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Vertex Metric-Based Dimension of Generalized Perimantanes Diamondoid Structure

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ABSTRACT Due to its superlative physical qualities and its beauty, the diamond is a renowned structure. While the green-colored perimantanes diamondoid is one of a higher diamond structure. Motivated by the structure's applications and usage, we look into the metric-based parameters of this structure. In this draft, we have discussed metric dimension and their generalizations for the generalized perimantanes diamondoid structure and proved that each parameter depends on the copies of original or base perimantanes diamondoid structure and changes with the parameter *n* or its number of copies.

INDEX TERMS Vertex metric dimension, generalized perimantanes diamondoids, diamond structure, resolving set.

I. INTRODUCTION

Due to its superlative physical qualities and its beauty, the diamond is a renowned structure. Polishing, drilling, cutting, and heatsink in electronics are numerous practical and industrial applications of diamond. Its hardness, exceptional thermal conductivity determines by its rigor of composition. A single molecule with macroscopic size makes it into a diamond crystal. The model depicting fundamental atomic groupings undergoes alterations while going from molecules to materials, both in terms of idea and actual manifestation, as well as insignificant computational processing [1].

There are four types of higher diamondoids and each have been assigned with four different colors and name. The assigned colors are yellow, red, blue, and green, while the names are, nonbranched rodlike zigzag catamantanes associated with yellow. The regularly branched catamantanes are linked with blue diamondoids, chiral diamondoids are red-colored and the green-colored are perimantanes diamondoids. All these series have been isolated and found from petroleum [2].

A chemical/molecular graph is a hydrogen-depleted molecular structure in which the edges represent bonds and the vertices represent atoms in the underlying organic chemical compounds. The chemical graph theory is the study of these chemical graphs [3], [4]. There are enough data available on this assumption and transformation from a chemical

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structure to a graph (vertex-edge-based structure). More detail can be found in the recent literature such as [5]–[10].

The notion of resolving set was proposed by the researcher in [11]. It is the first study to look at the notion of finding a graph's metric dimension using the definition of a resolving set. The least cardinality of a resolving set is the metric dimension. The impetus for inventing the notion of finding the set came from LORAN and sonar stations. After that, several academics took this concept and labeled it in a variety of ways. The idea of a resolving set is dubbed as a metric dimension in [12]'s study. While the researchers in [13], [14] renamed the same notion with metric foundation or resolving set in a purely theoretical fashion. A more advanced definition of a resolving set was developed in the last decade. Researchers of [15] in which the idea of one node faultiness from the resolving set is explained. The notion is referred to as a fault-tolerant resolving set, and it is a generalized form of the resolving set.

Many concepts and implementations sprang from the generalized approach of resolving set. In electronics, [16], a recent innovation reveals the implementation of locating set (and its extensions). A method for studying diverse polyphenyl structures, especially for the polymer industry [17]. In the broader view, this idea is used in combinatorial optimization [18], some complex games or robotic roving [19], image processing [11], pharmaceutical chemistry [20].

The job of determining a graph's resolving set is a nondeterministic polynomial-time hard problem (NP), with an

unknown algorithmic difficulty [20]–[23]. Metric dimension or resolving set has a large literature because of its many variations and applications in various disciplines. Only the most current and broad results will be discussed here. The Internet graph and its fault-tolerant structure are discussed in [24]. Quartz structure is studied with the concept of fault-tolerant locating number in [25]. Computer-related interconnection networks are studied in [26]. On the topic of fault-tolerant locating set, [27] discussed convex polytopes and found their exact fault-tolerant locating numbers. There is extensive literature available on this topic, we refer to see some recent research work on this definition [28]–[32]. A general graph of a kayak paddle and some other cycle-related graphs are dealt with in [33]. A cellulose network is studied in [34], they computed some upper bounds for the structure. A coronoid structure is reestablished in the form of a metric of a graph in [35]. A hydrocarbon-based class of structure was studied in [36] with the concept of locating numbers and also determined some other variants. The generalized class of the Harary family is studied with the definition of locating set [37]. Generalized Peterson graphs and multi-graphs are discussed in [3], with the concept of metric basis. The researchers in [38] studied this definition on the Cayley graph and find out the locating number for such a generalized class. Moreover, some recent studies and literature are available at [39]–[41]. On the topic of partition dimension we refer to the extensive study and some bounds on this topic, the literature is [42]–[45].

Given below are some basic concepts elaborated for further use in our main results.

Definition 1 [46]: ''Let *G* be a connected graph with vertex set $V(G)$, the distance between two vertices *v*₁, *v*₂ ∈ *V*(*G*) is the length of a shortest *v*₁ − *v*₂ path between them. It is denoted by $d(v_1, v_2)$."

Definition 2: Let *B* be an ordered simple subset from the vertex set of a graph, say $B(G) = \{v_1, v_2, \ldots, v_i\}$, consider $v \in V(G)$. The location (position, representation) *r* (*v*|*B*) of *v* according to the subset *B* is the *i*−tuple distances $(d(v, v_1), d(v, v_2), \ldots, d(v, v_i))$. Taking any two vertices of a graph *G*, if these vertices have different locations $r(v|B)$ according to chosen subset B , then B is considered as a locating set for the graph. The minimum possible members (or the cardinality) of locating set is the metric dimension of the graph and we will defined this with the symbol *dim* (*G*).

Definition 3: A pertinent chosen subset *B* which is locating set becomes fault-tolerant locating set B_f , if it fulfills the condition $B_f \vee v$ for each vertex $v \in \mathit{ls}_f$ and it remains locating set. The minimum possible members (or the cardinality) of fault-tolerant locating set is the fault-tolerant metric dimension of the graph and we will defined this with the symbol $dim_f(G)$.

Definition 4: Let $B_p = \{B_{p_1}, B_{p_2}, \ldots, B_{p_\ell}, \}$ be the partition of a connected graph. Now for a vertex $j \in V(G)$ the partition dimension with respect to B_p is $r(j|B_p)$ = $(d(j, B_{p_1}), d(j, B_{p_2}), \ldots, d(j, B_{p_\ell}))$, where $d(j, B_w)$ = $min{d(j, y) : y \in B_w}$ for $1 \leq w \leq \ell$. The distinct codes of the two vertices $i, j \in V(G)$ with respect to B_p , that is $r(i|B_p) \neq r(j|B_p)$, such a partition B_p is known as distinguishing partition of *G* and denoted by pd(*G*) [47].

In this draft, we have discussed metric dimension and their generalizations for the generalized perimantanes diamondoid structure and proved that each parameter depends on the copies of original or base perimantanes diamondoid structure and changes with the parameter *n* or its number of copies. The next section will present some main results, conclusions are drawn and at the end, references are given for more and deep insight into this topic and structure.

II. GENERALIZED PERIMANTANES DIAMONDOID STRUCTURE AND MAIN RESULTS

The structure shown in Figure [1,](#page-2-0) is a green-colored perimantanes diamondoid and one of a higher diamond structure. Its topological version is found in [1], [2] and motivated by the structures applications and usage, we look into the metric-based parameters of this structure. The perimantanes diamondoid structure has total $|V(D_n)| = 22n + 3$, number of vertices and total edges are $|E(D_n)| = 38n + 2$. The labelling of vertices and edges is described in the Figure [1](#page-2-0) and is utilized in the major results. Furthermore, vertex and edge are stated given below.

$$
V(D_n)
$$

= { a_i^j : $i = 1, 2, ..., 19, j = 1, 2, ..., n$ }
 $\cup \{b_i : i = 1, 2, ..., 3 (n + 1)\},$
 $E(D_n)$
= { $a_i^j a_{i+1}^j$: $i = 1, 2, 4, 5, 7, 9, 10, 12, 14, 15, 17, 18,$
 $j = 1, 2, ..., n$ } $\cup \{b_i b_{i+1} : i = 1, 4, 7, ..., 3n + 1,$
 $i = 2, 5, ..., 3n + 1$ } $\cup \{b_i a_2^j : i = 1, 4, 7, ..., 3n - 2,$
 $j = 1, 2, ..., n$ } $\cup \{b_i a_3^j, b_i a_4^j : i = 2, 5, 8, ...,$
 $3n - 1, j = 1, 2, ..., n$ } $\cup \{b_i a_5^j : i = 3, 6, 9,$
 $..., 3n, j = 1, 2, ..., n$ } $\cup \{b_i a_{15}^j : i = 4, 7, 10, ...,$
 $3n + 1, j = 1, 2, ..., n$ } $\cup \{b_i a_{16}^j, b_i a_{17}^j : i = 5, 8,$
 $11, ..., 3n + 2, j = 1, 2, ..., n$ } $\cup \{b_i a_{18}^j : i = 6, 9,$
 $12, ..., 3n + 3, j = 1, 2, ..., n$.

Presented below are the main results of this novel structure.

Lemma 1: Let D_n is a structure of perimantanes diamondoid with $n \geq 1$, and *B* is the vertex resolving set of D_n for $n = 1$. Then one of the possible resolving set is stated by

$$
B = \{b_1, b_4, a_7^1\}.
$$

Proof: To demonstrate that the resolving set for the structure of perimantanes diamondoid, shown in the Figure [1](#page-2-0) and labeled as D_1 , for its particular value of $n = 1$. Let *B* be a resolving set and stated by, $B = \{b_1, b_4, a_7^1\}$. Moreover its resolving set shown in the Figure [2.](#page-2-1) The unique locations of

FIGURE 1. Vertex-edge sets of perimantanes diamondoid structure D_n.

the full vertex set of D_1 , with regard to the elements of *B* are as follows.

$$
r\left(a_i^j|B\right) = \begin{cases} (3-i, 5-i, i), & \text{if } i = 1, 2; \\ (2, 4, i), & \text{if } i = 3; \\ (2, i, i + 1), & \text{if } i = 4; \\ (i-2, i, i+1), & \text{if } i = 5, 6; \\ (10-i, 10-i, i-7), & \text{if } i = 7, 8; \\ (i-7, i-7, i-7), & \text{if } i = 9, 10, 11; \\ (i-8, i-8, i-5), & \text{if } i = 12, 13; \\ (18-i, 16-i, i-13), & \text{if } i = 14, 15; \\ (4, 2, i-13), & \text{if } i = 16; \\ (i-13, i-15, i-12), & \text{if } i = 17, 18, 19. \end{cases}
$$
(1)

$$
r(b_i|B) = \begin{cases} (i-1, 3+i, i+2), & \text{if } i = 1, 2, 3; \\ (i, i-4, i-1), & \text{if } i = 4, 5, 6. \end{cases}
$$
 (2)

We also observe that each vertex has a different representation and satisfy the concept of resolving set based on the arguments given above in the form of identifications of the whole vertex set of D_1 , which leads to the conclusion that defined *B* is one of a possible resolving set with $|B| = 3$.

Lemma 2: Let D_n is a structure of perimantanes diamondoid with $n = 1$. Then

 $dim(D_1) = 3.$

Proof: We employ the double inequality method and refer to Lemma [1](#page-1-0) to establish that the metric dimension of D_1 is 3. This plainly demonstrates that the resolving set has cardinality 3, which is described by $B = \{b_1, b_4, a_7^1\}$.

Now on contrary we have $dim(D_1) = 2$. from the assertion $dim(D_1) \geq 3$. Consider the resolving set *B*^{\prime} with cardinality 2. This assumption is discussed in the following cases.

Case 1: Let a chosen subset *B*['] having two distinct elements, say $B' = \{b_1, b_2\}$. The contradiction will be arise due the vertices which have two distance to any of the chosen element of *B'*. Mathematically, it can be written as $r(a_i^1|B') =$ $r\left(a_j^1 | B'\right) = d\left(a_i^1, b_1\right) = 2.$ *Case* 2: Let a chosen subset *B*^{\prime} having two distinct elements, say $B' = \{b_1, b_3\}$. The contradiction will be arise due

FIGURE 2. Resolving set of Diamond structure D₁.

the vertices which have two distance to any of the chosen element of *B'*. Mathematically, it can be written as $r(a_i^1|B') =$ $r\left(a_j^1 | B'\right) = d\left(a_i^1, b_1\right) = 2.$

Case 3: Let a chosen subset *B*^{\prime} having two distinct elements, say $B' = \{b_2, b_3\}$. The contradiction will be arise due the vertices which have two distance to any of the chosen element of *B'*. Mathematically, it can be written as $r(a_i^1|B') =$ $r\left(a_j^1 | B'\right) = d\left(a_i^1, b_1\right) = 2.$

Case 4: Let a chosen subset *B*^{\prime} having two distinct elements, say $B' = \{a_i^1, a_j^j\}$ $\binom{1}{1}$, with distinct *i*, *j*. The contradiction will be arise due the vertices which have two distance to any of the chosen element of B' . Mathematically, it can be written as $r(a_{\alpha}^1 | B') = r(a_{\beta}^1 | B') = d(a_{\alpha}^1, a_i^1) = 2.$ $\beta^{|\mathbf{D}|}$ $\int -a \, (u_{\alpha}, u_i)$

Similarly, There is not just one option among the available combinations $|V|C_2 = \frac{|V(D_1)|!}{2!(|V(D_1)|-2)!} = \frac{(25)!}{2 \times (23)!} = 300$ of the entire vertex set of D_1 . This indicate that two metric dimension of D_1 is not possible. Hence; $dim(D_1) \geq 3$. Hence,

$$
dim(D_1) = 3. \tag{3}
$$

$$
\Box
$$

Theorem 1: Let D_n is a structure of perimantanes diamondoid with $n > 2$. Then

$$
dim(D_n) = n + 2. \tag{4}
$$

Proof: To show that $dim(D_n) = n+2$, we will applying the induction method on n showing the number of copies of base perimantanes diamondoid graph. The base case for $n = 1$ $n = 1$ proved in the Lemmas 1 and [2,](#page-2-2) now assume that the assertion is true for $n = \alpha$.

$$
dim(D_{\alpha}) = \alpha + 2. \tag{5}
$$

We will show that it is true for $n = \alpha + 1$. Suppose

$$
dim (D_{\alpha+1}) = dim (D_{\alpha}) + dim (D_1) - 2.
$$
 (6)

Using Equations [3](#page-2-3) and [5](#page-2-4) in Equation [6,](#page-2-5) we have

$$
dim (D_{\alpha+1}) = \alpha + 2 + 3 - 2,
$$

= \alpha + 3. (7)

FIGURE 3. Resolving set of Diamond structure D_n.

As a result, the conclusion holds for all positive integers $n \geq 1$.

Moreover, the generalize resolving set for the generalized structure of perimantanes diamondoid, shown in the Figure 3 and stated by in the set form B $\{b_1, b_{3n+1}, a_7^1, a_7^2, a_7^3, \ldots, a_7^n\}.$ This concludes.

Lemma 3: Let D_n is a structure of perimantanes diamondoid with $n \geq 1$, and B_f is the vertex fault-tolerant resolving set of D_n for $n = 1$. Then one of the possible vertex fault-tolerant resolving set is stated by

$$
B_f = \{b_1, b_3, b_4, b_6, a_7^1, a_{13}^1\}.
$$

Proof: To demonstrate that the fault-tolerant resolving set for the structure of perimantanes diamondoid, shown in the Figure [1](#page-2-0) and labeled as D_1 , for its particular value of $n = 1$. Let B_f be the fault-tolerant resolving set and stated by, $B_f = \{b_1, b_3, b_4, b_6, a_7^1, a_{13}^1\}$. Moreover its resolving set shown in the Figure [4.](#page-3-1) The locations of the full vertex set of D_1 , in regard to the elements of B_f are described as (8) and (9), as shown at the bottom of the page.

FIGURE 4. Fault-tolerant resolving set of Diamond structure D¹ .

We can look over the entire vertex set of perimantanes diamondoid structure D_1 , have unique locations and fulfilling the idea of fault-tolerant resolving set, which leads to the conclusion that defined B_f is one of a possible resolving set with $|B_f|$ $\vert = 6.$

Lemma 4: Let *Dⁿ* is a structure of perimantanes diamondoid with $n = 1$. Then

$$
dim_f(D_1)=6.
$$

Proof: We employ the double inequality approach and refer to Lemma [3](#page-3-2) to establish that the fault-tolerant metric dimension of D_1 is 6. This clearly demonstrates that the fault-tolerant resolving set defined as B_f = $\{b_1, b_3, b_4, b_6, a_7^1, a_{13}^1\}$, has cardinality 6.

Now we need to prove that $dim_f(D_1) \geq 6$. On contrary we suppose $dim_f(D_1) = 5$. For this, consider the fault-tolerant

$$
\begin{cases}\n(3-i, 5-i, 7-i, i, 8-i), & \text{if } i = 1, 2; \\
(2, 5-i, 4, 7-i, i, 8-i), & \text{if } i = 3; \\
(2, 6-i, i, 8-i, i+1, 7-i), & \text{if } i = 4; \\
(i-2, 6-i, i, 8-i, i+1, 7-i), & \text{if } i = 6; \\
(i-2, 2, i, i-2, i+1, 7-i), & \text{if } i = 6; \\
(10-i, 12-i, 10-i, i-2, i-7, 15-i), & \text{if } i = 7; \\
(10-i, 12-i, 10-i, 4, i-7, 15-i), & \text{if } i = 8; \\
(i-7, 12-i, i-7, 13-i, i-7, 14-i), & \text{if } i = 9; \\
(i-7, 13-i, i-7, i-14, i-7, 14-i), & \text{if } i = 10, 11; \\
(i-8, i-10, i-8, i-14, i-5, 13-i), & \text{if } i = 12, 13; \\
(18-i, 20-i, 16-i, i-14, i-13, 20-i), & \text{if } i = 14; \\
(18-i, 20-i, 16-i, 18-i, i-13, 20-i), & \text{if } i = 15; \\
(4, 20-i, 2, 18-i, i-13, 20-i), & \text{if } i = 16; \\
(i-13, 21-i, i-15, 19-i, i-12, 20-i), & \text{if } i = 17, 18; \\
(i-13, 4, i-15, 2, i-12, 20-i), & \text{if } i = 17, 18; \\
(i, 10-i, i-4, 6-i, i-1, 9-i), & \text{if } i = 1, 2, 3; \\
(i, 10-i, i-4, 6-i, i-1, 9-i), & \text{if } i = 4, 5, 6.\n\end{cases}
$$
\n(9)

resolving set B_f' with cardinality 5. This assumption is discussed in the following sections.

Let a chosen subset B_f' having five distinct elements, say $B'_f = \{a_i^1, a_1^j\}$ $\binom{1}{1}$, with distinct *i*, *j*. The contradiction will be arise due the vertices which have two distance to any of the chosen element of B_f' . Mathematically, it can be written as

 $r\left(a_{\alpha}^1 | B'_{f}\right) = r\left(a_{\beta}^1 | B'_{f}\right) = d\left(a_{\alpha}^1, a_{i}^1\right) = 2.$ Similarly, there is not a single case among the available combinations $|V|C_5 = \frac{|V(D_1)|!}{5!(|V(D_1)|-5)!} = \frac{2(25)!}{5!(20)!} =$ 53130 of the entire vertex set of D_1 . This indicate that five fault-tolerant metric dimension of D_1 is not possible. Hence; $dim_f(D_1) \geq 5$.

Hence,

$$
dim_f(D_1)=6.
$$

 \Box

Theorem 2: Let D_n is a structure of perimantanes diamondoid with $n \geq 2$. Then

$$
dim_f(D_n) = 2(n+2).
$$
 (10)

Proof: To show that $dim_f(D_n) = 2(n+2)$, we will applying the induction method on *n* showing the number of copies of base perimantanes diamondoid graph. The base case for $n = 1$ proved in the Lemmas [3](#page-3-2) and [4,](#page-3-3) now assume that the assertion is true for $n = \alpha$.

$$
dim_f(D_{\alpha}) = 2(\alpha + 2).
$$
 (11)

We will show that it is true for $n = \alpha + 1$. Suppose

$$
dim_f (D_{\alpha+1}) = dim_f (D_{\alpha}) + dim_f (D_1) - 5.
$$
 (12)

Using Equations 10 and [11](#page-4-0) in Equation [12,](#page-4-1) we have

$$
dim_f (D_{\alpha+1}) = 2(\alpha + 2) + 6 - 5,
$$

= 2(\alpha + 2) + 1. (13)

As a result, the conclusion holds for all positive integers $n \geq 1$.

Moreover, the generalize fault-tolerant resolving set for the generalized structure of perimantanes diamondoid, shown in the Figure [5](#page-4-2) and stated by in the set form B_f = {*b*₁, *b*₃, *b*_{3*n*+1}, *b*_{3*n*+3}, *a*₁¹, *a*₁², *a*₁³, *...*, *a*₁^{*n*}, *a*₁₁₃, *a*₁²₃, $a_{13}^3, \ldots, a_{13}^n$ }.

This concludes.

Lemma 5: Let D_n is a structure of perimantanes diamondoid with $n \geq 1$, and B_p is the vertex partitioning resolving set of D_n for $n = 1$. Then one of the possible vertex partitioning resolving set is stated by

$$
B_p = \{B_{p_1}, B_{p_2}, B_{p_3}, B_{p_4}\},
$$

\n
$$
B_{p_1} = \{b_1\}, B_{p_2} = \{b_4\}, B_{p_3} = \{a_7^1\}, B_{p_4}
$$

\n
$$
= V (D_n) \setminus \{b_1, b_4, a_7^1\}.
$$

Proof: To prove that the partitioning resolving set for the structure of perimantanes diamondoid, shown in the Figure [1](#page-2-0) and labeled as D_1 , for its particular value of $n = 1$. Let B_p be the partition resolving set and stated by,

FIGURE 5. Fault-tolerant resolving set of Diamond structure D_n.

 $B_p = \{B_{p_1}, B_{p_2}, B_{p_3}, B_{p_4}\}, B_{p_1} = \{b_1\}, B_{p_2} = \{b_4\}, B_{p_3} =$ ${a_7 \brace a_7 \brace a_8 \brace a_9 \brace a_9 \brace a_1 \brace a_1 \brace a_1 \brace a_1$. The locations of the full vertex set of D_1 , with regard to the elements of B_p are as follows.

$$
r\left(a_i^j|B_p\right)
$$
\n
$$
\begin{cases}\n(3-i, 5-i, i, 0), & \text{if } i = 1, 2; \\
(2, 4, i, 0), & \text{if } i = 3; \\
(2, i, i+1, 0), & \text{if } i = 4; \\
(i-2, i, i+1, 0), & \text{if } i = 5, 6; \\
(10-i, 10-i, i-7, 1), & \text{if } i = 7; \\
(10-i, 10-i, i-7, 0), & \text{if } i = 8; \\
(i-7, i-7, i-7, 0), & \text{if } i = 9, 10, 11; \\
(i-8, i-8, i-5, 0), & \text{if } i = 12, 13; \\
(18-i, 16-i, i-13, 0), & \text{if } i = 14, 15; \\
(4, 2, i-13, 0), & \text{if } i = 16; \\
(i-13, i-15, i-12, 0), & \text{if } i = 17, 18, 19.\n\end{cases}
$$
\n(14)

$$
r(b_i|B_p)
$$

=
$$
\begin{cases} (i-1, 3+i, i+2, z), & \text{if } i = 1, 2, 3; \\ (i, i-4, i-1, z), & \text{if } i = 4, 5, 6. \end{cases}
$$
 (15)

where

$$
z = \begin{cases} 1, & \text{if } i = 1, 4; \\ 0, & \text{otherwise.} \end{cases}
$$
 (16)

We can observe that each vertice has a different representation and satisfy the concept of partition resolving set seen from presentation of the entire vertex set of D_1 , in the form of representations which leads to the conclusion that defined *B^p* is one of a possible resolving set with $|B_p| = 4$.

Lemma 6: Let D_n is a structure of perimantanes diamondoid with $n = 1$. Then

$$
3 \leq pd(D_1) \leq 4.
$$

Proof: To show that the partition dimension of D_1 is either 3 or 4, We employ the method of twofold inequality and refer to Lemma [5.](#page-4-3) This clearly shows that the partition resolving set has the cardinality 4 and has been stated by $B_p = \{B_{p_1}, B_{p_2}, B_{p_3}, B_{p_4}\}, B_{p_1} = \{b_1\}, B_{p_2} = \{b_4\}, B_{p_3} =$ ${a_7 \brace a_7 \brace a_9 \brace a_9} = V(D_n) \setminus {b_1, b_4, a_7 \brace a_7}.$

This shows that the partition metric dimension of D_1 is four or less. Two partition dimension is reserved for the path graph with if and only if condition, so the structure of perimantanes diamondoid can not have two partition dimension and left with

$$
3 \le pd(D_1) \le 4. \tag{17}
$$

 \Box

Theorem 3: Let D_n is a structure of perimantanes diamondoid with $n > 2$. Then

$$
pd(D_n) \le n+3. \tag{18}
$$

Proof: To show that $pd(D_n) \leq n+3$, we will applying the induction method on n showing the number of copies of base perimantanes diamondoid graph. The base case for $n = 1$ proved in the Lemmas [5](#page-4-3) and [6,](#page-4-4) now assume that the assertion is true for $n = \alpha$.

$$
pd(D_{\alpha}) \leq \alpha + 3. \tag{19}
$$

We will show that it is true for $n = \alpha + 1$. Suppose

$$
pd(D_{\alpha+1}) \le pd(D_{\alpha}) + pd(D_1) - 3. \tag{20}
$$

Using Equations [17](#page-5-0) and [19](#page-5-1) in Equation [20,](#page-5-2) we have

$$
pd (D_{\alpha+1}) \leq \alpha + 3 + 4 - 3,
$$

= $\alpha + 4.$ (21)

As a result, the conclusion holds for all positive integers $n \geq 1$. This concludes.

III. CONCLUSION

In this draft, we have discussed metric dimension and their generalizations for the generalized perimantanes diamondoid structure and proved that each parameter depends on the copies of original or base perimantanes diamondoid structure and changes with the parameter *n* or its number of copies. Future direction can be considered as to discussed its other parameters which are also based on the metric of structure.

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