_l, x_r$ )	$y$ Coordinate ( $y_b, y_t$ )	$bl(x_l, y_b)$	$br(x_r, y_b)$	$tl(x_l, y_t)$	$tr(x_r, y_t)$	Fragment, $F_k$	Objects	Candidate skylines, $CS_{nsF_k}$
-5.6, 0	-2.6, 0	-5.6, -2.6	0, -2.6	-5.6, 0	0, 0	$F_1$	$o_{38}$	$o_{38}$
0, 5.3	-2.6, 0	0, -2.6	5.3, -2.6	0, 0	5.3, 0	$F_2$	$o_{45}$	$o_{45}$
5.3, 6.3	-2.6, 0	5.3, -2.6	6.3, -2.6	5.3, 0	6.3, 0	$F_3$	-	-
6.3, 9.6	-2.6, 0	6.3, -2.6	9.6, -2.6	6.3, 0	9.6, 0	$F_4$	$o_{42}$	$o_{42}$
...	...	...	...	...	...	...	...	...
0, 5.3	0, 6.3	0, 0	5.3, 0	0, 6.3	5.3, 6.3	$F_6$	$o_1, o_2, o_3$	$o_2, o_3$
...	...	...	...	...	...	...	...	...
9.7, 22.7	7.3, 24.7	9.7, 7.3	22.7, 7.3	9.7, 24.7	22.7, 24.7	$F_{28}$	$o_{26}, o_{27},$ $o_{29}, o_{30},$ $o_{31}, o_{32},$ $o_{33}, o_{40}$	$o_{29}, o_{32}$

of  $d$ -dimensional coordination, rectangles or objects demonstrating them by the minimum bounding  $d$ -dimensional rectangles. Each node of the  $R$ -tree relates to the  $MBR$  that confines its children. An  $R$ -tree of order  $(m, M)$  considering each leaf node can have up to  $M$  entries, while the minimum permitted number of entries is  $m \leq M/2$ . All leaf nodes of the  $R$ -tree are at the same level. The  $R$ -tree is constructed based on the algorithms proposed by [84]. In this work, the following algorithms are utilised during the construction of the  $R$ -tree: *Search*, *Insert*, *ChooseLeaf*, *SplitNode*, and *AdjustTree*.

Based on Fig. 7 the  $MBRs$  constructed are as shown in Fig. 8 while Fig. 9 shows the  $R$ -tree derived based on the  $MBRs$  of the fragments.

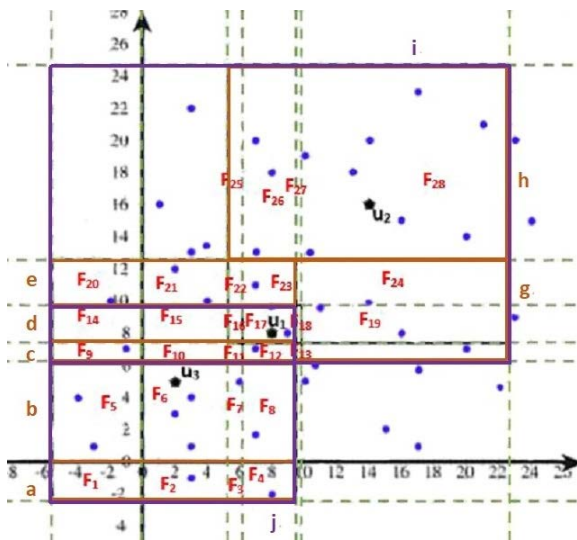


FIGURE 8. The  $MBRs$  derived based on the fragments of the  $S_{G_d}$  given in Fig. 7.

**E. DERIVE THE NON-SPATIAL SKYLINES**

This step performs the non-spatial dominance testing given in Definition 8 Non-spatial Dominance of the Fragment  $F_k$  towards the  $CS_{nsF_i}$  lists derived in the previous step, Construct the Fragments of a Search Region, presented in Part C

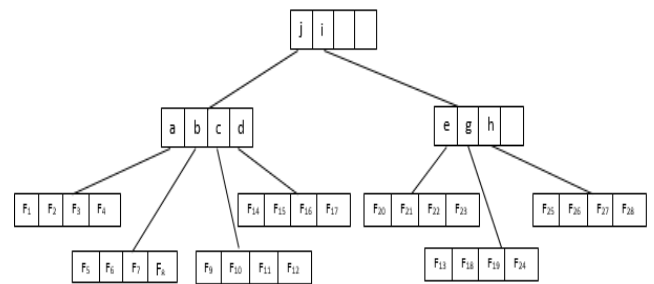


FIGURE 9. The  $R$ -tree derived based on the  $MBRs$  of the fragments given in Fig. 8.

to generate the non-spatial skylines of a given group of users. In other words, the pairwise comparisons are only performed between objects that are the candidate skylines of a fragment. The objects that non-spatially dominate the other objects, given the  $CS_{nsF_i}$  lists where  $i = \{1, 2, \dots, 28\}$  in Table 4 are  $o_{18}$  and  $o_{20}$ , thus  $Sky_{nsG_d} = \{o_{18}, o_{20}\}$ .

**F. DERIVE THE SPATIAL SKYLINES**

This step applies the spatial dominance testing given in Definition 5 Spatial Dominance towards the  $CS_{nsF_i}$  lists. First, the distance between each object,  $o_i$ , and each user,  $u_j$ , is determined, denoted as  $o_i - u_j$ . Given a group of  $l$  users, there will be  $l$  values of distances with regard to the object  $o_i$ , i.e.  $o_i - u_1, o_i - u_2, \dots, o_i - u_l$ . These values are treated as the values of dimensions to be used in the spatial skyline computation. Then, the total distance between an object,  $o_i$ , to each user,  $u_j$ , in the group of users is calculated and saved into a parameter named *Sum Distance* -  $o_i$ , i.e.  $Sum\ Distance - o_i = \sum_{p=1}^l o_i - u_p$ . The value of *Sum Distance* -  $o_i$  is used as a selection criterion in determining the object that should be considered in each iteration of the spatial skyline computation. The smallest value of *Sum Distance* -  $o_i$  indirectly indicates that most users in the group are close to the object  $o_i$  and has more chances to dominate the other objects.

An example is shown in Table 5 where the attributes  $o_i - u_1, o_i - u_2,$  and  $o_i - u_3$  are the distances of each object,

**TABLE 5.** The distance and sum distance of each object in  $CS_{nsF_t}$ .

Restaurant	$o_i - u_1$	$o_i - u_2$	$o_i - u_3$	Sum Distance - $o_i$
$o_1$	7.81	17.69	2	27.5
$o_2$	6.4	16.27	1.41	24.08
$o_3$	8.6	18.6	4.12	31.32
$o_4$	6.37	15.92	5.99	28.28
$o_5$	3.6	13.6	4	21.2
$o_6$	1.41	11.4	5.38	18.19
$o_7$	1	9.43	7.61	18.04
$o_8$	1.7	8.7	7.62	18.02
$o_9$	3.16	8.6	7.81	19.57
$o_{18}$	8	8.24	14.31	30.55
$o_{19}$	6.32	6	13	25.32
$o_{20}$	3.44	6.97	10.15	20.56
$o_{21}$	4.47	11.66	5.38	21.51
$o_{22}$	7.21	12.64	7	26.85
$o_{25}$	5.09	7.61	9.43	22.13
$o_{26}$	10.63	2.23	17.2	30.06
$o_{27}$	13.41	6.32	20.12	39.85
$o_{29}$	18.38	8.6	24.83	51.81
$o_{30}$	18.6	7.61	23.43	49.64
$o_{31}$	13.41	4	19.2	36.61
$o_{32}$	11.18	2.23	17.02	30.43
$o_{33}$	11.18	5	16.12	32.3
$o_{36}$	12.04	8.06	15.81	35.91
$o_{38}$	14.21	24.04	7.81	46.06
$o_{39}$	9.05	17.49	3.6	30.14
$o_{40}$	5.5	4.76	11.52	21.78
$o_{41}$	12.64	21.63	6.08	40.35
$o_{42}$	10	18.97	9.21	38.18
$o_{43}$	10	6.32	14.31	30.63
$o_{44}$	10.19	17.08	6.4	33.67
$o_{45}$	8.6	20.24	6.08	34.92

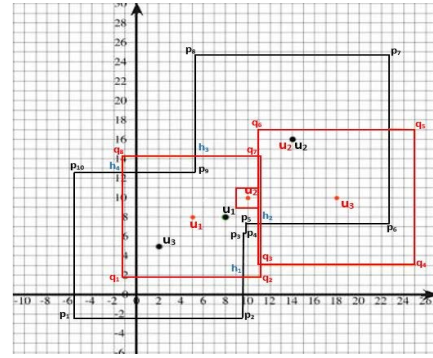
$o_i$ , and each user  $u_1$ ,  $u_2$ , and  $u_3$ , respectively; meanwhile the *Sum Distance -  $o_i$*  attribute presents the total distance of an object to each user in the group of users. Here,  $o_8$  will be the first object selected which is then followed by  $o_7$ . The objects that spatially dominate the other objects, given the  $CS_{nsF_t}$  lists in Table 4 are as listed in  $Sky_{S_{G_a}} = \{o_2, o_5, o_6, o_7, o_8, o_9, o_{20}, o_{25}, o_{26}, o_{32}, o_{40}\}$ .

### G. DERIVE THE FINAL SKYLINES

This is the final step that combines the results produced in the steps presented in Parts E and F above. Based on Definition 7 *Skylines of a Space*, the final skylines for a given group  $G_i$  is given by,  $Sky_{G_i} = Sky_{nsG_i} \cup Sky_{S_{G_i}}$ . Thus, the final skylines for the group  $G_a$ ,  $Sky_{G_a} = \{o_2, o_5, o_6, o_7, o_8, o_9, o_{18}, o_{20}, o_{25}, o_{26}, o_{32}, o_{40}\}$ .

### H. IDENTIFY THE OVERLAPPING REGION

This step constructs the overlapping region,  $O_R$ , between the search regions of two groups of users, say  $S_{G_i}$  and  $S_{G_j}$ . Here, it is assumed that the results of the skyline queries of a group of users, say  $G_i$ , have been derived. Thus, the overlapping region indicates that the region has been scanned and it is unwise to scan it again. Fig. 10 shows two search regions,  $S_{G_a}$ , the polygon with black border line and  $S_{G_b}$ , the polygon with red border line which represent the search region of group  $G_a$  and group  $G_b$ , respectively. If there are more than one search region that are available,  $\{S_{G_i}, S_{G_k}, \dots, S_{G_m}\}$ , then the overlapping area between  $S_{G_j}$  and each of the available search

**FIGURE 10.** The overlapping region between  $S_{G_a}$  and  $S_{G_b}$ .

region is analysed and the region with the highest percentage of overlapping area is selected in this step.

To identify the overlapping region, the following steps are performed:

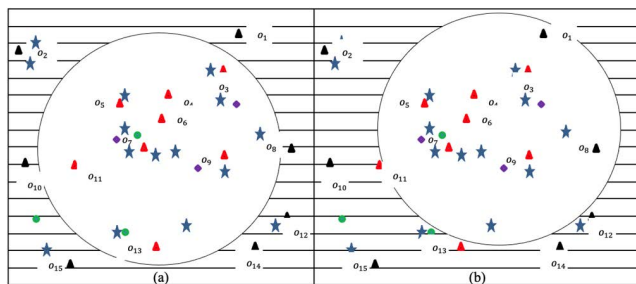
- (1) Get the polygon's vertices of  $S_{G_i}$ . Based on the example,  $S_{G_a} = \{p_1, p_2, p_3, p_4, p_5, p_6, p_7, p_8, p_9, p_{10}\}$ . Note that for simplicity, the coordinates of the vertices are omitted.
- (2) Get the polygon's vertices of  $S_{G_j}$ . Based on the example,  $S_{G_b} = \{q_1, q_2, q_3, q_4, q_5, q_6, q_7, q_8\}$ .
- (3) Get the vertices of  $S_{G_i}$  that are also in  $S_{G_j}$ . Based on the example,  $I_{G_a-G_b} = \{p_3, p_4, p_5, p_6, p_9\}$ .
- (4) Get the vertices of  $S_{G_j}$  that are also in  $S_{G_i}$ . Based on the example,  $I_{G_b-G_a} = \{q_1, q_6, q_7\}$ .
- (5) Get the coordinates where the edges of  $S_{G_i}$  and  $S_{G_j}$  meet. Based on the example,  $H = \{h_1(9.6, 1.8), h_2(11.2, 7.3), h_3(22.7, 17), h_4(5.3, 14.2), h_5(-1.2, 12.6)\}$ .
- (6) The overlapping region,  $O_R$ , is defined as a polygon derived based on the following vertices:  $I_{G_a-G_b} \cup I_{G_b-G_a} \cup H$ . Based on the example,  $O_R = \{h_1, p_3, p_4, p_5, h_2, p_6, h_3, q_6, q_7, h_4, p_9, h_5, q_1\}$ .

Once the  $O_R$  has been defined, the fragments derived in the earlier step are analysed. Those fragments that fall within the  $O_R$ ; are retrieved together with their candidate skylines,  $CS_{nsF_k}$ . To search for the fragments that are within the  $O_R$ , the  $R$ -tree constructed in Part D is traversed starting from the root node. The search is based on the depth-first search tree traversal algorithm. The following rules are applied: (i) If the visited node,  $v_n$ , is an internal node (example, the node with the  $a, b, c, d$  entries of Fig. 9), then each entry of the internal node,  $e_i.v_n$ , is examined (example  $a$ ). If the entry (example  $a$ ) overlaps with  $O_R$ , then rules (i) and (ii) are applied accordingly over the entry,  $e_i.v_n$ ; (ii) If the visited node,  $v_n$ , is a leaf node (example, the node with the  $F_1, F_2, F_3, F_4$  entries of Fig. 9), then each entry of the leaf node,  $e_i.v_n$ , is examined. If the entry overlaps with  $O_R$ , then it is one of the fragments that will be retrieved for further analysis. Hence, scanning this area is no longer necessary. While for the non-overlapping area, denoted as  $\neg O_R$ , the following steps as discussed above will be conducted: (4) *Construct the fragments of the non-overlapping region,*

i.e.  $\neg O_R = S_{G_j} - O_R$  (5 and 6) Derive non-spatial skylines and spatial skylines, respectively by considering both the lists  $CS_{O_R}$  and  $CS_{\neg O_R}$ , and (7) Derive the final skylines.

**V. THE ERSGU FRAMEWORK**

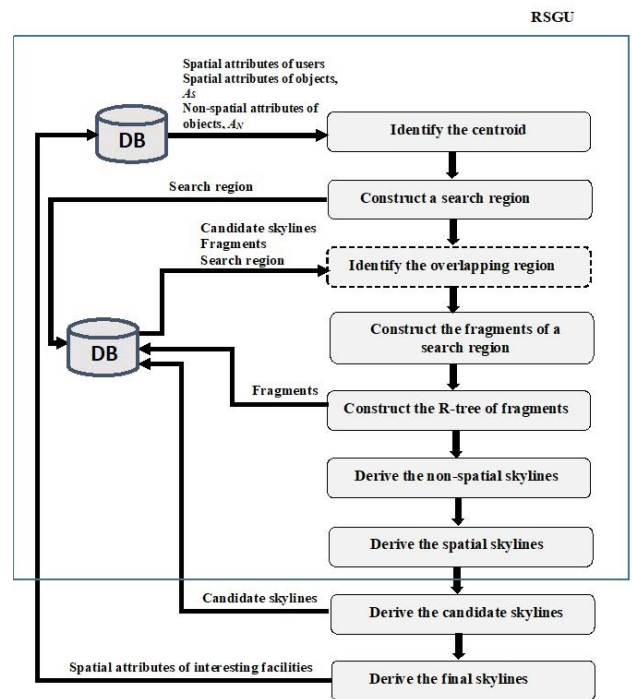
This section presents the *Extended Region-based Skyline for a Group of Users (ERSGU)* framework, an enhancement of the previous framework, *RSGU*, presented in Section IV. *RSGU* is designed for processing the skyline queries of a given group of users. *RSGU* attempts to avoid the process of rescanning the set of objects and the recomputation of skylines of a set of objects that is within a predetermined region that is known to have been previously visited by a group of users. As a result, the number of pairwise comparisons and skyline computation time can be greatly reduced. *ERSGU* has similar aims as *RSGU* with an additional feature. The skylines derived for a group of users by *ERSGU* are not only based on the evaluation of the spatial and non-spatial attributes of the objects that are within the predetermined region, but also the closeness of the objects to the desirable facilities or other interesting objects in the region. An example of object of interest is hotel while other desirable facilities/interesting objects are clinic, bus station, airport, etc. This would benefit the users, since they might want to visit a place where there are several other useful facilities or interesting objects (for simplicity, henceforth referred to as *interesting objects*) nearby. Hence, the skylines, which are the objects recommended to be visited by the group of users, are derived by analysing both the locations of the users, i.e. spatial attributes, as well as the spatial and non-spatial attributes of the objects along with the closeness of the objects to other interesting objects.



**FIGURE 11.** Example of (a) previously analysed region (b) current region with interesting objects/facilities.

Fig. 11 simulates a sample of situation considered in this paper which is similar to the sample of situation described in Fig. 2. Besides, the 15 distinct objects representing objects of interest ( $\blacktriangle$ ), several other interesting objects are also presented. The symbols ( $\star$ ), ( $\bullet$ ), and ( $\blacklozenge$ ) in the figure represent the distinct types of interesting objects. Fig. 11(a) presents the previously visited region that is derived based on the locations of a group of users in which the skylines of the region have been derived; while Fig. 11(b) presents the current visited region that is derived based on the locations of a different group of users in which the skylines are to be identified. As shown in Fig. 11(a) and Fig. 11(b) the sets of objects

that fall within the derived regions are  $\{o_3, o_4, o_5, o_6, o_7, o_9, o_{11}, o_{13}\}$  and  $\{o_1, o_3, o_4, o_5, o_6, o_7, o_8, o_9\}$ , respectively. Each set of objects is then compared not only based on the objects’ spatial and non-spatial attributes but also how close are they to the other interesting objects before the final skylines are recommended to the group of users. The object  $o_7$  for instance is near to several interesting objects compared to the objects  $o_4$  and  $o_6$ . From these figures, it is obvious that both regions cover similar area and contain several common objects, namely:  $o_3, o_4, o_5, o_6, o_7$ , and  $o_9$ . Hence, attempting to rescan the objects of previously visited region and recompute the skylines of the region (i.e. repeating the process of pairwise comparisons among objects) are undoubtedly unwise and costly. This is because not only the objects’ spatial and non-spatial attributes will be analysed again but also all the other interesting objects that are close to the objects need to be reexplored.



**FIGURE 12.** The extended region-based skyline for a Group of users (ERSFU) framework.

The *ERSFU* framework is presented in Fig. 12. It consists of nine main steps that are: (1) Identify the centroid, (2) Construct a search region, (3) Identify the overlapping region, (4) Construct the fragments of a search region, (5) Construct the *R*-tree of the fragments, (6) Derive the non-spatial skylines, (7) Derive the spatial skylines, (8) Derive the candidate skylines, and (9) Derive the final skylines. Steps (1) till (7) are the same steps as *RSGU*, while Step (8) *Derive the final skylines* of *RSGU* is renamed as (8) *Derive the candidate skylines* without any changes with regard to the step. Step (9) *Derive the final skylines* is the new step incorporated into *ERSFU*. Hence, only Step (8) and Step (9) of *ERSFU*

will be further discussed while the earlier seven steps are as explained in Section IV.

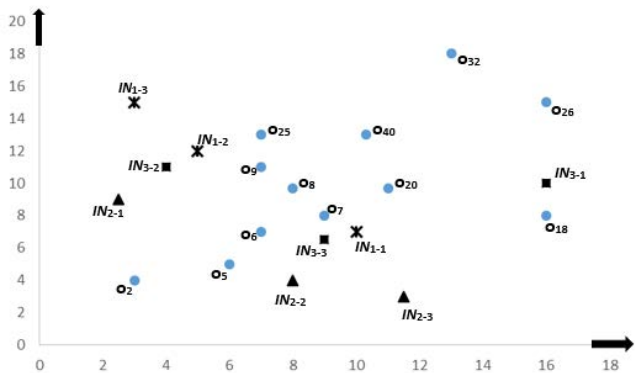
**A. DERIVE THE CANDIDATE SKYLINES**

This step combines the results produced in the steps presented in Section IV Parts E and F. Based on Definition 7, the candidate skylines for a given group  $G_i$  is given by,  $C - Sky_{G_i} = Sky_{nsG_i} \cup Sky_{sG_i}$ . Thus, referring to the same example given in Section IV, the candidate skylines for the group  $G_a$ ,  $C - Sky_{G_a} = \{o_2, o_5, o_6, o_7, o_8, o_9, o_{18}, o_{20}, o_{25}, o_{26}, o_{32}, o_{40}\}$ .

**B. DERIVE THE FINAL SKYLINES**

The final step of *ERSGU* attempts to derive the final skylines based on the list of candidate skylines produced in Step (8). The candidate skylines are those objects that are not dominated by any other objects based on the two conditions defined in Definition 12 Dominance. These conditions are (1)  $o_i$  non-spatially dominates  $o_j$  ( $o_i <_{ns} o_j$ ) and (2)  $o_i$  spatially dominates  $o_j$  ( $o_i <_s o_j$ ). The final skylines are derived based on the candidate skylines obtained above in which condition (3) of Definition 12 Dominance is satisfied, i.e.  $o_i$  dominates  $o_j$  with  $p$ -closest interesting objects ( $o_i <_{s-IN} o_j$ ).

To clarify this step, consider the group of users,  $G_a$ , and its search region,  $S_{G_a}$ , as given in Table 1 and Fig. 6, respectively. Assume that the interesting objects that fall within the region of  $S_{G_a}$  are as presented in Fig. 13 with their detail locations presented in Table 3. In this example, the following are assumed:



**FIGURE 13.** Locations of the interesting objects in the region  $S_{G_a}$ .

- 1) Only candidate skylines produced in Step (8) are presented in the figure, i.e.  $C - Sky_{G_a} = \{o_2, o_5, o_6, o_7, o_8, o_9, o_{18}, o_{20}, o_{25}, o_{26}, o_{32}, o_{40}\}$ . These objects are denoted as circles.

- 2) There are three types of interesting objects symbolised by star (type 1), triangle (type 2), and rectangular (type 3).

Each type of interesting object has three objects. The notation  $IN_{t-l}$  is used to denote the  $l$ -th interesting object of type  $t$ . Hence,  $IN_1 = \{IN_{1-1}, IN_{1-2}, IN_{1-3}\}$ ,  $IN_2 = \{IN_{2-1}, IN_{2-2}, IN_{2-3}\}$ , and  $IN_3 = \{IN_{3-1}, IN_{3-2}, IN_{3-3}\}$  represent the set of interesting objects of type 1, 2, and 3,

**TABLE 6.** The distances between each  $C - Sky_{G_a}$  and  $IN$ .

ID	$IN_{1-1}$	$IN_{1-2}$	$IN_{1-3}$	$IN_{2-1}$	$IN_{2-2}$	$IN_{2-3}$	$IN_{3-1}$	$IN_{3-2}$	$IN_{3-3}$
$o_2$	7.61	8.24	11	5.02	5	8.55	14.31	7.07	6.5
$o_5$	4.47	7.07	10.44	5.31	2.23	5.85	11.18	6.32	3.35
$o_6$	3	5.38	8.94	4.92	3.16	6.02	9.48	5	2.06
$o_7$	1.41	5.65	9.21	6.57	4.12	5.59	7.28	5.83	1.5
$o_8$	3.36	3.78	7.28	5.54	5.7	7.55	8	4.2	3.35
$o_9$	5	2.23	5.65	4.92	7.07	9.17	9.05	3	4.92
$o_{18}$	6.08	11.70	14.76	13.53	8.94	6.72	2	12.36	7.15
$o_{20}$	2.87	6.42	9.59	8.52	6.44	6.71	5	7.11	3.77
$o_{25}$	6.70	2.23	4.47	6.02	9.05	10.96	9.48	3.60	6.80
$o_{26}$	10	11.40	13	14.77	13.60	12.81	5	12.64	11.01
$o_{32}$	11.40	10	10.44	13.82	14.86	15.07	8.54	11.40	12.17
$o_{40}$	6	5.39	7.56	8.76	9.28	10.07	6.44	6.60	6.62

respectively. The locations of each of these objects are as given in Table 3.

Given a set of candidate skylines for a group of users,  $C - Sky_{G_i} = \{o_1, o_2, \dots, o_y\}$  and a set of  $p$  distinct type of interesting objects,  $IN = \{IN_1, IN_2, \dots, IN_p\}$ , the following steps are performed:

- (a) For each  $o_j \in C - Sky_{G_i}$  and for each set of  $t$  type of interesting objects,  $IN_t = \{IN_{t-1}, IN_{t-2}, \dots, IN_{t-n}\}$ , get the Euclidean distance between  $o_j$  and  $IN_{t-i}$ ,  $dist(o_j, IN_{t-i})$ . Hence, if there are  $y$  objects of candidate skylines and  $p$  types of interesting objects with each type having approximately  $n$  objects, then the number of distances that needs to be measured  $\approx y \times p \times n$ . Table 6 shows the distances measured for each  $o_j \in C - Sky_{G_a}$  and each  $IN$ . Here, there are  $12 \times 9 = 108$  distances that are captured.
- (b) For each  $o_j \in C - Sky_{G_i}$  and for each set of  $t$  type of interesting objects,  $IN_t = \{IN_{t-1}, IN_{t-2}, \dots, IN_{t-n}\}$ , the closest interesting object to the object  $o_j$  is determined based on Definition 11 Closest Property of a  $t$  Type of Interesting Objects. The distance of the closest interesting object, say  $IN_{t-i}$ , to  $o_j$  is then captured. This will produce an object  $o_j$  with  $p$  values of distances denoted as  $o_j (cd_1, cd_2, \dots, cd_p)$  with  $cd_h$  represents the closest distance of an interesting object of type  $h$  to the object  $o_j$ . Here, each  $p$  type can be regard as a dimension while  $cd_p$  is a value of the dimension  $p$ . By applying Definition 11 Closest Property of a  $t$  Type of Interesting Objects to the example given in Table 6, the closest object of each type to each  $C - Sky_{G_a}$  is as given in Table 7. Meanwhile, Table 8 presents the closest distance values of each type to each  $C - Sky_{G_a}$ . For instance, the closest interesting objects of types 1, 2, and 3 to  $o_2$  are  $IN_{1-1}$ ,  $IN_{2-2}$ , and  $IN_{3-3}$ , with distances 7.61, 5, and 6.5, respectively.
- (c) The final skylines are determined by performing the conventional skyline algorithm over the candidate skylines with  $p$  types of interesting objects as dimensions and distances as the values used in the pairwise comparisons. Here, Definition 15 Skylines of a Space is applied. For instance,  $o_6(3, 3.16, 2.06)$  is said to dominate  $o_2(7.61, 5, 6.5)$  since  $\forall t \in \{1, 2, 3\}, o_{6,t} < o_{2,t}$ . However,  $o_9(2.23, 4.92, 3)$  and  $o_{18}(6.08, 6.72, 2)$  are

**TABLE 7.** The closest object of each type to each  $C - Sky_{G_t}$ .

ID	$IN_{1-1}$	$IN_{1-2}$	$IN_{1-3}$	$IN_{2-1}$	$IN_{2-2}$	$IN_{2-3}$	$IN_{3-1}$	$IN_{3-2}$	$IN_{3-3}$
$o_2$	7.61	-	-	-	5	-	-	-	6.5
$o_5$	4.47	-	-	-	2.23	-	-	-	3.35
$o_6$	3	-	-	-	3.16	-	-	-	2.06
$o_7$	1.41	-	-	-	4.12	-	-	-	1.5
$o_8$	3.36	-	-	5.54	-	-	-	-	3.35
$o_9$	-	2.23	-	4.92	-	-	-	3	-
$o_{18}$	6.08	-	-	-	-	6.72	2	-	-
$o_{20}$	2.87	-	-	-	6.44	-	-	-	3.77
$o_{25}$	-	2.23	-	6.02	-	-	-	3.60	-
$o_{26}$	10	-	-	-	-	12.81	5	-	-
$o_{32}$	-	10	-	13.82	-	-	8.54	-	-
$o_{40}$	-	5.39	-	8.76	-	-	6.44	-	-

said not to dominate each other as  $o_9$  is better than  $o_{18}$  in type 1 and type 2 and worse than  $o_{18}$  in type 3. Consequently,  $o_5$  dominates  $o_2$ ,  $o_7$  dominates  $o_8$ ,  $o_9$ ,  $o_{18}$ ,  $o_{20}$ ,  $o_{26}$ ,  $o_{32}$ , an  $o_{40}$ ; while  $o_6$  is not dominated by any other objects. As a result, the final skylines,  $Sky_{G_a} = \{o_5, o_6, o_7\}$ .

**VI. RESULTS AND DISCUSSION**

**A. EXPERIMENTAL SETTINGS**

To fairly evaluate the performance and prove the efficiency of *RSGU* and *ERSGU*, several extensive experiments are designed. These experiments are conducted on Intel Core i7 3.6GHz processor with 32GB of RAM and Windows 8 professional. The implementation of *RSGU* and *ERSGU* was done on VB.NET 2013. The performance results of *RSGU* and *ERSGU* are compared to the *VR* algorithm proposed by [63]. To the best of our knowledge, the works that are closely related to our work are by [60], [61], [62], and [63]; with [63] being the most recent among the list above.

In validating the correctness of *RSGU*, the following is conducted: The *VR* algorithm and the *RSGU* framework were run over a given data set and a group of users,  $G_a$ , to derive a set of skylines,  $Sky_{G_a} - VR$  and  $Sky_{G_a} - RSGU$ , respectively. Then, given another group of users,  $G_b$ , the *VR* algorithm and the *RSGU* framework were run again and the set of skylines produced, namely:  $Sky_{G_b} - VR$  and  $Sky_{G_b} - RSGU$  are recorded. Intuitively, the correctness of *RSGU* is proven as the skyline objects produced by the *VR* algorithm,  $Sky_{G_a} - VR$ , is equal to the skyline objects produced by the *RSGU* framework, i.e.  $Sky_{G_a} - VR = Sky_{G_a} - RSGU$ . Similarly,  $Sky_{G_b} - VR = Sky_{G_b} - RSGU$ .

Meanwhile, the correctness of *ERSGU* framework could not be verified since both the *VR* algorithm and the *RSGU* framework that utilised the same evaluation criteria do not consider the closeness of the objects to other interesting objects. Undoubtedly, the skyline objects produced by the *ERSGU* framework are different from those produced by the *VR* algorithm and the *RSGU* framework.

Two types of data sets are used in the experiments, namely: *synthetic* and *TIGER* data sets. The *synthetic* data set is used to simulate several experimental settings to reflect all possible real settings. The *synthetic* data set is generated in such a way that all objects are independent with uniform distribution.

**TABLE 8.** The closest distance value of type 1, 2, and 3 to each  $C - Sky_{G_a}$ .

ID	$IN_1$	$IN_2$	$IN_3$
$o_2$	7.61	5	6.5
$o_5$	4.47	2.23	3.35
$o_6$	3	3.16	2.06
$o_7$	1.41	4.12	1.5
$o_8$	3.36	5.54	3.35
$o_9$	2.23	4.92	3
$o_{18}$	6.08	6.72	2
$o_{20}$	2.87	6.44	3.77
$o_{25}$	2.23	6.02	3.60
$o_{26}$	10	12.81	5
$o_{32}$	10	13.82	8.54
$o_{40}$	5.39	8.76	6.44

This is in line with the settings used in previous works [60], [61], [63]. In addition, in skyline queries, the *synthetic* data set is commonly used by previous researchers in evaluating the performance of their proposed approaches [70], [4], [81], [72], [73], [54], [51], [55], [59], [52]. On the other hand, the *TIGER* data set is a real data set from the line segment data of Long Beach. It is used by previous works that are related to spatial skyline queries [60], [61], [63]. The set of objects is prepared by extracting the midpoint of each road line segment. The data set contains 50,747 objects. There are 8 different types of objects, namely: *hospital*, *restaurant*, *church*, *school*, *institution*, *building*, *hotel*, and *populated place*. Each type of object has a different number of objects, a different list of non-spatial attributes, and spatial attributes. Skyline computation requires objects to be of same arity while the domination test is performed only on numerical values. Since *rate* and *price* are the only dimensions with numerical values, thus these dimensions are used as the evaluation criteria of the skyline computation.

Each experiment is run 10 times and the average value of these runs is reported. In deriving the set of skylines, it is assumed that lower values are preferable compared to higher ones. The performance measurement used in the experiments is processing time as it is the most commonly used measurement in evaluating the performance of skyline algorithm [63]. The processing time is evaluated for different parameter settings that are:

- (i) number of users in a group – which is varied with minimum 4 users and maximum 25 users,
- (ii) number of groups of users – which is varied with minimum 2 groups and maximum 32 groups,
- (iii) overlapping region – the overlapped regions between different groups of users are controlled with 20%, 40%, 60%, 80%, and 100% of overlapped,
- (iv) space size – the space area of the *synthetic* and *TIGER* data sets is as presented in Table 9,
- (v) density – for the *TIGER* data set, the density of the objects is set to 0.56%, 1.60%, 7%, 15%, and 34%,
- (vi) dimensionality – for both the *TIGER* and *synthetic* data sets, the number of dimensions is varied with 2, 4, 6, 8, and 10 dimension, and
- (vii) number of objects – the initial size of the Long Beach from the *TIGER* data set is 50,747 objects while the

**TABLE 9.** The parameter settings of the *synthetic* and *TIGER* data sets.

Parameter Settings	Data Sets	
	Synthetic	Long Beach Tiger
Number of dimensions	2, 4, 6, 8, 10	2, 4, 6, 8, 10
Number of Groups of Users	2, 4, 8, <b>16</b> , 32	2, 4, 8, <b>16</b> , 32
Number of Users in Groups	4, 8, <b>15</b> , 20, 25	4, 8, <b>15</b> , 20, 25
Percentage of Overlapping Region	20%, <b>40%</b> , 60%, 80%, 100%	20%, <b>40%</b> , 60%, 80%, 100%
Number of Objects	20,000, <b>50,000</b> , 80,000	<b>50,747</b>
Space	[0, 250]*[0, 250], [0, 500]*[0, 500], [0, 750]*[0, 750], <b>[0, 1000]*[0, 1000]</b>	[0, 250]*[0, 250], [0, 500]*[0, 500], [0, 750]*[0, 750], <b>[0, 1000]*[0, 1000]</b>
Density	-	0.56%, 1.60%, 7%, 15%, 34%

number of objects of the *synthetic* data set is varied with 50K as the minimum and 80K as the maximum number of objects.

Meanwhile, the location of each user is randomly generated within a given space size. These parameter settings are clearly shown in Table 9 with values in bold representing the default values.

## B. THE EXPERIMENTAL RESULTS

This section presents the experimental results of the *RSGU* and *ERSGU* frameworks that are designed with the aim at deriving skyline objects of a group of users in a predetermined region. The performance of both frameworks is measured with regard to processing time with different parameter settings as discussed in Part A and presented in Table 8. These results are compared to the results of *VR* [63], based on the *synthetic* and real data sets. Both the *RSGU* and *ERSGU* frameworks derived skyline objects by analysing the locations of the users, i.e. spatial attributes, as well as the spatial and non-spatial attributes of the set of objects that is within a predetermined region of the group of users. However, in the *ERSGU* framework, the closeness of an object to the desirable facilities or other interesting objects in the region is also considered. Both frameworks work by partitioning the region into smaller units called fragments in such a way that overlapping areas between the current and previous visited regions can be easily determined. The results of computing the skylines of each fragment, known as fragment skylines, are saved and utilised in the skyline computation of subsequent requests. Meanwhile, the *VR* algorithm is performed repeatedly for each group of user's query without exploiting the previous skyline computation results.

*Effect of Number of Groups of Users* – The number of groups of users is one of the factors that has significant effect on the performance of skyline algorithms in processing skyline queries of a group of users. In this section, the experimental results of the proposed solutions, *RSGU* and *ERSGU*,

and the previous algorithm, namely: *VR* [63] are illustrated, for both the *synthetic* and *TIGER* data sets with respect to processing time, by varying the number of groups of users from 2 – 32 as applied in [63]. The parameter settings for the *synthetic* data set are as follows: the number of dimensions is set to 6, the number of users in a group is set to 15 in a fixed space  $[0, 1000] \times [0, 1000]$  with 40% overlapping region, while the number of objects is set to 50K. Meanwhile, the parameter settings for the *TIGER* data set are as follows: the number of dimensions is maintained to 2, the number of users in a group is set to 15 in a fixed space  $[0, 1000] \times [0, 1000]$  with 40% overlapping region, while the number of objects is 50,747.

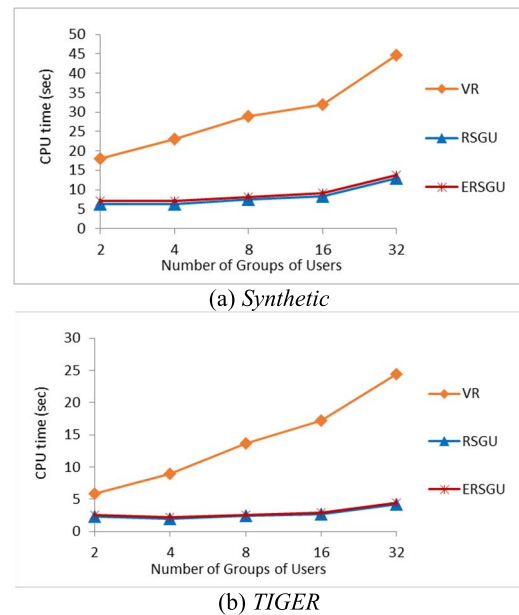
**FIGURE 14.** The results of processing time with varying number of groups of users.

Fig. 14(a) and 14(b) present the processing time achieved by the *RSGU*, *ERSGU*, and *VR* algorithm [63] based on the *synthetic* and *TIGER* data sets, respectively, with the number of groups sets from 2 to 32 groups. The processing time is calculated based on the following formula:

$$\sum_{i=1}^n \text{processing time}_i \quad (3)$$

where  $n$  is the number of groups in a run. For instance, if the number of groups is 4, the processing time is calculated as  $\sum_{i=1}^4 \text{processing time}_i$ . The *VR* algorithm is performed repeatedly for each group of user's query in which the predetermined region of a group is explored even though it has been analysed during the skyline computation of the earlier groups. Meanwhile, for both *RSGU* and *ERSGU*, only the fragment skylines that are related to the identified overlapping area need to be analysed.

Intuitively, when the number of groups increases, the processing time also increases which can be clearly seen through the performance of the *RSGU*, *ERSGU*, and *VR* algorithm. Nonetheless, both *RSGU* and *ERSGU* show a steady

performance for all runs with lesser processing time as compared to the VR algorithm. The processing time of *ERSGU* is slightly higher than *RSGU* since it has an additional evaluation criterion to be analysed, i.e. the closeness of an object to the desirable facilities or other interesting objects in the region. Similar trends as presented in Fig. 14(a) can be seen in Fig. 14(b). On the average, *RSGU* and *ERSGU* gained 70% and 67% improvements for the *synthetic* data set, respectively, and 72% and 69% for the *TIGER* data set, respectively; compared to the VR algorithm.

*Effect of Number of Objects* – In this study, the effect of number of objects on the performance of *RSGU*, *ERSGU*, and VR algorithm [63] is investigated. It is one of the important factors that has high impact on the skyline algorithms in deriving skyline objects. The parameter settings for the *synthetic* data set are as follows: the number of dimensions is fixed to 6, the number of groups is set to 16 groups with each group consisting of 15 users, the overlapping region is fixed to 40%, and the number of objects is varied with the following values: 20K, 50K, and 80K. Since the *TIGER* data set contains only 50,747 objects, hence it is excluded from this experiment.

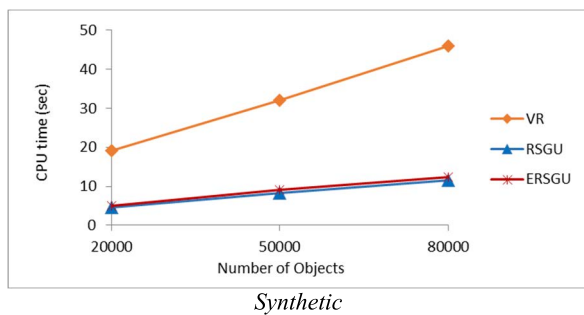


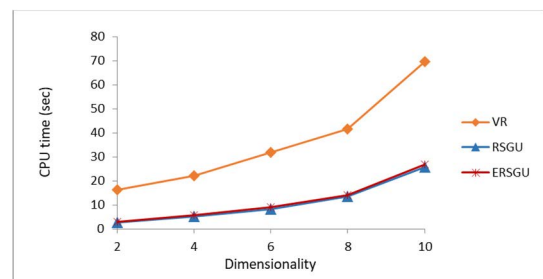
FIGURE 15. The results of processing time with varying number of objects.

Fig. 15 presents the processing time achieved by the *RSGU*, *ERSGU*, and VR algorithm [63] based on the *synthetic* data set, with the number of objects sets as 20K, 50K, and 80K. The processing time is calculated based on formula (3) with the number of groups,  $n = 16$ . The VR algorithm is performed repeatedly for each group of user’s query in which the predetermined region of a group is explored repeatedly even though it has been analysed during the skyline computation of the earlier groups. Meanwhile, for both *RSGU* and *ERSGU*, only the fragment skylines that are related to the identified overlapping area need to be analysed.

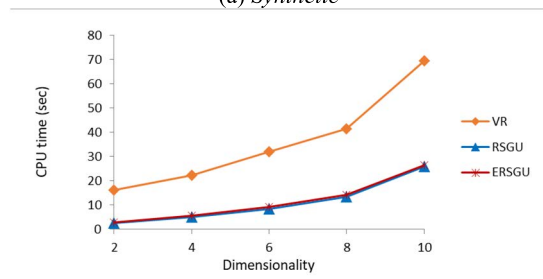
Obviously, when the number of objects increases, the processing time also increases which can be clearly seen through the performance of the *RSGU*, *ERSGU*, and VR algorithm. Nonetheless, both *RSGU* and *ERSGU* show a steady performance for all runs with lesser processing time as compared to the VR algorithm. The processing time of *ERSGU* is slightly higher than *RSGU* since it has an additional evaluation criterion to be analysed, i.e. the closeness of an object to the desirable facilities or other interesting objects in the region. On the average, *RSGU* and *ERSGU* gained 74% and 72%,

improvements, respectively, compared to the VR algorithm for the *synthetic* data set.

*Effect of Data Dimensionality* – Besides the number of groups of users and the number of objects, data dimensionality is also one of the factors that has substantial effect on the performance of skyline algorithms in identifying the skyline objects of a group of users. In this section, the experimental results of the proposed solutions, *RSGU* and *ERSGU*, and the previous algorithm, namely: VR [63] are illustrated, for both the *synthetic* and *TIGER* data sets with respect to processing time, by varying the number of dimensions from 2 – 10 dimensions. The parameter settings for the *synthetic* data set are as follows: the number of objects is fixed to 50K, the number of groups is set to 16 groups with each group consisting of 15 users, and the overlapping region is fixed to 40%. For the *TIGER* data set, the same parameter settings as above are used except that the number objects is maintained to its initial number, i.e. 50,747 objects.



(a) Synthetic



(b) TIGER

FIGURE 16. The results of processing time with varying dimensionality.

Fig. 16(a) and 16(b) present the processing time achieved by the *RSGU*, *ERSGU*, and VR algorithm [63] based on the *synthetic* and *TIGER* data sets, respectively, with number of dimensions varied from 2 – 10 dimensions. The processing time is calculated based on the formula (3) with the number of groups,  $n = 16$ . Obviously, when the number of dimensions increases, the processing time also increases which can be clearly seen through the performance of the *RSGU*, *ERSGU*, and VR algorithm. Nonetheless, both *RSGU* and *ERSGU* show a steady performance for all runs with lesser processing time as compared to the VR algorithm. The processing time of *ERSGU* is slightly higher than *RSGU* since it has an additional evaluation criterion to be analysed, i.e. the closeness of an object to the desirable facilities or other interesting objects in the region. On the average, *RSGU* and *ERSGU* gained 69% and 67% improvements, respectively, for the *synthetic* data

**TABLE 10.** The density rate of the types of objects in the *TIGER* data set.

Types of objects	% of the type of object in the whole population	No. of objects
Hospital	0.56	284
Restaurant	1.60	812
Church	7.00	3,552
School	15.00	7,612
Institution	34.00	17,254
Building	10.84	5,502
Hotel	13.00	6,597
Populated place	18.00	9,134

set and 69% and 67% improvements, respectively, for the *TIGER* data set compared to the *VR* algorithm.

*Effect of Density* – In this study, the effect of density on the performance of *RSGU*, *ERSGU*, and *VR* [63] is investigated. For this experiment, only the *TIGER* data set is considered. The *TIGER* data set consists of eight types of objects, namely: *hospital*, *restaurant*, *church*, *school*, *institution*, *building*, *hotel*, and *populated place*. This is represented in Table 10. The % of the type of object in the whole population and No. of objects reflect the density rate of a particular type of object in the area. For instance, the number of hospitals is 284 which is 0.56% of the whole population, i.e.  $0.56\% \times 50747$ . Meanwhile, *institution* is the densest objects with density rate = 34%. Intuitively, the higher the density rate; the higher is the processing time as more objects need to be analysed.

In this section, the experimental results of the proposed solutions, *RSGU* and *ERSGU*, and the previous algorithm, namely: *VR* [63] are illustrated, for the *TIGER* data set with respect to processing time, with various density rates as follows: 0.56% (*hospital*), 1.60% (*restaurant*), 7.00% (*church*), 15.00% (*school*), and 34.00% (*institution*). These density rates reflect the least dense to the densest that are available in the data set. The parameter settings used are as follows: the number of objects is 50,747, the number of groups is set to 16 groups with each group consisting of 15 users, the number of dimensions and the overlapping region is fixed to 2 and 40%, respectively.

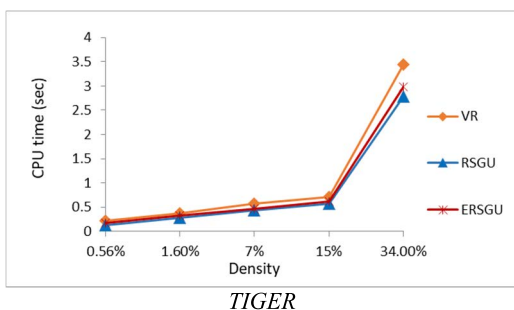
**FIGURE 17.** The results of processing time with varying density rate.

Fig. 17 shows the performance of *RSGU*, *ERSGU*, and *VR* algorithm [63] with regard to processing time. The performance of *RSGU*, *ERSGU*, and *VR* algorithm shows similar trends in which it starts to show a drastic increment when the

density rate is 15.00% until it reaches to 34.00%. This is due to the fact that the number of *institutions* (34.00%) is slightly more than twice the number of *schools* (15.00%). Despite that, both *RSGU* and *ERSGU* show a better performance for all runs with lesser processing time as compared to the *VR* algorithm. The percentage of improvement gained by *RSGU* and *ERSGU* is 21% and 14%, respectively. The processing time of *ERSGU* is slightly higher than *RSGU* since it has an additional evaluation criterion to be analysed, i.e. the closeness of an object to the desirable facilities or other interesting objects in the region. For instance, if the object of interest is *restaurant*, then the desirable facilities or other interesting objects are *hospital*, *church*, *school*, and *institution*.

*Effect of Space Size* – In this study, the effect of space size on the performance of *RSGU*, *ERSGU*, and *VR* algorithm is also investigated. In this section, the experimental results of the proposed solutions, *RSGU* and *ERSGU*, and the previous algorithm, namely: *VR* [63] are illustrated for both the *synthetic* and *TIGER* data sets with respect to processing time, by varying the space size as follows:  $[0, 250] \times [0, 250]$ ,  $[0, 500] \times [0, 500]$ ,  $[0, 750] \times [0, 750]$ , and  $[0, 1000] \times [0, 1000]$ . The parameter settings used for the *TIGER* data set are as follows: the number of objects is fixed to 50,747 objects, the number of groups is set to 16 groups with each group consisting of 15 users, and the overlapping region is set to 40%. For the *synthetic* data set, the same parameter settings as above are used except that the number objects is set to 50K objects.

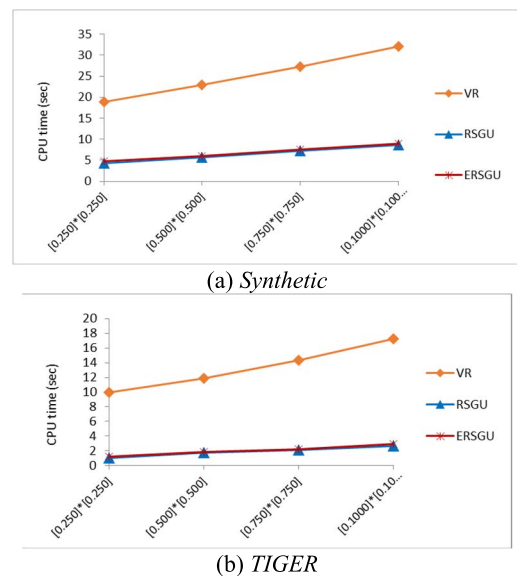
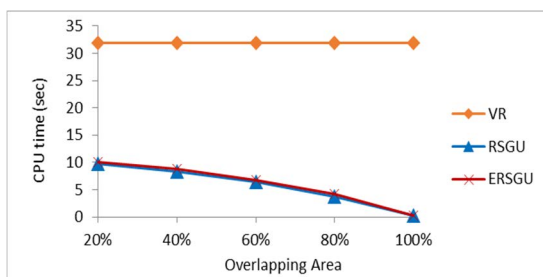
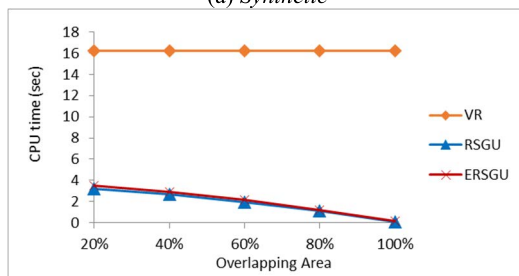
**FIGURE 18.** The results of processing time with varying space size.

Fig. 18(a) and 18(b) present the processing time achieved by the *RSGU*, *ERSGU*, and *VR* algorithm [63] based on the *synthetic* and *TIGER* data sets, respectively, with different space sizes. The processing time is calculated based on the formula (3) with the number of groups,  $n = 16$ . Obviously, the bigger the space size, the more objects it covered; hence the higher is the processing time. The *VR* algorithm is performed repeatedly for each group of user's query in



which the predetermined region of a group is explored repeatedly even though it has been analysed during the skyline computation of the earlier groups of users. Meanwhile, for both *RSGU* and *ERSGU*, only the fragment skylines that are related to the identified overlapping area need to be analysed. From the figure, both *RSGU* and *ERSGU* show a steady performance with a slight increment in each iteration. Based on this analysis, *RSGU* and *ERSGU* gained 73% and 74% improvements, respectively, for the *synthetic* data set, and 82% and 83% improvements, respectively, for the *TIGER* data set, compared to the *VR* algorithm.

*Effect of Overlapping Region* – Another factor that has a significant effect on the performance of skyline algorithms for a group of users is the overlapping region covered between the groups of users. In this section, the experimental results of the proposed solutions, *RSGU* and *ERSGU*, and the previous algorithm, namely: *VR* [63] are illustrated, for both the *synthetic* and *TIGER* data sets with respect to processing time, by varying the percentage of overlapping region from 20% – 100%. The parameter settings for the *synthetic* data set are as follows: the number of objects is fixed to 50K, the number of groups is set to 16 groups with each group consisting of 15 users, and the number of dimensions is fixed to 6. For the *TIGER* data set, the same parameter settings as above are used except that the number objects is maintained to its initial number, i.e. 50,747 objects in  $[0,1000]*[0,1000]$ , and each object is with 2 dimensions.

(a) *synthetic*(b) *TIGER*

**FIGURE 19.** The results of processing time with varying percentage of overlapping area.

Fig. 19(a) and 19(b) present the processing time achieved by the *RSGU*, *ERSGU*, and *VR* [63], based on the *synthetic* and *TIGER* data sets, respectively, with varying percentage of overlapping region. The processing time is calculated based on the formula (3) with the number of groups,  $n = 16$ . From these figures, both the *RSGU* and *ERSGU* show a steady

performance which reflects that when the percentage of overlapping region increases, the processing time decreases. This is due to the fact that both the *RSGU* and *ERSGU* exploit the skyline computation results of the previous groups of users that are associated to the identified overlapping region. Hence, the higher the percentage of overlapping region means the higher the percentage of the area that has been explored in the earlier skyline computations of the groups of users. In both solutions, *RSGU* and *ERSGU*, rescanning of the overlapping region and recomputation of skyline objects within the overlapping region are avoided. Interestingly, when the percentage of overlapping region is 100%, which implies that the region covered by the current group of users is the exact same region that has been explored in the skyline computations of the previous groups of users; shows the processing time taken is almost 0 because of the skyline results are the same. However, the *VR* algorithm shows a steady performance for all runs which clearly indicates that the overlapping region has no significant effect on the performance of the *VR* algorithm. This is true since the *VR* algorithm is performed repeatedly for each group of user's query in which the predetermined region of a group is explored repeatedly even though it has been analysed during the skyline computation of the earlier groups. Based on this analysis, *RSGU* and *ERSGU* gained 82% and 81% improvements, respectively, for the *synthetic* data set, and 87% and 86% improvements, respectively, for the *TIGER* data set, compared to the *VR* algorithm.

*Effect of Number of Users in a Group* – Another factor that has a major impact on the performance of skyline algorithms in processing the skyline queries of a group of users is the number of users in a group. In this section, the experimental results of the proposed solutions, *RSGU* and *ERSGU*, and the previous algorithm, namely: *VR* [63] are illustrated, for both the *synthetic* and *TIGER* data sets with respect to processing time by varying the number of users in a group from 4 – 25 as applied in the previous study [63]. The parameter settings for the *synthetic* data set are as follows: the number of objects is fixed to 50K, the number of groups is set to 16 groups in a fixed space  $[0, 1000]*[0, 1000]$  with 40% overlapping region, while the number of dimensions is fixed to 6. For the *TIGER* data set, the same parameter settings as above are used except that the number objects is maintained to its initial number, i.e. 50,747 objects, with the number of dimensions fixed to 2.

Fig. 20(a) and 20(b) present the processing time achieved by *RSGU*, *ERSGU*, and *VR* [63] based on the *synthetic* and *TIGER* data sets, respectively, with the number of users in a group sets to 4, 8, 15, 20, and 25. The processing time is calculated based on the formula (3) with the number of groups,  $n = 16$ . From these figures, both *RSGU* and *ERSGU* show similar performance for all runs with lesser processing time as compared to the *VR* algorithm. The processing time of *ERSGU* is slightly higher than *RSGU* since it has an additional evaluation criterion to be analysed, i.e. the closeness of an object to the desirable facilities or other interesting objects in the region. On the average, *RSGU* and *ERSGU* gained 72%

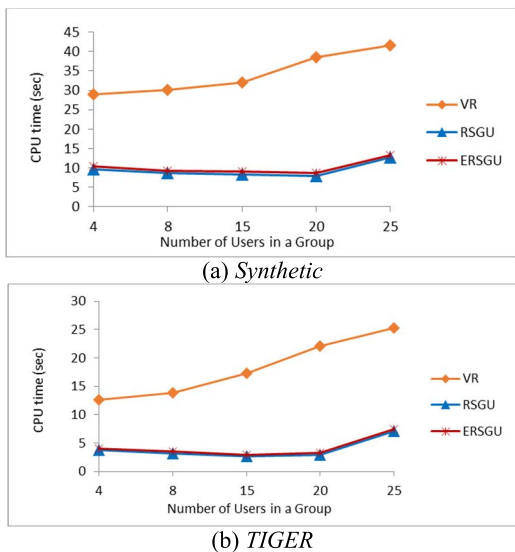


FIGURE 20. The results of processing time with varying number of users in a group.

and 70% improvements, respectively, for the *synthetic* data set and 76% and 78%, respectively, for the *TIGER* data set, compared to the VR algorithm.

*Number of Skyline Objects* – Since the number of skyline objects is a significant factor in validating the correctness of the skyline algorithms in processing skyline queries, thus investigating the performance of *RSGU*, *ERSGU*, and the previous algorithm, *VR* [63] with regard to the skyline objects derived by these solutions is inevitable. In this section, the experimental results of the proposed solutions, *RSGU* and *ERSGU*, and the previous algorithm, namely: *VR* [63], for both the *synthetic* and *TIGER* data sets with respect to the number of skyline objects derived by these solutions are illustrated. The parameter settings for the *synthetic* data set are as follows: the number of objects is fixed to 50K, the number of groups is set to 16 groups in a fixed space  $[0, 1000] \times [0, 1000]$  with 40% overlapping region, the number of dimensions is fixed to 6, while the number of users in a group is varied from 4 – 25. For the *TIGER* data set, the same parameter settings as above are used except that the number objects is maintained to its initial number, i.e. 50,747 objects, with the number of dimensions fixed to 2.

Fig. 21(a) and 21(b) present the number of skyline objects derived by the *RSGU*, *ERSGU*, and the *VR* algorithm [63], based on the *synthetic* and *TIGER* data sets, respectively. From these figures, it is obvious that the number of skyline objects derived by *RSGU* and *VR* is the same for all runs which verified the correctness of *RSGU*. This is because in deriving the skyline objects, both *RSGU* and *VR* utilised the same evaluation criteria, namely: spatial and non-spatial attributes of the objects. However, as expected the number of skyline objects produced by *ERSGU* is different from the number of skyline objects obtained by both the *RSGU* and *VR* algorithm since *ERSGU* has an additional evaluation criterion to be analysed, i.e. the closeness of an object to the desirable facilities or other interesting objects in the region.

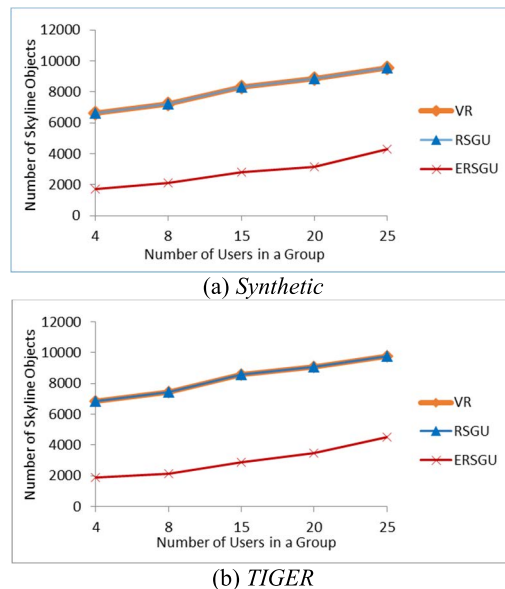


FIGURE 21. The results of number of skyline objects with varying number of users in a group.

### C. TIME COMPLEXITY ANALYSIS

This section presents the time complexity analysis of the proposed frameworks, *Region-based Skyline for a Group of Users (RSGU)* and *Extended Region-based Skyline for a Group of Users (ERSGU)*. Since *ERSGU* utilises a different set of evaluation criteria in deriving the final skyline objects, hence we only report its time complexity without comparing it to the baseline method. On the other hand, the time complexity of *RSGU* is compared to a baseline method (*BM*) which utilises the conventional skyline algorithm in deriving the skyline objects for a group of users. In this method, each group of users is treated separately. This approach is assumed by many methods including the *VR* algorithm [63].

The time complexity analysis is based on the following: Given a data set  $D = \langle R, U, O \rangle$ , where  $U = \{u_1, u_2, \dots, u_n\}$  is a list of  $n$  users and  $O = \{o_1, o_2, \dots, o_m\}$  is a list of  $m$  objects. Let  $G_p = \{u_1, u_2, \dots, u_p\}$  be a group of  $p$  users where  $G_p \subset U$  in region  $R_p$ . Assume that  $O_r = \{o_1, o_2, \dots, o_r\}$  is the list of objects that is within the region  $R_p$ . Table 11 and Table 12 present the time complexity at each step of the *RSGU* and *ERSGU*, respectively.

In order to compare the time complexity of *RSGU* against the baseline method (*BM*), assume the following:

Given a group of users,  $G_p = \{u_1, u_2, \dots, u_p\}$ , where  $G_p \subset U$ , and the candidate skylines of  $G_p$  in region  $R_p$  denoted as  $CS_{G_p}$  with  $c$  cardinality. Find the skylines of a group of users  $G_q = \{u_1, u_2, \dots, u_q\}$  in region  $R_q$ , i.e.  $CS_{G_q}$ , where  $G_q \subset U$ ,  $G_q \neq G_p$ , and  $R_q \cap R_p \neq \emptyset$ . Assume that  $O_r = \{o_1, o_2, \dots, o_r\}$  is the list of objects that is within the region  $R_p$ ;  $O_s = \{o_1, o_2, \dots, o_s\}$  is the list of objects that is within the region  $R_q$ ,  $u$  is the number of candidate skylines of  $CS_{G_p}$  that are in the overlapping area  $O_R$ , and  $v$  is the number of objects that are in the non-overlapping area  $\neg O_R$ . Obviously,  $u \leq r$ ,  $u \leq s$ ,  $u \leq c$ , and  $u + v < s$ .

TABLE 11. The time complexity of RSGU.

Steps	Time Complexity	Remarks
(1) Identify the centroid	$O(1)$	The centroid is identified using Equation (1).
(2) Construct a search region	$O(r) + O(p) + O(1)$	<ol style="list-style-type: none"> <li>1. Identify the closest object to the centroid: <math>r</math> iterations, where <math>r</math> is the number of objects within the region <math>R_p</math>.</li> <li>2. Construct a search region of each user in the <math>G_p</math>: <math>p</math> iterations, where <math>p</math> is the number of users in <math>G_p</math>.</li> <li>3. Construct a search region for <math>G_p</math>: 1 iteration, by performing union over the search region of each user.</li> </ol>
(3) Identify the overlapping region	$O(1) + O(\log M n)$	<ol style="list-style-type: none"> <li>1. Identify the overlapping region based on the polygon's vertices: 1 iteration.</li> <li>2. Traverse the <math>R</math>-tree to identify the fragments that are within the overlapping region: <math>O(\log M n)</math>, where <math>M</math> is the maximum number of entries and <math>n</math> is the minimum number of entries in a node.</li> </ol>
(4) Construct the fragments of a search region	$O(2p \log 2p) + O(2p \log 2p) + O(k) + O(w^2)$	<ol style="list-style-type: none"> <li>1. Sort the <math>x</math>-coordinates of each search region: <math>O(2p \log 2p)</math> with 2 <math>x</math>-coordinates for each <math>p</math> user.</li> <li>2. Sort the <math>y</math>-coordinates of each search region: <math>O(2p \log 2p)</math> with 2 <math>y</math>-coordinates for each <math>p</math> user.</li> <li>3. Construct the <math>k</math> fragments: <math>k</math> iterations.</li> <li>4. Derive the non-spatial candidate skylines for <math>k</math> fragments: <math>\sum_{i=1}^k w_i(w_i - 1)/2</math> iterations, where <math>w_i</math> is the number of objects of the <math>i</math>-th fragment; this is simplified as <math>w^2</math> iterations.</li> </ol>
(5) Construct the $R$ -tree of the fragments	$O(\log M n)$	The $R$ -tree is constructed based on the algorithms proposed by [30]: $O(\log M n)$ , where $M$ is the maximum number of entries and $n$ is the minimum number of entries in a node.
(6) Derive the non-spatial skylines	$O(c^2)$	Apply the non-spatial dominance over the candidate skylines of the $k$ fragments: $c(c - 1)/2$ iterations, where $c$ is the number of candidate skylines of the $k$ fragments.

The time complexity of both methods is analysed based on the following steps: (i) derive the final skylines of  $G_p$  and (ii) derive the final skylines of  $G_q$ . Here, we assumed that the

TABLE 11. (Continued.) The time complexity of RSGU.

(7) Derive the spatial skylines	$O(pr) + O(r) + O(r \log r) + O(r^2)$	<ol style="list-style-type: none"> <li>1. Calculate the distance between each user of <math>G_p</math> and each object of <math>O_r</math>: <math>p \times r</math> iterations.</li> <li>2. Calculate the <math>Sum\ Distance - o_i</math> of each object of <math>O_r</math>: <math>r</math> iterations.</li> <li>3. Sort the <math>Sum\ Distance - o_i</math> of the <math>O_r</math> objects: <math>O(r \log r)</math></li> <li>4. Apply the spatial dominance over the objects of <math>O_r</math>: <math>r(r - 1)/2</math> iterations.</li> </ol>
(8) Derive the final skylines	$O(1)$	Union between the non-spatial skylines and spatial skylines.

TABLE 12. The time complexity of ERSQU.

Steps	Time Complexity	Remarks
Steps (1) till (7) are the same steps as RSGU.		
(8) Derive the candidate skylines	$O(1)$	Union between the non-spatial skylines and spatial skylines.
(9) Derive the final skylines	$O(yh) + O(yh) + O(y^2)$	<ol style="list-style-type: none"> <li>1. Calculate the distance between each candidate skyline <math>o_j</math> and a set of <math>p</math> distinct type of interesting objects with <math>h</math> total number of objects: <math>y \times h</math> iterations, where <math>y</math> is the number of candidate skylines.</li> <li>2. Identify the closest interesting object of each type to the object <math>o_j</math>: <math>\sum_{i=1}^p y \times h_i</math> iterations, where <math>h_i</math> is the number of objects of type <math>p_i</math>.</li> <li>3. Perform the conventional skyline algorithm over the candidate skylines with the closest distance of each type as the evaluation criteria: <math>y(y - 1)/2</math> iterations.</li> </ol>

region  $R_p$  is an unexplored region while  $R_q \cap R_p \neq \emptyset$  implies that some parts of the region  $R_q$  have been analysed in step (i). To simplify the comparisons, only the steps followed by  $BM$  and  $RSGU$  with different time complexities are analysed.

- (i) To derive the final skylines of  $G_p$  using the baseline method, similar steps to the steps 1, 2, 6, 7, and 8 are performed. The objects analysed in Step 6 are those objects that are within the region  $R_p$ ; hence the number of objects analysed is  $r$  with  $r(r - 1)/2$  pairwise comparisons [52]. The time complexity to derive the final skylines of  $G_p$  with the baseline method,  $T(BM_{G_p})$ , is given below:

$$T(BM_{G_p}) = T(\text{Step 6}) \approx O(r^2)$$

Meanwhile, utilising the  $RSGU$ , steps 1, 2, 4, 5, 6, 7, and 8 are performed. The time complexity to derive the final skylines of  $G_p$  with  $RSGU$ ,  $T(RSGU_{G_p})$ , is as

given below:

$$\begin{aligned} T(RSGU_{G_p}) &= T(\text{Step 4}) + T(\text{Step 5}) + T(\text{Step 6}) \\ &= [O(2p \log 2p) + O(2p \log 2p) \\ &\quad + O(k) + O(w^2)] + O(\log Mn) \\ &\quad + O(c^2) \approx O(w^2) + O(c^2) \end{aligned}$$

Since  $w + c < r$ , thus  $T(RSGU_{G_p}) < T(BM_{G_p})$ .

- (ii) To derive the final skylines of  $G_q$  using the baseline method, similar steps to the steps 1, 2, 6, 7, and 8 are performed. The objects analysed in Step 6 are those objects that are within the region  $R_q$ ; hence the number of objects analysed is  $s$  with  $s(s-1)/2$  pairwise comparisons [52]. The time complexity to derive the final skylines of  $G_q$  utilising the baseline method,  $T(BM_{G_q})$  is as follows:

$$T(BM_{G_q}) = T(\text{Step 6}) \approx O(s^2)$$

Meanwhile, utilising the *RSGU*, steps 1, 2, 3, 4, 5, 6, 7, and 8 are performed.

$$\begin{aligned} T(RSGU_{G_q}) &= T(\text{Step 3}) + T(\text{Step 4}) + T(\text{Step 5}) \\ &\quad + T(\text{Step 6}) = O(\log Mn) \\ &\quad + [O(2p \log 2p) + O(2p \log 2p) + O(k) + O(v^2)] \\ &\quad + O(\log Mn) + O(u^2) \approx O(v^2) + O(u^2) \end{aligned}$$

Step 4 is performed over the non-overlapping area involving  $v$  objects, while Step 6 is performed over the overlapping area in which  $u$  candidate skylines of  $CS_{G_p}$  are analysed.

Since  $v + u < s$ , thus  $T(RSGU_{G_q}) < T(BM_{G_q})$ .

- (iii) The total time complexities for both methods are as follows:

$$\begin{aligned} T(BM) &= T(BM_{G_p}) + T(BM_{G_q}) \\ &\approx O(r^2) + O(s^2) \\ T(RSGU) &= T(RSGU_{G_p}) + T(RSGU_{G_q}) \\ &\approx O(w^2) + O(c^2) + (v^2) + O(u^2) \end{aligned}$$

It is obvious that  $T(RSGU) < T(BM)$  since  $w + c < r$  and  $v + u < s$ .

## VII. CONCLUSION

In this paper, we proposed the *Region-based Skyline for a Group of Users (RSGU)* and *Extended Region-based Skyline for a Group of Users (ERSGU)* frameworks that are designed to derive skyline objects which are point of interests (*PoIs*) to be recommended to a group of users. The skyline objects are derived by analysing both the locations of the users, i.e. spatial attributes, as well as the spatial and non-spatial attributes of objects that are within a predetermined region of the group of users. Two main aims have been set that are:

- (i) to avoid the process of rescanning the set of objects within a predetermined region that is known to have been previously visited by a group of users and (ii) to avoid the recomputation of skylines of a set of objects within a predetermined region that has been analysed in earlier computations of previously visited group of users. Meanwhile, *ERSGU* framework considers the closeness of the objects to other interesting objects in its solution as its additional feature. Several experiments have been conducted and the results show that both the *RSGU* and *ERSGU* outperform the work by [63] with respect to CPU time. There are several further enhancements that can be made based on the findings presented in the paper. These include (i) proposing an approach to continuously derive skyline objects by considering the movement of the users, (ii) incorporating the *Group-based skyline* in finding the optimal combinations of skyline objects [86] for a group of users, and (iii) embedding the proposed frameworks into a Travelling Recommender System similar to the work in [87].

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