

Received May 18, 2021, accepted May 31, 2021, date of publication June 3, 2021, date of current version June 15, 2021. Digital Object Identifier 10.1109/ACCESS.2021.3086044

On-the-Fly Parallel Processing IP-Core for Image Blur Detection, Compression, and Chaotic Encryption Based on FPGA

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ABSTRACT This paper presents a 3 in 1 standalone FPGA system which can perform color image blur detection in parallel with compression and encryption. Both blur detection and compression are based on the 3-level Haar wavelet transform, which is used as a common building block to save the resources. The compression is based on performing the hard thresholding scheme followed by the Run Length Encoding (RLE) technique. The encryption is based on the 128-bit Advanced Encryption Standard (AES), which is considered one of the most secure algorithms. Moreover, the modified Lorenz chaotic system is combined with the AES to perform the Cipher Block Chaining (CBC) mode. The proposed system is realized using HDL and implemented using Xilinx on XC5VLX50T FPGA. The system has utilized only 25% of the available slices. Furthermore, the system can achieve a throughput of 3.458 Gbps, which is suitable for real-time applications. To validate the compression performance, the system has been tested with all the standard 256×256 images. It is shown that depending on the amount of details in the image, the system can achieve 30dB PSNR at compression ratios in the range of (0.08-0.38). The proposed system can be integrated with digital cameras to process the captured images on-the-fly prior to transmission or storage. Based on the application, the blurred images can be either marked for future enhancement or simply filtered out.

INDEX TERMS AES, blur, chaos, compression, DWT, encryption, FPGA, Haar, HDL, RLE.

I. INTRODUCTION

Nowadays, the digital images are used as a source of information in many fields, such as the social media, education, research, medical examinations, and surveillance systems. Accordingly, extensive research efforts have been conducted throughout the past years to facilitate the use of the digital images in a both efficient and secure way. The three main research fields are: image compression, encryption, and image processing. The image blur detection is one of the important applications in the field of image processing, and it will be presented along with the compression and encryption throughout this paper.

The associate editor coordinating the review of this manuscript and approving it for publication was Sedat Akleylek¹⁰.



FIGURE 1. The general structure of the lossy image compression.

A. IMAGE COMPRESSION

The image compression is one of the most popular applications that is mainly used for saving both the storage space and the network's bandwidth. The general structure of the lossy image compression is shown in Fig. 1 [1]. First, the source encoder is used to transform the input image into a sparse representation. The source encoder can be constructed using a variety of transforms, such as the Discrete Wavelet Transform (DWT), the Discrete Cosine Transform (DCT), or the Discrete Fourier Transform (DFT) [2]. The DWT is a

computationally efficient algorithm, and it has been adopted in the JPEG-2000 standard [1]. Despite the efficient compression ratios that can be achieved by the DWT, there are some limitations, such as the oscillations, the shift variance, the aliasing, and the lack of directionality [3]. The complex wavelet transform can overcome these limitations, but at the expense of the computational cost [3]. Moreover, the fractional wavelet packet transform can be used with signals that contain high fractional frequency components, but its computational cost is also higher than the DWT [4]. Therefore, the DWT will be more suitable for real-time applications that require an efficient implementation area and high operating speed.

After the source encoder, the hard thresholding is performed to reduce the number of non-zero coefficients in the transformed image. Then the quantization is used to decrease the number of bits for each non-zero coefficient. Finally, the entropy encoder is used to perform a lossless compression technique, such as the Huffman encoder, the arithmetic encoder, or the Run Length Encoder (RLE). The Huffman encoder is used in the JPEG standard and it can provide good compression results, however the RLE is more suitable for real-time applications [1].

B. IMAGE ENCRYPTION

In addition to the compression process, the images must be also encrypted before sending them over insecure communication channels. Since the encryption process destroys the correlation between the pixels, it is usually done after the compression. A lot of encryption standards are available, and there is always a tradeoff between the encryption strength and the computational cost.

The Advanced Encryption Standard (AES) is an ISO/IEC 18033-3:2010 standard for symmetric block cipher. The AES is one of the most powerful encryption standards, and it has been used in a lot of protocols, such as the IEEE802.11 wireless Local Area Network (LAN) and the IEEE802.15.4 wireless sensor networks [5]. The AES encryption has been frequently used in the Cipher Block Chaining (CBC) mode, which needs a Pseudo Random Number Generator (PRNG) to generate the Initialization Vector (IV).

The PRNGs are based on deterministic functions that are implemented in the digital domain using the computers or the Field Programmable Gate Arrays (FPGAs). In the previous years, the chaos theory has played an important role in the design of the PRNGs, image hashing [60], and modulation schemes [61]. The Rössler [6] and Lorenz [7] are two of the most famous chaotic systems, and they have been utilized in a lot of encryption applications [8], [9]. However, these two systems are based on multiplication operations that affect the hardware performance, regarding the area and speed. Hence, modified versions, called the modified Rössler and the modified Lorenz, were presented in [10], [11]. These modified versions do not utilize any hardware expensive multiplier.

In [12], six different multiplier-less chaotic PRNGs were implemented using FPGA, and their performance was



FIGURE 2. The general structure of image restoration.

compared using the NIST suite. It has been found in [12] that the modified Lorenz can pass all the NIST tests and can provide the best hardware performance as well. Hence, it will be suitable for low area and high-speed applications.

The modified Lorenz was given by (1), where a, b, and c are the system's parameters while X, Y, and Z are the system's variables. The modified Lorenz hardware implementation was presented in detail in [12].

$$\dot{X} = a \cdot (Y - X), \tag{1a}$$

$$\dot{Y} = -sgn(X) \cdot (Z - b), \tag{1b}$$

$$\dot{Z} = sgn(X) \cdot X - c \cdot Z. \tag{1c}$$

C. IMAGE BLUR DETECTION

The image blur detection is one of the fundamental applications in the field of image processing, and it is usually performed at the start of any image restoration algorithm as shown in Fig. 2 [13]. Similar to the image compression, the blur detection requires image transformation at the beginning.

A lot of no-reference image blur detection techniques have been proposed in the literature [14]. The blur metrics of these techniques are based on various approaches, such as the Haar DWT [15], the Sobel edge detection [16]–[19], the image power spectrum [20], [21], the DCT [22], [23], and a hybrid of curvelet, wavelet, and cosine transforms [24]. The simplest of these approaches is the Haar DWT [15], and it can be implemented on the hardware level without using a lot of resources.

Tong *et al.*'s [15] Haar-based blur detection algorithm can discriminate blurred images with high accuracy. This algorithm was tested on a database that include 2355 images, and the reported accuracy was 98.6% [15]. Furthermore, this algorithm can be combined with neural networks to assess the overall image quality [25]–[27]. In [27], the output of Tong *et al.*'s algorithm was considered the best performing feature that can be used for image quality assessment using neural networks. Also, Tong *et al.*'s algorithm can be used to filter out blurred images in many applications, for example the lifelogging wearable cameras [28].

D. THE FPGA-BASED IMPLEMENTATION

Image processing techniques in addition to the compression and encryption are usually implemented using software solutions that run on general-purpose processors. However, due to the sequential nature of these processors, the execution time will not be suitable for real-time applications, especially with high resolution images. This problem can be solved by using the Hardware (HW) solutions, such as the FPGAs, which can perform a lot of operations in parallel. The FPGAs are used in many applications, such as the PRNG [12],



FIGURE 3. The general structure of the FPGA-based image processing systems.

Neural Networks [29], encryption [30], edge detection [31], and compression [32].

In the past decade, the FPGAs have been used in the field of image processing due to their powerful parallel processing capabilities [33]. They can exploit the temporal and spatial parallelism of many image processing algorithms. The spatial parallelism refers to the utilization of multiple copies of the hardware to perform multiple tasks in parallel. On the other hand, the temporal parallelism refers to the utilization of a pipelined hardware unit that divides the processing task into several cascaded stages.

The FPGAs can be integrated with the digital cameras as shown in Fig. 3 [34], [35]. First, the CMOS image sensor is used to capture the image. Next, the Analog to Digital Converter (ADC) is used to convert the analog values of the image's pixels into digital values. Then the Camera Link Interface (CLI) is used to send the digital image to the FPGA. Finally, the FPGA is used to perform all the necessary processing tasks on-the-fly. The word on-the-fly means as soon as the data is received as explained in [36].

Since the FPGA's internal memory will not be sufficient to store the full uncompressed image, an external off-chip memory will be utilized as shown in Fig. 3. To emphasize this point, consider one of the most powerful FPGAs, the Xilinx Virtex 7 XC7V2000T, which includes only 46512kb Block Random Access-Memory (BRAM) [29]. This powerful FPGA can store only one 1408 \times 1408 color image with 24-bits per pixel. Hence, the off-chip memory will be needed to store high resolution images, which are in the range of tens of Megapixels.

E. MOTIVATIONS AND OBJECTIVES

Over the past few years, the wearable lifelogging cameras have been used in a lot of applications. They are used to track our daily activities by automatically capturing hundreds of images throughout the day. This vast number of images need to be stored instantly in the cloud. Hence, the compression and encryption must be done on-the-fly before transmitting the images to the cloud. Also, since a lot of the captured images may suffer from blur, most probably they will be deleted later on by the life-loggers [28]. Instead of deleting the blurred images at a later stage, a better option is to filter out the images directly after they are captured by the camera. In this way, only the unblurred images will be sent to the cloud. Accordingly, the network's bandwidth and the cloud's storage space will be saved. So, performing the blur detection on-the-fly and in parallel with the compression and encryption will be very beneficial to the lifelogging cameras. There is no solution presented in the literature, software, or hardware, that can perform the image blur detection in parallel with compression and encryption. In most systems, the blur detection is usually performed at a later stage after transmitting the compressed images to the cloud [28].

Most of the recent compression-encryption systems are implemented using software, and they suffer from long computational time [37]–[44]. Therefore, they are not suitable for real-time applications. Only few old compression-encryption systems are implemented using FPGAs, but they do not utilize a strong encryption scheme [45]–[47]. The fastest FPGA implementation is based on the DCT compression and the stream cipher [47]. The stream cipher is not as strong as the block cipher technique, for example the AES-CBC.

The main contributions of this paper are based on the following two aspects:

- To the best of the authors' knowledge, this paper presents the first HW implementation of Tong *et al.* [15] Haar-based blur detection algorithm.
- (2) The blur detection is integrated with image compression and encryption as a 3 in 1 parallel HW solution, which is suitable for real-time applications. Depending on the application, the blurred images can be either marked for future enhancement or simply filtered out. The compression is based on both the RLE and the Haar DWT. The encryption is implemented using one of the most secure algorithms: the 128-bit AES, which is combined with the modified Lorenz chaotic PRNG to perform the CBC mode. The proposed system can be integrated with digital cameras to process the captured images prior to transmission or storage. Once the Npixels image is fully delivered to the frame buffer, the proposed system will just read the image in only Nclock cycles, and all the tasks will be performed onthe-fly in a pipelined manner.

The paper is organized as follows. Section I presents the introduction. Section II presents the HW implementation of the blur detection algorithm. Section III presents the HW implementation of the full system. Section IV presents the results. Section V presents the conclusion.

II. THE HW IMPLEMENTATION OF THE IMAGE BLUR DETECTION

A. THE HAAR DWT

The 1-level 2d Haar transform is illustrated as shown in Fig. 4(a) [48]. Every four neighboring pixels in the original image, such as A, B, C, and D, are transformed into X, Y, Z, and W. X is called the average coefficient while Y, Z, and W are called the vertical, horizontal, and diagonal detail coefficients, respectively. The transformed matrix will be divided into four bands: Low-Low (LL), High-Low (HL), Low-High (LH), and High-High (HH). The LL band will contain the average coefficients while the other bands will contain the detail coefficients. The 3 decomposition levels are obtained by performing the 1-level 2d transform recursively on the LL band as shown in Fig. 4(b).



FIGURE 4. a) The 1-level 2d DWT, b) The 3-level 2d DWT (pyramid structure), c) The implementation of the 1-level 2d DWT using the convolutional method, d) The HW implementation of the 1-level 2d DWT, and e) The HW implementation of the 3-level 2d DWT.

The 1-level 2d DWT is built using the conventional convolutional method, which consists of a High Pass Filter (HPF) and a Low Pass Filter (LPF) followed by down samplers, as shown in Fig. 4(c) [49]. The Haar transform's LPF and HPF are based on computing the sum and difference between two consecutive inputs as given by [48]:

$$G[Z] = \frac{1}{2} \left(1 + Z^{-1} \right), \qquad (2a)$$

$$H[Z] = \frac{1}{2} \left(1 - Z^{-1} \right).$$
 (2*b*)

The hardware implementation of the 1-level 2d Haar transform is presented in Fig. 4(d). The registers are used to

provide the delay elements in (2), the hardwired shifters are used to perform the multiplication by 1/2, and the enable signals, E_1 and E_2 , are used to control the down sampling. The fractional bit that arises after the multiplication by 1/2 will be truncated to maintain the same number of bits throughout the design.

The 3-level 2d DWT is implemented using cascaded 1-level 2d DWT units as shown in Fig. 4(e). The enable signals, E_1 to E_6 , are used to control the down sampling of the three levels. E_1 is always enabled while E_2 , E_3 , E_4 , E_5 , and E_6 are enabled once every 2, 4, 8, 16, and 32 clock cycles, respectively. To use these cascaded units properly, the image



FIGURE 5. A demonstration of the proposed scanning technique on a 4×4 image, which is stored row by row as indicated by the index of each pixel.

must be scanned in a certain order, which is different from the conventional row by row technique.

Once the N pixel image is available in the frame buffer, it will be scanned as shown in the example in Fig. 5. In this example, a 4×4 image is stored row by row as indicated by the memory location written in each block. The image scanning is performed as follows. First, the image is divided into four quadrants as depicted by the dashed squares. Then, each quadrant is divided again into four quadrants. This process is performed until the smallest quadrant, which contains only 1 pixel, is reached. Finally, every four equal sized quadrants are scanned in this order: upper left, upper right, lower left, then lower right. For the example in Fig. 5, the memory addresses must be generated in this order: 0, 1, 4, 5, 2, 3, 6, 7, 8, 9, 12, 13, 10, 11, 14, 15.

Accordingly, a 4-bit counter will be used to generate the 16 different addresses, but the counter's bits will be rearranged in a different order. Suppose that the counter's bits are normally arranged as $A_3A_2A_1A_0$, then to generate the required sequence, the bits must be reordered as $A_3A_1A_2A_0$. This means that starting from the Least Significant Bit (LSB), the even bits will be inserted before the odd bits. This concept can be applied to any n-bit counter. The proposed scanning technique can be applied on square images with any size (N pixels) to perform the 3-level DWT in only N clock cycles. For the case of non-square images, the zero-padding method can be used to change the input image to a square.

B. THE IMAGE BLUR DETECTION

Tong *et al.*'s blur detection algorithm [15] is based on classifying the image edges into four types: Dirac, A-step, G-step, and Roof. The classification is based on the variation in the pixels' intensity as shown in Fig. 6(a) [15]. Blurred images usually do not have Dirac or A-step edges. Instead, they have a lot of G-step and Roof edges. Accordingly, the algorithm can detect blurred images by identifying the type of edges in the image.

The first step of the algorithm is to perform the 3-level 2d Haar wavelet transform. The second step is to construct an edge map for each level using (3a) for i = 1, 2, 3. To improve the hardware performance, the edge maps can be



FIGURE 6. a) The graphical description of different edge types. b) The rules for discriminating the edges based on the maximum edge values at the 3 decomposition levels.

approximated using (3b).

$$Emap_i = \sqrt{LH_i^2 + HL_i^2 + HH_i^2}, \qquad (3a)$$

$$Emap_{i} \approx |LH_{i}| + |HL_{i}| + |HH_{i}|. \tag{3b}$$

The third step of the algorithm is to divide the edge maps using partitioning windows. Since the size of the edge map is scaled down by a factor of 4 after every decomposition, the partitioning windows must be scaled down by the same factor. Therefore, the sizes of the partitioning windows used for level 1, 2, and 3 are 8×8 , 4×4 , and 2×2 , respectively. In this way, the three edge maps will have the same number of partitions.

The fourth step is to find the maximum value, $Edge_{max i}^{j}$, in each partitioning window where *j* is the index of the partition and *i* is the decomposition level. If $Edge_{max1}^{j}$ or $Edge_{max2}^{j}$ or $Edge_{max3}^{j}$ is higher than a certain threshold (*Th*), then the partition with index *j* will have an edge point. The type of the edge can be identified using the rules presented in Fig. 6(b). Furthermore, in case the edge is identified as a G-step or Roof, but $Edge_{max1}$ is lower than *Th*, then it will be considered as an unsharp edge point. The final step is to divide the total number of Dirac and A-step edges (N_{DA}) by the total number of edges (N_{edges}). If this ratio is lower than a very small threshold (e.g., 0.05), then the image will be considered blurred. The blur extent is calculated by dividing the total number of unsharp G-step and Roof edges (N_{BRG}).

The hardware implementation of the blur detector unit is presented in Fig. 7. The blur detector unit will receive the detail coefficients of levels 1, 2, and 3 from the DWT unit every 4, 16, and 64 clock cycles, respectively.

First, the sum of the absolute values of *HH*, *HL*, and *LH* is computed for each level to generate *Edge*₁, *Edge*₂, and *Edge*₃.

The computation of the Edge value, (|LH| + |HL| + |HH|), is implemented using 3 cascaded adders/subtractors as



FIGURE 7. The HW implementation of the blur detection algorithm.

 $(0 \pm LH \pm HL \pm HH)$ where the adder/subtractor's control signal is driven by the sign bit of the detail coefficient.

Next, the registers connected to $Edge_1$, $Edge_2$, and $Edge_3$ will be enabled only if the generated edge value is larger

than the stored value. In this way, the registers will store the maximum value they receive. Furthermore, to allow the registers to be updated only when the correct edge value is generated, the comparators of levels 1, 2, and 3 will be ANDed with the input flags F_1 , F_2 , and F_3 , which are enabled once every 4, 16, and 64 clock cycles, respectively.

The proposed image scanning technique allows the system to scan a new 16×16 block from the input image every 256 clock cycles. Hence, the maximum values of the current edge map partition, *EdgeMax*₁, *EdgeMax*₂, and *EdgeMax*₃, will be available in the registers at the end of every 256 clock cycles. After finding these maximum values, the registers will be reset using *resw* before proceeding to the next partition. The edge type detector will receive the maximum values and will update the counters according to the rules in Fig. 6(b). Finally, two dividers are used to compute the required ratios, one to check whether the image is blurred or not, and the other to calculate the blur extent. The *resBD* signal is used to reset the counters and the dividers before processing a new image.

III. THE PROPOSED SYSTEM

The HW architecture of the proposed system is presented in Fig. 8. The proposed architecture performs the color image blur detection in parallel with the compression and encryption in a pipelined fashion. The image is sent to the system pixel by pixel where each color pixel *RGB* consists of 24 bits.

The main data path of the system is explained as follows. First, the input color pixel is transformed from RGB into



FIGURE 8. The block diagram of the proposed system.

HI

sel_{COEF}(5:4)

res

res

HH

sel_{COEF}(3:2)

res

10

01

10

01

HH

10

01

2-bit

counter

2

Q

sel_{COEF}(1:0)

res

En

En₁



FIGURE 9. The HW implementation of the sampler.

YCbCr components. Then each color component will pass through a separate 3-level DWT unit. Next, the blur detection unit will process the detail coefficients of the Y color component only. Finally, three compression-encryption units are used to encode the 3 color components in parallel. Each compression-encryption unit will use the 10 output channels of the corresponding 3-level DWT unit in a sequential manner. Therefore, a multiplexing unit is needed to interface the 10 output channels of the 3-level DWT unit with the compression-encryption unit. The encrypted outputs are provided in words of 32-bits. Also, flag signals are generated to indicate that a new encrypted data is available at the outputs.

The IV is generated from the modified Lorenz PRNG which uses a 65-bit seed value. The modified Lorenz is implemented using the design in [12], which can past all the NIST tests [50]. This PRNG generates 24-bits per clock cycle; therefore, a 128-bit shift register is needed to receive the generated bits sequentially, and then provide all the 128-bits in parallel.

Finally, the main control unit is used to synchronize and control all the blocks in addition to incrementing the address generator. The *address* is used to read the image from the external frame buffer. Moreover, the control unit uses other internal signals as inputs, such as, *address*, F_1 , F_2 , F_3 , res_W , sel_{COEF} , F_Y , F_{Cb} , and F_{Cr} ; however, they are not shown in Fig.8 just for the sake of simplicity. The control unit will be explained in detail in a subsequent section.

A. THE SAMPLER

The sampler circuit is used to generate the enable signals of the 3-level 2d DWT unit in addition to the flag signals of the blur detector. The sampler circuit consists of an 8-bit counter and a set of comparators as shown in Fig. 9. The flag signals, F_1 and F_2 , are the same as the sampler signals, E_3 and E_5 , respectively.

B. THE COLOR CONVERTER

The color converter is used to change the input RGB components to YCbCr components where Y is the luminance while



component is a prerequisite for the blur detection. The conversion is performed as shown in (4) [51]. The conversion requires 3 multipliers and 3 adders for each component to compute Y, C_b, and C_r in parallel. Thus, the converter utilizes a total of nine 8-bit multipliers and nine 8-bit adders. The 8-bit multipliers will multiply the input component (R, G, or B) with the corresponding numerator in (4), (e.g., $R \times 65$). Accordingly, the output of the multiplier will have 16 bits. Only the 8 Most Significant Bits (MSBs) of the multiplier will be taken, which is equivalent to dividing by 256 and then truncating the fraction part. Furthermore, the adders are connected in a tree structure to reduce the critical path.

applied on grayscale images only; therefore, obtaining the Y

$$Y = 16 + \frac{65}{256}R + \frac{129}{256}G + \frac{25}{256}B,$$
 (4*a*)

$$Cb = 128 - \frac{38}{256}R - \frac{74}{256}G + \frac{112}{256}B, \qquad (4b)$$

$$Cr = 128 + \frac{112}{256}R - \frac{94}{256}G - \frac{18}{256}B.$$
 (4c)

C. THE MULTIPLEXING UNIT

The MUX unit and the selector are shown in Fig. 10. This unit consists of three cascaded 4-input multiplexers, which are used to select between the coefficients of the 3 levels. The select signals of the multiplexers are generated from the 2-bit counters. The *Ensel* is used to enable or disable the counters.

The $Mode_{sel}$ is used to choose the counting mode. If Modesel is set high, the counters will operate normally



FIGURE 11. The HW implementation of the compression-encryption unit.

from 0_{10} to 63_{10} . Otherwise, the *sel_{COEF}* will change, such that the last 10 coefficients will pass in this order: *LH*₁, *HL*₁, *HH*₁, *LH*₂, *HL*₂, *HH*₂, *LH*₃, *HL*₃, *HH*₃, *LL*₃.

D. THE COMPRESSION-ENCRYPTION UNIT

The HW implementation of the compression-encryption unit is presented in Fig. 11. First, Mux1 is used to perform the hard thresholding by selecting between In and 0_{10} . The selection is based on whether In is greater than the compression threshold (comp_{TH}) or not. Next, an 8-bit counter will count the number of zero bytes that pass through MUX₁. Accordingly, Z_{COUNT} and the output of Mux₁ will form the RLE 16-bit word. For example, {00₁₆, 00₁₆, 00₁₆, 00₁₆, 0F₁₆,} will be encoded as $\{04_{16}, 0F_{16}\}$. The RLE word will be inserted into a 128-bit shift register (Reg₁), which is enabled using the RLE_{FLAG} . The main problem of the RLE can be explained using the following example. Suppose that $comp_{TH}$ is set to 3_{10} , and the input sequence has 14 consecutive coefficients greater than or equal to $comp_{TH}$ (e.g., $In = \{03_{16}, 04_{16}, 05_{16} \dots 0F_{16}, 10_{16}\}).$ Accordingly, this sequence will be encoded as $\{00_{16}, 03_{16},$ $00_{16}, 04_{16} \dots 00_{16}, 10_{16}$. This implies that the 14 input bytes will be encoded in 28 bytes, which has a negative effect on the compression performance. To solve this problem, the In bytes will be inserted directly into another 112-bit shift register (Reg₂), which is enabled by En_{RLE} .

In addition, 2 zero bytes will be concatenated with Reg₂ to complete the 128-bits and indicate that they are not run length encoded. Accordingly, the 14 input bytes will be encoded in

16 bytes as $\{03_{16}, 04_{16}, \dots, 0F_{16}, 10_{16}, 00_{16}, 00_{16}\}$, which is in this case better than the RLE. After filling Reg1 and Reg₂, MUX₂ will select between these 2 registers as follows. If Reg₁ is filled before Reg₂, then MUX₂ will select Reg₂ as explained in the previous example. Otherwise, the RLE will be the better option, and MUX₂ will select Reg₁. To control MUX₂, 2 flip flops are used to hold the present state of Reg₁ and Reg₂ (P_1, P_2) . The 3-bit counter is used to count the number of RLE words that enter Reg₁. Similarly, the 4-bit counter is used to count the number of In bytes that enter Reg₂. Reg₁ will be filled after receiving 8 RLE words while Reg₂ will be filled after receiving 14 In bytes. Once the register is filled, its present state will be set to logic 1. Since the In bytes are received every clock cycle, the AES unit will have at least 14 clock cycles to encrypt the input block. The AES will encrypt the 128-bit block in only 11 clock cycles as will be explained in the next subsection. Finally, a logic unit is used to control all the blocks. The En_{RLE} is used to enable the compression-encryption unit. The Mode_{RLE} is used as an indicator for the last In byte. The Last BLOCK is used to start the encryption of the last 128-bit block even if Reg1 and Reg2 are not yet filled.

E. THE AES

The HW implementation of the 128-AES-CBC is shown in Fig. 12. All the details of the AES are available in [52]. The encryption is performed as follows. The initialization vector (IV) will be updated and loaded into Reg₁ at the start



FIGURE 12. The HW implementation of the 128-bit AES-CBC.

of every new image. Then the data in Reg_1 will be used as a feedback and XORed with the input to perform the CBC mode. After feeding back the content of Reg_1 , the AES will start the encryption process, which consists of 11 rounds. Each round will have a unique Round Key (*Rkey*), which will be generated using the key expander.

In the proposed system, each round will be performed in 1 clock cycle. In the first clock cycle, the *Rkey* will be added to the input, and the result will be stored in Reg2. In the next 9 clock cycles, the data stored in Reg₂ will be updated after applying a series of operations: the sub bytes substitution using the S-box, then the shift rows, followed by the Mix columns, and finally adding the *Rkey*. In the 11th clock cycle, the Mix columns operation will be skipped, and the final result will be stored in the output register Reg₁. The output flag F_Y is used to indicate that the encrypted block is now available in Reg1. The encrypted block will be provided through *Out*_Y in 32-bit words using a 4-input multiplexer, which is controlled by a 2-bit counter.

As explained above, the operations performed in the initial round is different from the middle 9 rounds and the last round as well. Hence, multiplexers are used to select the required operations in each round. The multiplexers will be controlled using the LSB and the MSB of an 11-bit left circular shift register. Initially, the shift register will be loaded with 1_{10} ; thus, the LSB will be high while the MSB will be low. Accordingly, the data coming from the input side will be selected and only the Add round key operation will be performed. In the next 9 cycles, the bits will be moving in the shift register towards the left, and both LSB and MSB will be

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at logic 0. Hence, all the AES operations will be performed. In the 11th cycle, the MSB will be high while the LSB will be low; therefore, the Mix columns will be skipped.

The Add round key block is implemented using a group of XOR gates. The sub bytes block is implemented using 16 parallel S-boxes, which are implemented using Look Up Tables (LUTs). The shift rows block is implemented using hardwired shifters; hence, no computational resources will be utilized. Finally, the Mix columns block is implemented using a series of XOR operations as explained in [52].

The key expander is implemented as shown in Fig. 13. The key expansion process is performed as follows. First, the input key will be loaded in a 128-bit register, which will provide the *Rkey*. Next, the 32 LSBs of the 128-bit register will pass through a series of operations before they are combined back with all the 128-bits. The 32 LSBs will first undergo a hardwired byte cyclic rotation. Then 4 parallel S-boxes will be used to substitute the rotated bytes. Accordingly, the 4 S-boxes will provide a total of 32 bits, only the 8 MSBs will be XORed with the *Rcon* signal, and then they will be concatenated back with the remaining 24 LSBs.

To generate the *Rcon* signal, first a value of 1_{10} will be loaded in an 8-bit register. Next, the register will be updated by shifting the stored bits to the left and then adding the signal *con*. The signal *con* is constructed using *Rcon* as follows $\{0, 0, 0, Rcon(7), Rcon(7), 0, Rcon(7), Rcon(7)\}$.

F. THE MAIN CONTROL UNIT

The main control unit is implemented using a Finite State Machine (FSM) as shown in Fig. 14. To simplify the state



FIGURE 13. The HW implementation of the AES key expander. a) The Key expander block diagram. b) The Rcon function.

diagram, the outputs written in each state are the ones that are set to logic 1. As illustrated in Fig. 14, The FSM will jump from the initial state, St_0 , to St_1 in case the input *start* is set to logic 1. In St_1 , the FSM will start to increment the address of the frame buffer to read the image. The FSM will spend 1 clock cycle in St_1 until the first pixel is read from the frame buffer. The next state, St_2 , is used to add another 1 clock cycle delay due to the first pipeline stage in the system, see *pipeline*₁ in Fig. 8. In case the system is not pipelined, the FSM will jump directly from St_1 to St_3 .

In St_3 , the FSM will enable the selector and the sampler to start the image processing. The FSM will remain in St_3 for 4 clock cycles until the first set of coefficients are generated from level₁ as indicated by the input flag F_1 .

Next, the FSM will jump to St_4 to enable the RLE. Since the coefficients of level₂ are not generated yet, the FSM will enable the RLE for 3 times every 4 clock cycles to encode the 3 detail coefficients of level₁ only. The FSM will remain in St_4 until the first set of coefficients are generated from level₂, as indicated by the input flag F_2 . Then the FSM will jump to St_5 . Similarly, the FSM will remain in St_5 until the first set of coefficients are generated from level₃, and then will jump to St_6 .

In St_6 , the RLE will be enabled all the time to encode the coefficients of all the 3-levels sequentially. The FSM will remain in St_6 until the image is fully read from the frame buffer and then will jump to St_7 . The FSM will wait for 1 clock cycle in St_7 until the counters of the blur detector are updated, and then will jump to St_8 to start the division process of the blur detector. Also, in St_7



FIGURE 14. The FSM of the control unit. The inputs and outputs are written in blue and black colors, respectively.



FIGURE 15. The standard test images.

and St_8 , the last 10 coefficients will be multiplexed to the compression-encryption unit in the order explained in the previous subsection. Accordingly, the *Mode_{sel}* will be set to logic 0. The FSM will remain in St₈ until the last coefficient from the transformed image is encoded as indicated by $(sel_{COEF} == 63_{10})$.

Next, the FSM will jump to St_9 or St_{10} according to whether the AES units are ready to process the last block or not. Suppose that one of the AES units has started to encrypt a given block and there are still 2 non-zero coefficients left in the transformed image. In this case the AES buffer will be ready after 2 clock cycles only while the AES unit will still need another 9 clock cycles to finish the encryption process. Accordingly, the FSM will move to St_9 and wait until all the AES units are done. In St_{10} , the system will encrypt the last block in the image. St_{11} and St_{12} are delay states to ensure that the 128-bit encrypted block is fully transferred through the 32-bit output port of the AES unit. Finally, the FSM will wait in St_{13} until a new image is received, and then will repeat the same process again.

IV. RESULTS

The proposed system is realized using HDL and implemented on XC5VLX50T FPGA using Xilinx. The simulation process is performed in three steps. First, the input image is written in a Memory Initialization File (MIF) using MATLAB. Next, the HDL design is tested with the MIF, and then the encrypted image is written in a text file using Xilinx simulator. Finally, the generated text file is decrypted and decompressed using MATLAB to verify the compressionencryption process. The standard test images are summarized in Fig. 15. The proposed system is evaluated using both color and grayscale images to ensure a fair comparison with most of the recent compression-encryption systems. In case of grayscale images, the YCbCr conversion will be skipped.

A. THE HDL SIMULATION RESULTS

The HDL simulation results are illustrated in 3 parts. The first part will verify the YCbCr conversion, the 3-level DWT, and

 TABLE 1. The post placement and routing results.

Logic utilization	Used	Available	Utilization					
Slices	1866	7200	25%					
Registers	2955	28800	10%					
LUTs	5736	28800	19%					
Frequency	144.11 MHz							
Throughput	3.458 Gbps							

the blur detection. The second part will verify the AES. The last part will verify the compression scheme. The simulation is performed on Lena 256×256 color image.

The first part of the simulation is presented in Fig. 16. The Y component of the color converter, which is the grayscale image, is presented in Fig. 16(a). To see how the edge maps are generated, the simulation will focus on a block of 16×16 pixels near Lena's hat as depicted in Fig. 16(a). This block of pixels is magnified in Fig. 16(b) just for the sake of illustration. Figures 16(c) to 16(e) show the generated edge maps for the 3 levels. In case the edge threshold is set to 10, the position of the edge points will follow the edges in Lena's hat. The Xilinx simulation results for the edge map of level 2 is presented in Fig. 16(f)

The second part will verify the AES using the example in [52]. The HDL simulation results, shown in Fig. 17, match exactly the results in [52].

The last part of the simulation, which verifies the compression operation, is shown in Fig. 18 where the $comp_{TH}$ is set to 3_{10} . The input sequence consists of 7 consecutive bytes with magnitude less than or equal to the $comp_{TH}$ followed by the non-zero coefficient 4_{10} . Accordingly, the output RLE word will be 0704₁₆.

B. THE HW PERFORMANCE

The proposed system is fully implemented using slice registers and LUTs. No DSP48Es or block RAMs are utilized. The post placement and routing results are summarized in



FIGURE 16. The HDL simulation results of the blur detector. a) The test image, b) The cropped part, c) Edge map 1, d) Edge map 2, e) Edge map 3, and f) The output of Xilinx's simulator for Edge map 2.

	Round 0	Round 1		Round 2		Round 3	Ro	und 4	Round	5
clk										
R _{KEY0}	09cf4f3c	2a6c	7605	7359	f67f	6d7a88	3b X	db0bad00) 11f9	15bc
R _{KEY1}	abf71588	23a3	3939)	5935	807a	(1e237e	44 X	b671253b	Caf2	b8bc
R _{KEY2}	28aed2a6	8854	2cb1	(7a96	o943	4716fe	3e)	a8525b7f	7c83	9d87
R _{KEY3}	2b7e1516	a0fa	fe17)	(f2c2	95f2	3d8047	7d)	ef44a541		c6f8)
Word ₀	193de3be	(a49c	7ff2)	aa8f	5f03	486c4e	ee X	e0927fe8	f100	6f55
Word ₁	a0f4e22b	689f	852b)	(61dd	e3ef	(671d9d	lod X	c86353c0	√ c192	4cef
Word ₂	9ac68d2a	6b5b	ea43 🔿	(82d2	4ad2	(4de3p1	.38)	d9b13550	X 7cc8	8b32
Word ₃	e9f84808	026a	5049)	6832	469a	d65f58	e7 X	85b8be01		d50c

	Round5	Round 6	Round 7	Round 8	Round 9	Round 10
clk						
R _{KEY0}	11f915bc	ca0093fd	(4ea6dc4f)	7f8d292f	575c006e	b6630ca6
R _{KEY1}	X caf2b8bc	(dbf98641)	(84a6 <mark>4fb</mark> 2)	312bf560	28d12941	e13fDcc8
R _{KEY2}	X 7c839d87	110b3efd	(5f5fc9f3)	b58dbad2	19fadc21	(c9ee2589)
R _{KEY3}	X d4d1c6f8	(6d88a37a)	(4e54f70e)	ead27321	ac7766f3	(d014f9a8)
Word ₀	X f1005f55	260e2e17	(5a4142b1)	ea835cf0	eb40f21e	(3925841d)
Word ₁	X c1924cef	(<u>3d41þ77d</u>)	(1949dc1f)	0445332d	592e3884	O2dc09fb
Word ₂	X 7cc88b32	(<u>e864</u> 72a9)	(<u>a3e01965</u>)	655d98ad	8ba113e7	(dc118597)
Word ₃	X 5db5d50c	(fdd28b25)	(<u>7a8c</u> 040c)	(<u>8596</u> 0c5	1bc342d2	(<u>196a</u> 0b32

FIGURE 17. The AES HDL simulation results for the testcase in [52], which shows the encrypted words after each round.



FIGURE 18. The HDL simulation of the compression process with a threshold value equal to 3.

Table 1. The proposed system has utilized 25% of the available slices. This is a very efficient utilization relative to the

functionalities offered by the system: color image blur detection, compression, and AES-CBC encryption. Furthermore,

System	Image	Detect Blur	Compression	Cipher	Device	Slice LUT	Slice Reg	Slices	DSP48E	Clock (MHz)	Throughput (Mbps)
[45]	Gray	Х	Haar DWT SLCCA	128-AES (ECB)	Altera EPF10K250A	NA	NA	NA	NA	40	330
[46]	Gray	Х	JPEG	TEA	Virtex-6	2130	1456	NA	9	87.290	698.32
[47]	Gray	Х	DCT	Stream cipher	XC5VLX330	2058	1536	NA	0	NA	1648
[62]	NA	Х	Х	Keccak-f[400] 128-AES	Virtex-6	NA	NA	809	NA	332.98	2131.07
This work	Color	\checkmark	Haar DWT RLE	128-AES (CBC)	XC5VLX50T	5736	2955	1866	0	144.11	3458

TABLE 2. The device logic utilization and throughput for different FPGA-based compression-encryption systems.

the timing analysis shows that the maximum clock frequency is 144.11 MHz. Since the system processes a 24-bit color pixel per clock cycle, the throughput will be 3.458Gbps. This throughput allows the system to process 1-MP image in only 7.276ms. Therefore, the rate will be 137 Frames Per Second (FPS) for 1-MP or 30 FPS for 4.5-MP. Accordingly, the proposed system is well suited for real-time applications.

The HW performance of different FPGA image compression-encryption systems are summarized in Table 2. The proposed system is the only one capable of integrating the blur detection with image compression-encryption. Furthermore, the proposed system can process color images, and not just grayscale images as the other systems. Moreover, the proposed system has the highest throughput, which is even twice as fast as a simple DCT compression and stream cipher system [47]. Also, the throughput of the proposed system is 62% higher than the FPGA-based system presented in [62], which accommodates a single AES-128 core with Keccak-f[400]. Regarding the HW implementation area, the proposed system has the largest slice utilization. This is due to the 128-AES-CBC, which improves the security level at the expense of utilizing more resources. The AES-CBC is more secure than the TEA, the stream cipher, and the AES-ECB. Also, the implementation area has increased due to the YCbCr conversion and the blur detection, which are not available in the other systems.

Most of the image compression-encryption systems are implemented using software solutions. Accordingly, the proposed FPGA system is also compared with other CPU based systems, which use new and non-standardized encryption methods, as shown in Table 3.

Like the previous discussion, the proposed system is superior in terms of speed. Only [37] and [40] work on color images, however their Execution Time (ET) is not suitable for real-time applications. According to the results in Table 3, the proposed system is 13.9 times as fast as the best CPU based implementation [57].

Furthermore, the ET of the proposed FPGA system is compared with the software implementation of image blur detection. In [58], Tong *et al.* blur detection algorithm was implemented on Jetson TK1 using different approaches: the single-core Sequential approach [15], the multi-core parallel



FIGURE 19. The blur detection ET for 1024×1024 color images using different implementations of Tong *et. al.*'s algorithm. The SW solutions were implemented on Jetson Tk1.

approach [59], and the GPU parallel approach [58]. The comparison is presented in Fig. 19, which shows that the FPGA implementation is 1.773 times as fast as the GPU-based implementation.

C. THE BLUR DETECTION RESULTS

The blur detection feature of the proposed system is tested on 256 \times 256 images with different blur levels. The blur threshold of the algorithm is set to 10. The input image is considered blurred if the ratio of N_{DA} to N_{edges} is less than 5%, otherwise, the image is considered unblurred. The simulation results are shown in Fig. 20 where the images on the left are classified by the system as unblurred while the other images are classified as blurred. Furthermore, by examining the same image at different blur levels, it is clear that the estimated blur extent can be used to rank the images according to the amount of blur.

D. THE COMPRESSION-ENCRYPTION RESULTS

The compression performance of the system is assessed by calculating the Peak Signal-to-Noise Ratio (PSNR) versus the Compression Ratio (CR). The CR is calculated as shown in (5). The PSNR is calculated as shown in (6a), (6b), and (6c) where *MSE* is the Mean Square Error, $M \times N$ is the dimensions of the image, $f_1(x, y)$ is the plain image, and $f_2(x, y)$ is the reconstructed image. In color images, the PSNR is calculated for each color component, (Red, Green, and Blue), and then the average value is calculated as shown in (6c). In lossy

	Device		Image	Dotoot		256×256	512×512	1024×1024	
System	Processor RAM Clock type Blur Compression		Compression	ET (s)	ET (s)	ET (s)			
[37]	Core i3- 4005U	2	1.7	Color	Х	DWT	12.18	47.42	NA
[38]	Core i7-6500U	4	3.1	Gray	Х	Comp. sensing	7.07	NA	NA
[39]	Core i7	8	3.4	Gray	Х	Comp. sensing	6.2	NA	NA
[40]	NA	4	3.3	Color	Х	Comp. sensing	1.1168	3.9593	NA
[41]	Core i7-4770K	16	3.5	Gray	Х	JPEG	0.98	2.64	8.33
[42]	Core i5	4	3.1	Gray	Х	DWT, SPIHT	0.773	3.56	NA
[43]	Core i7-3667U	8	2.5	Gray	Х	Comp. sensing	0.4769	1.4712	NA
[44]	Core i7-6700	8	2.9	Gray	Х	Comp. sensing	0.2387	0.731	NA
[53]	Core i5-6500	8	3.2	Gray	Х	Comp. sensing	0.03178	0.10335	NA
[54]	Core i3	4	2.13	Gray	Х	Comp. sensing	0.01983	0.054	NA
[55]	Core2 Duo	2	2.4	Gray	Х	Chinese Remainder	0.018	0.05	0.268
[56]	Core i3–2120	2	3.3	Gray	Х	Comp. sensing	NA	0.0278	NA
[57]	Core 2 Duo	4	1.86	Gray	Х	DCT, Huffman	NA	0.02632	NA
This work	nis FPGA XC5VLX50T		0.144	Color	\checkmark	DWT, RLE	0.00046	0.00182	0.007276

TABLE 3. The Execution Time (ET) comparison with different CPU-based image compression-encryption systems.

image compression, the PSNR is inversely proportional to the CR. It is important to keep the PSNR above certain limit to achieve good image quality. The reconstructed image quality is considered good if the PSNR values are between 30dB and 50dB. The compression is unacceptable if the PSNR is below 20dB [37].

$$CR = \frac{size \ of \ cipher \ image}{size \ of \ plain \ image}$$
(5)

$$MSE = \frac{1}{M \times N} \sum_{y=1}^{M} \sum_{x=1}^{N} (f_1(x, y) - f_2(x, y))^2 \qquad (6a)$$

$$PSNR = 10 \times \log_{10} \left(\frac{255^2}{MSE}\right) \tag{6b}$$

$$PSNR_{RGB} = \frac{PSNR_{Red} + PSNR_{Green} + PSNR_{Blue}}{3}$$
(6c)

The simulation results of the compression-encryption process are shown in Fig. 21. The CR ratio is adjusted to approximately 20% for the three 256×256 test images, "Baboon", "Barbara", and "Lena". The images are successfully ciphered by the AES as shown in the figure. Regarding the quality of the restored images, the "Baboon" has the lowest PSNR while "Lena" has the highest PSNR. This is due to the amount of details in the image. Regarding the "Baboon" image, which has a lot of details, the compression threshold must be set high to reach the 20% CR at the expense of lowering the PSNR.

For further investigation, the PSNR vs the CR is evaluated for all the standard test images. For each image, the compression threshold is increased until the PSNR drops to 30dB. Since the compression threshold is restricted to integer values, it is hard to get exactly the 30dB point. Accordingly, the point closest to the 30dB will be marked for each image. The PSNR vs the CR for all the color and gray images are presented in Fig. 22 and Fig. 23, respectively. Each image is tested with two resolutions, 256×256 and 512×512 as shown by the red and blue curves, respectively. By inspecting the 30dB point for each test image, it is found that the CR is greatly affected by the amount of details in the image. For example, in the 256×256 images, the "House" has the best CR, which is 0.08 and 0.11 for color and grayscale, respectively. On the other hand, the "Baboon" has the worst CR, which is 0.38 and 0.75 for color and grayscale, respectively.

Moreover, the results show that the CR is better in color images compared to the grayscale images. This is due to the YCbCr conversion, which puts most of the details in the Y component as the human visual system is more sensitive to the luminance than chrominance. By inspecting the 30dB point in the 256×256 "House" and "Baboon", it is found that the CR of the color image is better than the grayscale image by 27% and 49%, respectively.

Furthermore, the results show that the compression performance is improved at higher resolutions. For example, by examining the 30dB point of the "Plane" color image, the CR at resolution 512×512 is better than 256×256 by 43.75%.

The PSNR vs CR of the proposed system is compared with other systems as shown in Fig. 24. The results show that the proposed system can achieve high PSNR at low CR. By inspecting the lowest CR in the figure, which is the 25%, it is found that the PSNR of the proposed system is better than the other systems by more than 8%. However, the compression scheme in [40] can provide better PSNR at CRs higher than 50%. This is due to the truncation rounding scheme of the YCbCr converter, which limits the maximum PSNR of the proposed system to 33dB.

In addition to the PSNR, the Mean Structural Similarity (MSSIM) index is also used to measure the similarity







Plane (unblurred)

between the original and the restored images from three

aspects: brightness, contrast, and structure. The MSSIM

can be calculated using the equations presented in [38].

for each image.

Blur extent= 21.8045% FIGURE 20. The Simulation results of the blur detection for different 256 × 256 test images. The calculated blur extent is written

Table 4 compares the MSSIM of the proposed system with one of the most recent compression-encryption systems [38] at 25% CR.



Three different 256×256 grayscale images are used in work table 4. The results show that the MSSIM of the proposed "I

work is close to the system presented in [38]. In the case of "Lena" and "Peppers", the MSSIM is 9% lower than the



FIGURE 23. The PSNR vs CR of the proposed system for several gray images. The red and blue curves represent the 256 \times 256 and 512 \times 512 images, respectively.



FIGURE 24. The compression performance comparison for different compression-encryption systems.

 TABLE 4. The MSSIM index of the proposed system at 25% CR.

Image	This Work	[38]
Lena (256×256)	0.5710	0.6211
Pepper (256×256)	0.5733	0.6354
Man (256×256)	0.5695	0.5375

system in [38]. On the other hand, in the case of the "Man" image, the MSSIM of the proposed system is 6% higher than the system in [38].

V. CONCLUSION

In this paper, the color image blur detection has been integrated with compression and encryption as a 3 in 1 parallel HW architecture. The 128-bit AES-CBC has been combined with the modified Lorenz chaotic PRNG to provide a highly secure encryption scheme. To reduce the resources, the Haar DWT has been used as a common building block for both blur detection and compression. Furthermore, to achieve high speed, the entropy encoder has been implemented using the RLE technique. The proposed system has been implemented on FPGA (XC5VLX50T) using only 25% of the available slices. The system can process 4.5MP images at a rate of 30 FPS, which is more than 100% faster than all the available FPGA-based image compression-encryption systems. In addition, the system has been tested with all the standard 256×256 images. It is shown that depending on the amount of details in the image, the system can achieve 30dB PSNR at CRs in the range of (0.08-0.38). Furthermore, it is shown that the PSNR of the proposed system at 25% CR is higher than most of the newly published compression-encryption systems by more than 8%.

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