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Performance Comparison of QoS Deployment Strategies for Cellular Network Services

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ABSTRACT Differentiated quality-of-Service (QoS) techniques are widely used to distinguish between different service classes and prioritize service needs in mobile networks. Mobile Network Operators (MNOs) utilize QoS techniques to develop strategies that are supported by the mobile network infrastructure. However, QoS deployment strategy can ensure that the radio resources provided by the base station are easily consumed if it is not used correctly or when different techniques are used all together. In this paper, we propose a scheduling algorithm and compare two different QoS deployment strategies for prioritized User Equipment (UEs) (with higher scheduling rates and dedicated bandwidth) that MNOs can use in the current infrastructure, using a commercial real-time Long Term Evolution (LTE) network in different test scenarios. Moreover, we expose the real-time user experience in terms of uplink throughput and analyze results of the UE's real-time key performance indicators (KPIs) in detail. Experiment results are evaluated considering the implications of different QoS support types on network coverage and capacity planning optimizations. Our results demonstrate that even though pre-configured resource allocations can be given to prioritize UEs, the experience of all the UEs can be affected unexpectedly in the presence of many UEs who have received different QoS deployment support. Our experimental observations have revealed that location of the UEs with respect to Base Station (BS) and the availability of dedicated bandwidth UEs inside cell may have implications on the apriori defined resource allocation strategies of the other UEs.

INDEX TERMS QoS, differentiated services, mobile networks, experiments.

I. INTRODUCTION

In mobile networks, quality-of-service (QoS) refers to the capability of providing better service for the selected traffic types of a network under the same underlying technology. One of the main reasons of using QoS techniques for Mobile Network Operators (MNOs) is to detect and differentiate the prioritized services. Therefore, preferential services such as real-time applications can have higher priority over other services in MNO infrastructure. The emerging wireless applications in 5G require ultra-low latencies with ultra-high reliability [1]. Some users that demand to differentiate for the same type of services creates the problem. Inevitably, this will create an opportunity for extra income for the service provider [2]. Although 5G architecture offers promising key

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technologies to revolutionize the spectrum usage and energy efficiency, QoS support makes great sense for the more efficient use of 5G radio resources [3]. QoS allows the user equipment (UE) services to run according to their importance where the most critical services can be served first. Inside the infrastructure, QoS identifiers are used as marks that define the tolerable packet loss rate, packet delay budget, etc. of a mobile network service. The mobile network is aware of the services that are marked with QoS differentiation and handles them differently in accordance with the QoS deployment strategy of the MNO. This differentiation can only be done for the services used by certain UEs that are pre-configured in Core Network (CN), transport network and radio access network (RAN).

In QoS deployment, MNOs simply assign a priority level to the type of traffic that is to be managed and then specify how the RAN will behave for these different types of traffic

based on [QoS](#page-0-0) deployment strategy. In this case, the main focus of [MNOs](#page-0-0) is to select the most suitable techniques to manage the underlying network resources efficiently. Thus, there is a trade-off between simplifying the complexity of network and providing the best [QoS](#page-0-0) support to the [UE](#page-0-0) services. An accurate [QoS](#page-0-0) deployment strategy should prevent unequal distribution of mobile network's resources to [UEs](#page-0-0) while satisfying the requirements.

In this paper, we discuss two main [QoS](#page-0-0) deployment strategies that can be performed in the nodes of a real mobile network. The first one is to provide more scheduling rates in the [RAN](#page-0-0) side for the [UE](#page-0-0) services based on their [QoS](#page-0-0) values. This method is based on assigning a different and higher [QoS](#page-0-0) value to the [UE](#page-0-0) services demanded by the prioritized [UEs](#page-0-0) so that this [QoS](#page-0-0) value can be scheduled at higher rates in the [Media Access Control \(MAC\)](#page-0-0) scheduler of the [Base Station \(BS\).](#page-0-0) [UEs](#page-0-0) differentiated with this method are priority [UEs.](#page-0-0) The second one is to run the [UE](#page-0-0) services with a guaranteed bit rate so that their resources do not fall below a certain bandwidth requirement. The [QoS](#page-0-0) requirement in this case is to control the amount of allocated bandwidth for each [UE](#page-0-0) service and allow them to consume bandwidth based on their service requirements and mission. In this paper, we concentrate on experimental validations of these different [QoS](#page-0-0) deployment strategies to analyze the change behaviour of [UE Key Performance Indicators \(KPIs\)](#page-0-0) when diverse set of [QoS](#page-0-0) requirements co-exist under the same network cell.

The rest of the paper is organized as follows: Section [II](#page-1-0) gives the related works and main contributions of the paper. In Section [III,](#page-2-0) we detail the system model, concepts and different [QoS](#page-0-0) deployment scenarios used throughout the text. In Section [IV,](#page-7-0) we give the details of the experimental components and our experimental results for existence of different [QoS](#page-0-0) types [UEs](#page-0-0) inside the cell coverage as well as point out some of the main outcomes of the conducted experimental scenarios. Finally, in Section [V](#page-13-0) we give the conclusions of the paper.

II. RELATED WORK AND MOTIVATION

We review the state-of-the-art works in three main parts.

A. QoS SUPPORT

There are various works that investigate [QoS](#page-0-0) support for different mobile network deployment scenarios. For [QoS](#page-0-0) [Long Term Evolution \(LTE\)](#page-0-0) networks, the article in [4] proposes an analytical [QoS](#page-0-0) model in terms of context load, processing load and memory access rate in various elements. [QoS](#page-0-0) provisioning solutions are presented in [8], [9] and the [QoS](#page-0-0) needs of critical communications are detailed in [10] for 5G networks. The study in [14] proposes a resource allocation scheme with content caching in [Software-Defined](#page-0-0) [Networking \(SDN\)](#page-0-0) based networks. For [Network Function](#page-0-0) [Virtualization \(NFV\)](#page-0-0) cases, the authors in [15] study the placement of [Virtual Network Functions \(VNFs\)](#page-0-0) to provide better [QoS](#page-0-0) for [MNOs.](#page-0-0) The article in [16] investigates load balancing solutions to optimize the [QoS](#page-0-0) of a cloud radio

access network. The authors in [17] study user grouping strategies while considering the diverse [QoS](#page-0-0) requirements of users in [Non-Orthogonal Multiple Access \(NOMA\)](#page-0-0) systems. A game theoretic approach for interference-limited cellular environments is presented in [11]. The [QoS](#page-0-0) support in [LTE](#page-0-0) [heterogeneous network \(HetNet\)](#page-0-0) is investigated in [5], [6]. In [12], a framework for implementing a [QoS-](#page-0-0)aware energy and jitter efficient scheduling methodologies in downlink for [HetNet](#page-0-0) is developed. For [HetNet,](#page-0-0) the authors in [13] exploit the network cooperation and propose two joint radio resource management schemes with energy savings while satisfying the system [QoS](#page-0-0) performance. A cluster-based resource allocation scheme is studied in [7] to resolve the resource allocation problem with [QoS](#page-0-0) guarantees for ultra dense networks. All these studies have specialized focus on how [QoS](#page-0-0) structures will take place in the future. However, sufficient practical information on how [MNOs](#page-0-0) will follow a strategy in new generation mobile networks is still missing in those works.

[The 3rd Generation Partnership Project \(3GPP\)](#page-0-0) standardization has also defined several [UE](#page-0-0) categories for [LTE](#page-0-0) [20]. It is basically defined to categorize both uplink and downlink capabilities of [UE.](#page-0-0) Depending on the [UE](#page-0-0) capability, [BSs](#page-0-0) can connect to each [UE](#page-0-0) more effectively. In [LTE](#page-0-0) standardization and networks, [QoS](#page-0-0) between [UE](#page-0-0) and [packet data](#page-0-0) [network \(PDN\)](#page-0-0) gateway is applied using ''bearers'' which represent a set of network configurations so that prioritization of the traffic is handled based on the desired level of [QoS](#page-0-0) guarantees. Moreover, in [LTE](#page-0-0) networks [QoS Class](#page-0-0) [Identifiers \(QCI\)](#page-0-0) is defined that consists of basic classes which are classified as ''default'', ''expedited forwarding'' and ''assured forwarding''. For example, [QCI](#page-0-0) for [Guaranteed](#page-0-0) [Bit Rate \(GBR\)](#page-0-0) bearers is between 5 and 9 and for non[-GBR](#page-0-0) bearers is between 1 and 4. [QCIs](#page-0-0) simply deal with [UEs](#page-0-0) that are requesting different services so that the [LTE](#page-0-0) schedulers can set the priorities among them. The [UE](#page-0-0) categories are designed for specific use cases. In 5G networks, QoS model is based on QoS Flows and a standardized [5G QoS Identi](#page-0-0)[fier \(5QI\)](#page-0-0) to [QoS](#page-0-0) characteristics mapping is given in 3GPP Release 15 [31]. The complexity of [UE](#page-0-0) categories are also defined using [Transmission Mode \(TM\)](#page-0-0) in [3GPP](#page-0-0) [21]. For example in 5G networks, in general low numbered [UE](#page-0-0) categories are designed for [massive Machine-Type of Commu](#page-0-0)[nications \(mMTC\)](#page-0-0) use cases, whereas high numbered [UE](#page-0-0) categories are especially designed for considering [Enhanced](#page-0-0) [Mobile Broadband \(eMBB\)](#page-0-0) or [Ultra-Reliable Low-Latency](#page-0-0) [Communication \(URLLC\)](#page-0-0) use cases. In a [mMTC](#page-0-0) scenario, massive number of [UEs](#page-0-0) with NB-IoT capabilities accessing cellular network are simulated together using a [QoS-](#page-0-0)aware priority-based scheduling strategy in [18]. Based on [3GPP](#page-0-0) [QoS](#page-0-0) rule, an algorithm design that prioritizes [GBR](#page-0-0) and non-[GBR](#page-0-0) bearers of difference [QCIs](#page-0-0) is designed in [19]. Different from the traditional [UEs](#page-0-0) that access the network, our study provides a hint of the [QoS](#page-0-0) deployment strategy for [MNOs](#page-0-0) and how the management of [UEs](#page-0-0) with differentiated [QoS](#page-0-0) services can be accomplished for next generation mobile services

with core network configuration assistance for different user types.

B. SCHEDULER DESIGN

The problem of [QoS](#page-0-0) support using various scheduling algorithms has been proposed for cellular networks with different objectives, such as throughput, latency, fairness, energy, etc in [22]–[24]. A throughput maximizing method using max-weight based scheduling algorithm is proposed in [23]. Out of different studied scheduling algorithms, [Proportional](#page-0-0) [Fair \(PF\)](#page-0-0) scheduling algorithm has attracted higher interest among the academic as well as industrial community and is widely adopted. The authors in [22] have extended [PF](#page-0-0) scheduling to assign higher priority to [Resource Blocks \(RBs\)](#page-0-0) that yield above average spectral efficiency. The authors in [24] have compared both [QoS-](#page-0-0)aware and [QoS-](#page-0-0)non-aware scheduling algorithms for multi-tenant [SDN-](#page-0-0)based infrastructure of cellular networks. However, these studies have mostly focused on either simulations or theoretical works of demonstrating the benefits of scheduling to maximize a given objective under given constraints. In this work, we formulate the [QoS](#page-0-0) deployment in [MNO](#page-0-0) environment using an optimization problem. In addition, we also propose a scheduling algorithm as a solution to the defined optimization problem when there are various kinds of differentiated [UEs](#page-0-0) with [QoS](#page-0-0) prioritization. We have also investigated the existence of dedicated bandwidth [UEs](#page-0-0) in the network environment to observe the [QoS](#page-0-0) behaviour of all types of [UEs.](#page-0-0)

C. EXPERIMENTAL TRIALS

There are also various works on the experimental performance analysis of [QoS](#page-0-0) in the literature. In [27], the authors study the [QoS](#page-0-0) performance of [LTE](#page-0-0) networks under different load scenarios. In [25], the focus is to analyze the correlation between [UE](#page-0-0) position, network load and [QoS](#page-0-0) performances for video specific services. For 5G case, a heterogeneous [QoS-](#page-0-0)driven resource allocation policy for mmWave in massive [Multiple Input Multiple Output \(MIMO\)](#page-0-0) is presented in [28]. Experimental evaluation of a utility based decision-making approach for wireless mobile broadband networks is investigated in [26]. In [18], the authors aim to analyze the performance of the [QoS](#page-0-0) aware [Narrow Band](#page-0-0) [IoT \(NB-IoT\)](#page-0-0) networks. [QoS](#page-0-0) changes with different physical configurations can also be provided and an example case is the electrical tilt [29]. However, different from those studies, in this paper we evaluated the [QoS](#page-0-0) deployment methods that can be implemented in [MNOs](#page-0-0) infrastructure and showed what [QoS](#page-0-0) strategy can be selected in a real network when commercial [UEs](#page-0-0) exist rather than focusing on validations based on simulation environments. In our previous work [30], we have provided an experimental work that enables differentiated [QoS](#page-0-0) for different types of [LTE](#page-0-0) users. In this paper, we extend this analysis by adding dedicated bandwidth [UEs](#page-0-0) into the experimental setup to observe the end-to-end [QoS](#page-0-0) change with different types of [UEs,](#page-0-0) namely, normal, priority [UEs](#page-0-0)

bandwidth.

with higher scheduling rates and priority [UEs](#page-0-0) with dedicated

D. OUR CONTRIBUTIONS

The motivation of this paper stems from the fact that most of the literature work mentioned above does not observe the co-existence of [UEs](#page-0-0) with different set of [QoS](#page-0-0) requirements in real operational network. As a matter of fact, there are various efforts to quantify the [QoS](#page-0-0) improvements of various schedulers on [LTE](#page-0-0) users. For example in [LTE](#page-0-0) standardization efforts, differentiated [QoS](#page-0-0) has been considered and 9 standardized [QCIs](#page-0-0) have also been characterized (and 12 more [QCIs](#page-0-0) are added for 5G networks [10], [31]). However, none of the previous works have observed the effects of scheduling and priority weighting on the performance of priority and normal [UEs](#page-0-0) active throughput when real-users in a real network operation scenario is activated and deployed over [MNOs](#page-0-0) infrastructure. Moreover, most of the previously available research works have concentrated on validations via simulations but not using real world experimental trial. Our contributions in this paper can be summarized as follows: (i) We build a real-world test network environment to observe the end-to-end [KPI](#page-0-0) performance values. The experiments were run with three different [UE](#page-0-0) types (namely normal as well as two prioritized [UEs](#page-0-0) with dedicated bandwidth requirements and higher scheduling rates) that are created inside the [LTE](#page-0-0) network infrastructure. (ii) We detail some of the characteristics, limitations and benefits of utilizing prioritized [UEs](#page-0-0) with dedicated bandwidth requirements and higher scheduling rates inside cellular infrastructure. (iii) We have shown experimentally that deploying proposed [QoS](#page-0-0) strategies for diverse set of [UEs](#page-0-0) requires careful network capacity and coverage planning, which are detailed in discussions and main takeaways section of the paper. As a summary, Table [1](#page-3-0) provides a summary comparison between various techniques discussed above and the proposed approach in this paper.

Notations: Throughout the paper, bold letters represent vectors, i.e., **x** is a vector, and its i-th element is denoted by *xⁱ* . The sets are denoted by upper case calligraphic symbols. 0_M is the all-zeros column vector of size M. $||x||_1$ denotes the L1 norm of vector *x*.

III. SYSTEM MODEL AND CONCEPTS

Fig. [1](#page-3-1) shows a high level general diagram of the considered scenario where there are normal and prioritized [UEs](#page-0-0) distributed around the cells which are connected to core network to provide connectivity services to all [UEs](#page-0-0) inside the coverage area of a cell. In Table [2,](#page-4-0) the key notations symbols used throughout the paper are summarized. In our considered experimental scenario, we assume that $\mathcal{N} = \{1, 2, ..., N\}$ represents the set of multiple [UEs](#page-0-0) with diverse set of [QoS](#page-0-0) requirements with *N* [UEs](#page-0-0) in a given cell. Denote $\mathcal{N}_p \subset \mathcal{N}$ as the set of prioritized [UEs.](#page-0-0) Moreover, denote $\mathcal{N}_p^d \subset \mathcal{N}$ as the set of prioritized [UEs](#page-0-0) with dedicated bandwidth requirements of K_f [RBs](#page-0-0) with N_p^d prioritized [UEs](#page-0-0) of this type using the carrier frequency $f, \mathcal{N}_p^s \subset \mathcal{N}$ as the set of prioritized [UEs](#page-0-0)

FIGURE 1. Cellular network with eNodeBs providing services for normal, priority and dedicated UEs.

with N_p^s prioritized [UEs](#page-0-0) having *k* times higher scheduling rate than normal [UEs](#page-0-0) and $\mathcal{N}_n \subset \mathcal{N}$ as the set of normal UEs with *N_n* normal [UEs](#page-0-0) as shown in Fig. [1.](#page-3-1) Note that $\mathcal{N}_p = \mathcal{N}_p^s \cup \mathcal{N}_p^d$ and $\mathcal{N} = \mathcal{N}_p \cup \mathcal{N}_n$. Let $\mathcal{T} = \{1, 2, ..., T\}$ denote the set of the observation time where T is the duration of observation. Moreover, we also would like to point out that the analyzed problem considers a single shared [Serving Gateway \(S-GW\)](#page-0-0) resource serving multiple connected [BSs](#page-0-0) as shown in Fig. [1](#page-3-1) as well.

We assume that there are *M* available [RBs](#page-0-0) in a given cell at $t \in \mathcal{T}$. We denote the [RB](#page-0-0) set $\mathcal{M} = \{1, 2, ..., M\}$. Let the binary variable $q_{m,n}^t$ indicate whether [UE](#page-0-0) $n \in \mathcal{N}$ is assigned to [RB](#page-0-0) $m \in \mathcal{M}$ or not (i.e., if *n*-th [UE](#page-0-0) is assigned to *m*-th [RB,](#page-0-0) then $q_{m,n}^t = 1$ else $q_{m,n}^t = 0$ at $t \in \mathcal{T}$ i.e. during [Transmission Time Interval \(TTI\)](#page-0-0) (usually equal to 1 millisecond). Hence, each [UE](#page-0-0) $n \in \mathcal{N}$ is assigned to only one [RB](#page-0-0) so that

$$
\sum_{m \in \mathcal{M}} q_{m,n}^t = 1, \quad \forall n \in \mathcal{N}.
$$
 (1)

We define $\mathbf{\Delta}^{\mathbf{t}} = [\mathbf{\Psi}_1^t \ \mathbf{\Psi}_2^t \ \dots \ \mathbf{\Psi}_n^t] = (\mathbf{\Psi}_n^t, \mathbf{\Psi}_{-n}^t)$ as the $M \times N$ [UE](#page-0-0) assignment matrix of all [RBs](#page-0-0) at $t \in \mathcal{T}$. Here $\Psi_n^t = [q_{1,n}^t \quad q_{2,n}^t \quad \dots \quad q_{M,n}^t]^T$ denotes *n*-th [UE'](#page-0-0)s $M \times 1$ [RB](#page-0-0) assignment vector and Ψ^t_{-n} denotes the assignment vector of all [UEs](#page-0-0) other than the *n*-th [UE.](#page-0-0) Denote

 $\mathbf{S}_n^t = [s_{1,n}^t, s_{2,n}^t, \dots, s_{M,n}^t]^T$ as $M \times 1$ vector of the obtained [Modulation Coding Scheme \(MCS\)](#page-0-0) index values $s_{m,n}^t$ of *n*-th [UE](#page-0-0) and *m*-th [RB](#page-0-0) at $t \in \mathcal{T}$. Let us also denote \mathbf{R}_n^t = $[r_{1,n}^t, r_{2,n}^t, \ldots, r_{M,n}^t]^T$ as $M \times 1$ vector of the obtained [Trans](#page-0-0)[port Block Size \(TBS\)](#page-0-0) values $r_{m,n}^t$ of *n*-th [UE](#page-0-0) and *m*-th [RB](#page-0-0) at $t \in \mathcal{T}$. Hence, $\mathbf{R}^t = [\mathbf{R}_1^t \ \mathbf{R}_2^t \ \dots \ \mathbf{R}_n^t]$ and $\mathbf{S}^t =$ $[\mathbf{S}_1^t \ \mathbf{S}_2^t \ \ldots \ \mathbf{S}_n^t]$ are the *M* × *N* matrix of [TBS](#page-0-0) and [MCS](#page-0-0) index values at $t \in \mathcal{T}$ respectively. Define $\Upsilon_n^t = (\Psi_n^t)^T \times \mathbf{R}_n^t$ as the achieved data rate at $t \in \mathcal{T}$ for UE $n \in \mathcal{N}$.

Note that [UEs](#page-0-0) exchange information with its corresponding eNodeB in a particular region with an assigned [TBS](#page-0-0) value in a given [TTI.](#page-0-0) The maximum value of assigned [TBS](#page-0-0) for a given eNodeB for each [UE](#page-0-0) is identified by an integer value of α in this paper. For this reason, each [UE](#page-0-0) can get at most [TBS](#page-0-0) value of α .

A. PROBLEM FORMULATION

The problem definition can be described as follows: Given a network state $S = (\Psi_n^t, \Psi_{-n}^t)$ where (Ψ_n^t, Ψ_{-n}^t) is a combination of each RB assignment in the set of M to each UEs in the set N , we look for the optimal values of assignments to minimize a cost function $f(\Psi_n^t, \Psi_{-n}^t)$:

$$
f(\Psi_n^t, \Psi_{-n}^t) = -\sum_{n \in \mathcal{N}} U_n^t,\tag{2}
$$

TABLE 2. Symbols used throughout the paper.

Symbol	Meaning
$N, {\cal N}$	Number of UEs, UE set
$N_p^d,\, \mathcal{N}_p^d$	Number of prioritized UEs with dedicated bandwidth
	requirements of K_f RBs, corresponding prioritized UE set
N_p^s, N_p^s	Number of prioritized UEs with k times higher scheduling
	rate than normal UEs, corresponding prioritized UE set
$\frac{N_n, N_n}{M, M}$	Number of normal UEs, Normal UE set
	Number of RBs, RB set for all UEs
\mathcal{M}' , \mathcal{M}'' , \mathcal{M}'''	the set of RBs assigned to dedicated,
	priority and normal UEs respectively
$\overline{\Delta^{\text{t}}}$	$M \times N$ RRH assignment matix of all UEs over all RBs
$\mathbf{\Psi}_n^t$	$M \times 1$ RB assignment vector for <i>n</i> -th UE at $t \in \mathcal{T}$
$\boldsymbol{q}_{m,n}^{t}$	if m-th RB is assigned to n-th UE at TTI t 1: $=$
	0: else
\mathbf{R}_n^t	$M \times 1$ vector of the obtained TBS values
	of n -th UE in M different RBs
$r_{m,n}^t$	TBS value obtained at m-th RB for n-th UE.
\mathbf{S}_n^t	$M \times 1$ vector of the obtained MCS index values
	of n -th UE in M different RBs
$s_{m,n}^t$	MCS index value obtained at m-th RB for n-th UE.
α	Maximum achievable TBS value by each UE
$U_{n,m}^t$	utility metric obtained for the
	n -th UE using the m -th resource
λ_n^t	the average data rate of the n -th UE at time t.
$R_{m,n}^t$	obtained instantaneous data rate of the
	n -th UE at m -th resource and time t.
\boldsymbol{k}	the scheduling constant to provide
	higher data rate for UEs
τ	time constant of smooth filter
	in PF scheduler of LTE network
$\overline{\Delta t}$	allocation interval in PF scheduler of LTE network
K_f	maximum number of RBs allocated by prioritized
	UEs with dedicated bandwidth requirements
	on carrier frequency f
Υ^t_n	the achieved data rate at $t \in \mathcal{T}$ for
	UE $n \in \mathcal{N}$ and is equal to $(\mathbf{\Psi}_n^t)^T \times \mathbf{R}_n^t$

where U_n^t is the utility of the *n*-th UE at $t \in \mathcal{T}$. In order to accomplish this, each UE's utility needs to be maximized by choosing appropriate RB assignments. Using assigned [TBSs](#page-0-0) as the maximization parameter, the utility function of *n*-th UE is expressed as follows:

$$
U_n^t = \sum_{m \in \mathcal{M}} U_{m,n}^t = \sum_{m \in \mathcal{M}} \left(q_{m,n}^t \times r_{m,n}^t \right),
$$

$$
\forall n \in \mathcal{N}, \quad \forall t \in \mathcal{T}, \quad (3)
$$

where the term $U_{m,n}^t = q_{m,n}^t \times r_{m,n}^t$ is the obtained [TBS](#page-0-0) value of the [UE](#page-0-0) $n \in \mathcal{N}$ for [RB](#page-0-0) $m \in \mathcal{M}$ at $t \in \mathcal{T}$ ^{[1](#page-4-1)}. Then, the optimization problem can be described as follows: Our goal is to maximize the sum of total [TBS](#page-0-0) utility of all [UEs](#page-0-0) with the decision variables: (i) *Assignment problem:* the assignment of UE to each RB is represented by the variables $q_{m,n}^t$. (ii) *[UEs](#page-0-0) types [QoS](#page-0-0) satisfaction problem:* the resource allocation between two different prioritized [UEs](#page-0-0) and normal [UEs](#page-0-0) are characterized by Υ_n^t and dedicated BW UEs by $||\Psi_n^t||_1.$

For the considered shared mobile architecture, we use the following formulation for our optimization problem:

$$
\underset{\Delta^t}{\text{minimize}} \quad f(\Psi_n^t, \Psi_{-n}^t) \tag{4}
$$

subject to
$$
\sum_{m \in \mathcal{M}} q_{m,n}^t = 1, \ \forall n \in \mathcal{N}, \forall t \in \mathcal{T},
$$
 (4a)

$$
\{q_{m,n}^t\} \in \{0, 1\}, \quad \forall m \in \mathcal{M}, \forall n \in \mathcal{N}, \qquad (4b)
$$

$$
||\Psi_n^t||_1 = K_f, \quad \forall n \in \mathcal{N}_p^d,
$$
 (4c)

$$
\Upsilon_i^t = k \times \Upsilon_j^t,
$$

$$
\forall i \in \mathcal{N}_p^s, \forall j \in \mathcal{N}_n, \forall t \in \mathcal{T}, \qquad (4d)
$$

$$
r_{m,n}^t \le \alpha, \quad \forall n \in \mathcal{M}, \forall n \in \mathcal{N}, \forall t \in \mathcal{T}. \quad (4e)
$$

where the objective function given in [\(4\)](#page-4-2) is to maximize the total sum of [TBS](#page-0-0) values over all [UEs,](#page-0-0) [RBs](#page-0-0) and observation duration $t \in \mathcal{T}$. Constraint in [\(4a](#page-4-2)) illustrates that each UE is assigned to only one RB, [\(4b](#page-4-2)) denotes the binary value constraint of [UE](#page-0-0) assignment per [RB,](#page-0-0) [\(4c](#page-4-2)) gives the dedicated bandwidth requirements of prioritized UEs at carrier frequency f , [\(4d](#page-4-2)) gives the constraint imposed for prioritized UEs with *k* times higher scheduling rate than normal [UEs](#page-0-0) and [\(4e](#page-4-2)) yields the maximum achievable [TBS](#page-0-0) value by each [UE](#page-0-0) at a given [TTI](#page-0-0) $t \in \mathcal{T}$.

The time scale of the operation of the maximization operate at faster scale at the BSs, than the various gateways at the core or transport networks due to dynamic nature of the propagation environment. The optimization problem can be solved for each [TTI](#page-0-0) in a given time frame depending on the use cases and the requirements. Note that in the optimization problem defined above, the main decision process at the scheduling layer is to decide the number of [RBs](#page-0-0) assigned to each [UEs](#page-0-0) based on their time-frequency locations, [MCS](#page-0-0) and transmission power. In our optimization problem, we are dealing with a binary optimization problem. The optimization variables are binary, hence this problem is an [Integer Linear](#page-0-0) [Problem \(ILP\)](#page-0-0) [33] (and also c.f. Problem (P1) in [34]). However, [ILPs](#page-0-0) are difficult to solve and known to be NP-hard and with exponential execution time, preventing solutions even for reasonable problem sizes. On the other hand, there are also several approximations (e.g. continuous methods) for binary optimization in the literature [35]. One approach is to recognize that after relaxation, the problem is a linear program (LP) and enforce the binary constraints after solving the LP.

Moreover, solving [\(4\)](#page-4-2) is challenging because (i) of the existence of coupling behaviour in [UE](#page-0-0) assignments for each [RB](#page-0-0) problem and large-scale [QoS](#page-0-0) satisfaction requirements problem for different [UE](#page-0-0) types and carrier frequency, (ii) the achievable maximum [TBS](#page-0-0) value per [TTI](#page-0-0) depends on many factors such as wireless link, channel bandwidth, number of [UEs,](#page-0-0) locations of [UEs,](#page-0-0) interference, that may not be controllable, (iii) the globally optimal [UE](#page-0-0) assignment per [RB](#page-0-0) decision solution for given demands depends on [QoS](#page-0-0) requirements of multiple number and types of [UEs.](#page-0-0) This is

¹Note that [TBS](#page-0-0) values are extracted from a table mapping obtained using [MCS](#page-0-0) index and number of [RBs](#page-0-0) values according to 3GPP specification to determine how many bits can be transmitted per one [TTI](#page-0-0) for [Physical](#page-0-0) [Downlink Shared Channel \(PDSCH\)](#page-0-0) [32]. The [TBS](#page-0-0) index, $i \in \{0, \ldots, 26\}$, is a function of modulation and coding scheme as given in Table 8.6.1-1 of [32]

non-tractable in large-scale and has high computational complexity.

Therefore, in the following section, we will discuss a heuristic approach to the optimization problem given in [\(4\)](#page-4-2). In our scheduling design methodology, [UEs-RBs](#page-0-0) assignments are done while providing the necessary [QoS](#page-0-0) guarantees for different [UE](#page-0-0) types. This scheduler design is also used throughout our experiment trial to observe its implications in real-life scenarios.

B. HEURISTIC SCHEDULER DESIGN

To find a solution for the optimization problem defined in [\(4\)](#page-4-2), we study a scheduler design and propose an algorithm in this section. Schedulers are one of the core components of LTE systems utilized in eNodeBs for resource management among [UEs](#page-0-0) and network performance optimization. In its basic functionality, schedulers allocate resources to [UEs](#page-0-0) based on their [Channel Quality Indicator \(CQI\)](#page-0-0) and [QoS](#page-0-0) requirements which can be defined by [MNOs.](#page-0-0)

Some key challenges of providing end-to-end [QoS](#page-0-0) support for [LTE](#page-0-0) users are: First, to map different [QoS](#page-0-0) classes across different domains (such as [RAN,](#page-0-0) transport and core networks). Second, to provide the appropriate scheduling methodologies that can enable [QoS](#page-0-0) differentiation among users. In eNodeBs of 4G systems, one of the main scheduler methodology is using [PF](#page-0-0) scheduling algorithm. It provides a balance between fairness and overall spectrum efficiency. The performance metric of [PF](#page-0-0) algorithm for the *n*-th [UE](#page-0-0) can be written as

$$
R_{m,n}^t = \frac{U_{m,n}^t}{\lambda_n^t},\tag{5}
$$

where $U_{m,n}^t = q_{m,n}^t \times r_{m,n}^t$ is the instantaneous achievable data rate (or [TBS\)](#page-0-0) of the *n*-th user, at m-th [RB](#page-0-0) and time $t \in \mathcal{T}$, λ_n^t denotes the average data rate of the n^{th} [UE](#page-0-0) until time $t \in \mathcal{T}$ and it can be calculated by

$$
\lambda_n^t = \left(1 - \frac{1}{\tau}\right) \times \lambda_n^{(t - \Delta t)} + \sum_{\forall m \in \mathcal{M}} \frac{q_{m,n}^{(t - \Delta t)} \times U_{m,n}^{(t - \Delta t)}}{\tau}, \quad (6)
$$

where $\tau > 1$ denotes the time constant of smooth filter which controls the system latency and is the past window length, Δt is the [TTI,](#page-0-0) which is the period of allocation. Note that window length value τ gives a trade-off between throughput and latency. Higher window value results in higher throughput since the scheduler waits for the large peaks. This in turn increases the latency. Lower window value indicates low waiting period for throughput peaks, whereas it decreases the latency [36]. To solve the above optimization problem in an experimental set-up, we perform configuration updates on the existing scheduler methods of eNodeBs. In our experimental trials, we have used the following utility metric,

$$
\bar{R}_{m,n}^t = w_n^t R_{m,n}^t,\tag{7}
$$

where w_n^t is the weight assigned to each user $n \in \mathcal{N}$ based on their priority status. For prioritized UEs with *k* times

higher scheduling rate than normal UEs $w_i^t / w_j^t = k$, $\forall i \in$ $\mathcal{N}_p^s, \forall j \in \mathcal{N}_n, \forall t \in \mathcal{T}$. Without loss of generality, we assign $w_i^f = 1$, $\forall i \in \mathcal{N}_n$. Hence during experimental trials, the utility metric in [\(7\)](#page-5-0) becomes,

$$
\bar{R}_{m,n}^t = \begin{cases} R_{m,n}^t, & \text{if } n \in \mathcal{N}_n, \mathcal{N}_p^d \\ k \times R_{m,n}^t, & \text{if } n \in \mathcal{N}_p^s \end{cases}
$$
(8)

Algorithm 1 Pseudo Code of the Heuristic Scheduling Algorithm Used in Experimental Setup

 $\mathbf{Input:}\ \mathbf{S^t},\ k,\ \mathcal{M},\ \mathcal{N},\ K_f$ **Output:** [RBs-UEs](#page-0-0) $M \times N$ assignment matrix, Δ^t at $t \in \mathcal{T}$ **Initialization:** $\mathcal{M}' = \emptyset$, $\mathcal{M}'' = \emptyset$, $\mathcal{M}''' = \emptyset$, $\Psi_n^t = \mathbf{0}_M$. 1: **procedure** SCHEDULE 2: **Foreach** (UE-n $\in \mathcal{N}$) // Calculate RB assignment vector for n // 3: **Compute:** $U_{m,n}^t$ and λ_n^t , $\forall m \in \mathcal{M}$ using \mathbf{S}^t 4: **if** $n \in \mathcal{N}_p^d$ \triangleright Iterate until min. bandwidth 5: **while** $(|\mathcal{M}'| < K_f)$ do // Allocate free RBs for Dedicated UEs // 6: **Run:** $(\Psi_n^t, FLAG) =$ $\overline{ALLOCATE}(\mathcal{M}',\mathcal{M},\mathcal{M}'',\mathcal{M}''')$ 7: **if** !FLAG **then** 8: **break** // Requirement satisfied // 9: **end if** // Reallocate RBs of Normal UEs // 10: **Run:** $(\Psi_n^t, FLAG) =$ $\stackrel{\cdot\cdot}{REALLOCATE}(\mathcal{M}',\mathcal{M}''')$ 11: **if** !FLAG **then** 12: **break** // Requirement satisfied // 13: **end if** // Reallocate RBs of Priority UEs // 14: **Run:** $(\Psi_n^t, FLAG) =$ $\stackrel{\cdot\cdot}{REALLOCATE}(\mathcal{M}',\mathcal{M}'')$ 15: **if** !FLAG **then** 16: **break** // Requirement satisfied // 17: **else** //Share RBs between dedicated UEs// 18: **Run:** $\Psi_n^t = RB_ALLOCALOCATION$ $(\mathcal{M}', \mathcal{U}_{m,n}^t, \lambda_n^t, k, \mathcal{N})$ 19: **break** // Requirement satisfied // 20: **end if** 21: **end while** 22: **else** //Calculate RB assignment vector for// //normal and priority UEs// 23: **Run:** $\Psi_n^t =$ $RB_ALLOCALION$ $(\mathcal{M}, U_{m,n}^t, \lambda_n^t, k, \mathcal{N})$ 24: **end if** 25: **end foreach** 26: **end procedure**

Algorithm [1](#page-5-1) and corresponding function calls in Algorithm [2](#page-6-0) summarizes the pseudo code of the [RB](#page-0-0) allocation strategy for each [UE](#page-0-0) with different [QoS](#page-0-0) requirements. Note that in line #18 of Algorithm [1,](#page-5-1) dedicated BW [UEs](#page-0-0) uses [\(8\)](#page-5-2) similar to normal [UEs](#page-0-0) after obtaining its required RBs.

, *k*, N) 26: **Find:** (*m* , *n* $=$ arg max_{*m*∈*M*,*n*∈*N*} { $\overline{\overline{R}}$ ^{*tt*}_{*m*,*n*}} using [\(8\)](#page-5-2) 27: **Update:** Ψ_n^t using m^* 28: **Return:** (Ψ_n^t, FLAG) 29: **end procedure**

In summary, we propose a simple heuristic scheme to solving the optimization problem at hand with low computational complexity that can work on the desired time scale. The algorithm is simple, but an essentially greedy scheme (not necessarily the optimal one). It is based on time-level optimization. Each [RB](#page-0-0) can only be allocated to a single user in any [TTI.](#page-0-0) At any given $t \in \mathcal{T}$, the first priority is to assign the available resources to priority UEs (i.e. dedicated BW UEs) in order to satisfy the minimum throughput requirements of dedicated BW priority UEs. No extra RBs are given to dedicated BW UEs. When the number of RBs were enough for dedicated BW UEs, the remaining RBs would be used to maximize the throughput of remaining priority UEs and normal UEs. After this requirement is satisfied, we try to maximize the remaining priority UEs' (with higher scheduling rates) throughput or equivalently maximize the number of scheduled priority UEs in the current $t \in \mathcal{T}$. [\(8\)](#page-5-2) tries to enforce that if the priority UEs with higher scheduling rates are scheduled, they should transmit *k* times more data bits than the normal UEs. In the case of infeasibility, i.e. the problem does not lead into

any optimal solution (no resources available for additional priority users (i.e. dedicated BW UEs) at any time $t \in \mathcal{T}$), the heuristic solution serves as many dedicated BW UEs as possible, excluding the remaining UEs from RB allocations. Hence, the algorithm always converges even though it may not be feasible for some UEs.

C. COMPLEXITY ANALYSIS

The heuristic scheme presented in Algorithm [1](#page-5-1) is motivated by the complexity of the schemes in the optimal solution. With $M \times N$ variables, the complexity of the optimization in [\(4\)](#page-4-2) is $O(2^{M \times N})$. In Algorithm [1,](#page-5-1) the worst case would be when dedicated BW request cannot be satisfied with allocation of all RBs. So the number of RBs are not enough for the remaining UEs (i.e. priority users with higher scheduling rates and normal users). In this case, lines #5-21 in Algorithm [1,](#page-5-1) #2-10, and #14-22 in Algorithm [2](#page-6-0) are executed. When M and N are large, the complexity is due to the two while iterations (Line #5 in Algorithm [1](#page-5-1) and line #2 in Algorithm [2\)](#page-6-0) and also finding ''arg max'' (line #26 in Algorithm [2.](#page-6-0) The complexity for performing arg max is proportional to number of values being sorted. The overall time complexity of Algorithm [1](#page-5-1) with Algorithm [2](#page-6-0) is therefore $\mathcal{O}(M^2 \times N)$.

D. PRACTICAL SETTINGS AND CHARACTERISTICS OF QoS **DEPLOYMENTS**

In Fig. [1'](#page-3-1)s architecture, to provide end-to-end [QoS](#page-0-0) support for all types of [UEs,](#page-0-0) [RAN,](#page-0-0) transport and core network equipment need to be configured appropriately for each type of defined [UEs](#page-0-0) with a given set of [QoS](#page-0-0) requirements. End-to-end [QoS](#page-0-0) deployment strategy should be assured and managed by entities that have a global mobile network topology view. The basis of [QoS](#page-0-0) support in [LTE](#page-0-0) networks is accomplished via [Evolved Packet System \(EPS\)](#page-0-0) bearers [37]. An [EPS](#page-0-0) bearer builds a logical channel between [UE](#page-0-0) and [PDN.](#page-0-0) [EPS](#page-0-0) bearers can be classified as [GBR](#page-0-0) and non[-GBR](#page-0-0) bearers depending on scheduling and queue management policy. In [GBR](#page-0-0) bearers, a permanent network resource is allocated whereas in non-[GBR](#page-0-0) this does not exist. In our experimental scenarios, prioritized [UEs](#page-0-0) are created by defining different non[-GBR QCI](#page-0-0) profile at [Home Subscriber Server \(HSS\)](#page-0-0) where [QCI](#page-0-0) levels are assigned statically. This assigned [QCI](#page-0-0) gives higher priority to prioritized [UEs](#page-0-0) with high scheduling rate than normal [UEs.](#page-0-0) During operation in first step, [UE](#page-0-0) service requests [QoS](#page-0-0) value from the [CN](#page-0-0) via [Control Plane \(CP\)](#page-0-0) signaling during service initiation. Then, the [CN](#page-0-0) assigns the pre-defined [QoS](#page-0-0) value that will be used in the [User Plane \(UP\)](#page-0-0) session of the [UE](#page-0-0) service (such as voice, [eMBB\)](#page-0-0). Therefore, [QoS](#page-0-0) values for [UE](#page-0-0) services in [MNO](#page-0-0) environment are assigned by the [CN.](#page-0-0) [RAN](#page-0-0) equipment is just executing the [QoS](#page-0-0) policies (e.g. via scheduling) in radio access segment depending on the [QoS](#page-0-0) assignments done at [CN.](#page-0-0)

For our experimental scenarios, infrastructure provider plans to provide a dedicated wireless resource allocation for its customers (e.g. to a national bank that has

nationwide branches). The intended use case is to migrate [Automated Teller Machines \(ATMs\)](#page-0-0) from g.shdsl [38] which is an old technology using fixed line access into a wireless access device. For this reason, the required [Upload \(UL\)](#page-0-0) RB allocation for an [ATM](#page-0-0) is selected to be K_f . In this case, the location, received signal quality and number of the [ATMs](#page-0-0) inside the coverage cell are quite important so that [BSs](#page-0-0) can service to those [ATMs](#page-0-0) with fixed bandwidth wireless access while also having minimum impact on existing normal [UEs.](#page-0-0) For this reason before selecting the experimental fields and corresponding cells, prior studies on how many dedicated bandwidth [ATMs](#page-0-0) can be accommodated in a certain region need to be studied via appropriate network planning tools while considering the expertise of network planning experts. Other prioritized [UEs](#page-0-0) can have higher scheduling rates than normal [UEs.](#page-0-0) These [UEs](#page-0-0) are generally the public and private enterprise customers of [MNOs.](#page-0-0) To give insights into different [QoS](#page-0-0) deployment strategies used in our experimental set-up, we provide Table [3,](#page-7-1) which gives a summary of the characteristics, limitations and benefits of the experimented different [QoS](#page-0-0) deployment strategies, namely both dedicated and higher scheduling rate policies that provide cellular services to [UEs.](#page-0-0)

Note that in our analysis results given in the next section, we have not run simulations of the proposed heuristic scheduling algorithm to compare it with the optimal solution of the original optimization problem given in [\(8\)](#page-5-2). Our contributions are mainly focusing on experimental analysis and results as opposed to pure simulation-based analysis results.

Note that experimental works that are performed in real live networks of operators on differentiated QoS trials in LTE networks are not common in the literature and is not a trivial task especially when real users are using the existing operational infrastructure. Additionally, other practical restrictions such as the hardware/software limitations, regulatory restrictions (compliance requirements, customized tariffs to different users, etc), security, marketing demands (price per differentiated users, service usage needs, etc.) need to be taken into account in practical QoS deployments and algorithm design. Hence, results and related experiences (main takeaways discussions, lessons learned, trade-off analysis or challenges experienced) on field trials are quite valuable insights into investigation of the achievable performances of different QoS deployment strategies in parallel with UEs having different QoS policies under real conditions and with real equipment limitations.

FIGURE 2. The location of the BS, handover region and coverage areas where cell center and cell edge UEs' KPIs (for normal and both prioritized UEs) are observed.

IV. PERFORMANCE ANALYSIS

A. DETAILS OF THE EXPERIMENTAL SETUP

Fig. [2](#page-7-2) shows the schematic illustration of the network topology (the locations of LTE networks' cell-edge and cell-center test sites) used throughout the real network experiments in Cekmekoy region of city of Istanbul in Turkey. The experiments were run during different times in two days (ranging between 14:50 to 23:40 local time). Two feature enabled cells namely *PCI* − 209 and *PCI* − 378 are used for observations. Normally, for [Reference Signal Received Power \(RSRP\)](#page-0-0) values below −65 dBm [LTE](#page-0-0) users are considered to be at near distance locations to connected eNodeB. For ranges between −75 dBm and −85 dBm [LTE](#page-0-0) users are at middle distance locations to connected eNodeB (practical value is around -80 dBm or slightly better and this is also good for $f =$ 800 MHz (low bandwidth) conditions)). For ranges between −100 dBm and −120 dBm, [LTE](#page-0-0) users are considered to be in far (edge) distance to connected eNodeB (practical value

is −105 dBm or −110 dBm where noise limited conditions exist for $f = 1800$ MHz (high bandwidth)).

Our considered scenario for the experimental set-up is as follows: 11 monitored [UEs](#page-0-0) are connecting to eNodeB sequentially in which there are three types of [UEs](#page-0-0) with different [QoS](#page-0-0) requirements. These types are Normal [UEs,](#page-0-0) prioritized [UEs](#page-0-0) with dedicated bandwidth (shortly named as dedicated BW [UE\)](#page-0-0) and prioritized [UEs](#page-0-0) (shortly named as priority [UE\)](#page-0-0) with high scheduling rate. In addition 11 monitored [UEs,](#page-0-0) there are also normal commercial [UEs](#page-0-0) in the real network connected to this site whose throughput values are not monitored but their presence has direct effect on observed throughput of both priority, dedicated BW and normal [UEs.](#page-0-0)

During our experiments, we have configured prioritized [UEs](#page-0-0) to be $k = 1.5$ times higher scheduling rate than nor-mal [UEs.](#page-0-0) To create priority [UEs](#page-0-0) to be $k = 1.5$ times higher scheduling rate, a different non[-GBR QCI](#page-0-0) profile is defined in [HSS](#page-0-0) and those users are statically assigned to this [QCI](#page-0-0) level. This [QCI](#page-0-0) value has a higher priority than the [QCI](#page-0-0) value of normal [UEs.](#page-0-0) In [RAN,](#page-0-0) eNodeB is configured by applying a weight ratio of $w_i^t / w_j^t = 1.5$, $\forall i \in \mathcal{N}_p^s, \forall j \in \mathcal{N}_n, \forall t \in$ T resource allocation inside [PF](#page-0-0) scheduler when [E-UTRAN](#page-0-0) [Radio Access Bearer \(E-RAB\)](#page-0-0) is established. Therefore, resource allocations considering the [QCI](#page-0-0) values are also taken into account in eNodeB at the same time. As part of experimental [UE](#page-0-0) equipment, we have used 10 [QCI-](#page-0-0)6 SIM cards assigned as dedicated bandwidth [UEs,](#page-0-0) 10 [QCI-](#page-0-0)7 SIM cards assigned as priority [UEs,](#page-0-0) 1 [QCI-](#page-0-0)8 SIM card assigned as normal [UEs](#page-0-0) with 11 identical brand [UE](#page-0-0) terminals.

In all the experimental tests, eNodeBs are configured to operate in $f = 1800$ MHz carrier frequency with 20 Mhz bandwidth at cell centers and in $f = 800$ Mhz carrier frequency in cell edges with 10 MHz bandwidth. Therefore, for requirements of dedicated BW [UEs,](#page-0-0) K_f = 20 [RBs](#page-0-0) is selected for $f = 1800$ Mhz and $K_f = 10$ [RBs](#page-0-0) is selected $for $f = 800$ Mhz. During experiment, [UEs](#page-0-0) download data via$ FTP server as shown in Fig. [1.](#page-3-1) Then in all test scenarios, while normal [UEs](#page-0-0) are downloading data, $N_p = 10$ priority test UEs (depending on [UE](#page-0-0) type) enter into the cell area and download data simultaneously under the same connected eNodeB. After a certain amount of test period, priority [UEs](#page-0-0) quit the cell. Depending on the test scenario, those normal and priority [UEs](#page-0-0) may be on the cell center or cell edge. During our experiments, we have configured dedicated BW [UEs](#page-0-0) bandwidth requirements to be $K_f = 20$ RBs. All experiments are done in full-buffered traffic mode to force the scheduler of eNodeB work in full performance capacity. Hence, high-traffic areas are selected for experiments. Content size of 5 Gbytes for [UL](#page-0-0) are used via FTP for demonstrating the [UL](#page-0-0) throughput variations. In summary, the system level parameters used throughout the experiments are detailed in Table [4.](#page-8-0)

B. EXPERIMENTAL RESULTS

In this section, we present some of the experimental evaluation results to provide end-to-end [QoS](#page-0-0) support for [LTE UEs.](#page-0-0)

TABLE 4. Experimental parameters and their corresponding values.

In our experiments, we only show achieved experimental throughput results in [UL](#page-0-0) direction without loss of generality using the scheduling algorithm in Algorithm [1.](#page-5-1)

Fig. [3](#page-9-0) shows the [UL](#page-0-0) performance of [UEs](#page-0-0) in two different experimental scenarios to enable [QoS](#page-0-0) support in live [LTE](#page-0-0) networks. In Fig. [3a](#page-9-0), all monitored [UEs](#page-0-0) are in cell center whereas in Fig. [3b](#page-9-0) normal [UE](#page-0-0) is in cell center and $N_p^d = 10$ dedicated BW [UEs](#page-0-0) are in cell edge. In all of two scenarios, at first while normal [UE](#page-0-0) is generating traffic, 10 more dedicated BW [UEs](#page-0-0) are starting to generate traffic one-by-one sequentially. In Fig. [3a](#page-9-0), it is observed that all of dedicated BW [UEs'](#page-0-0) [UL](#page-0-0) throughput values stay above 4 Mbps until the arrival of sixth dedicated BW [UE](#page-0-0) (where no free [RBs](#page-0-0) can be obtained from normal [UEs](#page-0-0) since $K_f = 20$ [RBs\)](#page-0-0) whereas normal [UE](#page-0-0) throughput values diminish to zero value after fifth dedicated BW [UE](#page-0-0) enter into the coverage area. Note that dedicated BW [UEs](#page-0-0) get lower [RBs](#page-0-0) than minimum required [RB](#page-0-0) of $K_f = 20$ due to non-availability of free [RBs](#page-0-0) allocated by normal [UEs.](#page-0-0) For this reason, dedicated BW [UEs](#page-0-0) share [RBs](#page-0-0) with other dedicated BW [UEs](#page-0-0) as observed in Fig. [3a](#page-9-0). Most of the throughput values of all dedicated BW [UEs](#page-0-0) are slightly lower than 2.4 Mbps at the end of experimental observation period in Fig. [3a](#page-9-0) where each dedicated BW [UE](#page-0-0) obtains K_f = M/N_p^d = 10 [RBs](#page-0-0) since N_p^d = 10. In Fig. [3b](#page-9-0), we can observe that after all dedicated BW [UEs](#page-0-0) are put into cell edge, the amount of dedicated resources to normal [UEs](#page-0-0) diminishes fast. This is due to low signal quality dedicated BW [UEs](#page-0-0) entering into the coverage area of the cell and suppressing [RBs](#page-0-0) utilization of normal [UE.](#page-0-0) Normal [UEs'](#page-0-0) throughput value decreases after arrival of fifth dedicated BW [UE](#page-0-0) into the cell.

The mobility of users is one of the dimensions that can have an impact on the QoS of UEs. High mobility scenarios can make resource allocations more challenging and cause

FIGURE 3. UL performance of users (a) All users are in cell center (both $f = 1800$ MHz) (b) Normal user is in cell center $(f = 1800$ MHz), users with minimum dedicated bandwidth are in cell edge $(f = 800$ MHz) [Figures are best viewed on colors].

FIGURE 4. Throughput change of a mobile dedicated BW [UE](#page-0-0) during inter-eNodeB handover (f = 1800 MHz).

service disruption when providing real-time services for some users. Our experimental results also take into account the mobility of the users (mostly pedestrian UEs) during real-world deployment of the QoS services. To experiment with mobility behaviour of dedicated BW [UEs,](#page-0-0) Fig. [4](#page-9-1) shows the throughput variation of a dedicated BW [UE](#page-0-0) during inter

FIGURE 5. UL performance of users (a) All [UEs](#page-0-0) are in cell center (both $f = 1800$ MHz) (b) Normal [UE](#page-0-0) is in cell center $(f = 1800 \text{ MHz})$, priority users are in cell edge $(f = 800 \text{ MHz})$ [Figures are best viewed on colors].

eNodeB handover between two feature-enabled [BSs.](#page-0-0) Dedicated BW [UE](#page-0-0) can obtain more than two times the throughput in *PCI* − 378 cell (with average throughput value of 19716 kbps) compared to *PCI* − 209 cell (with average throughput value of 9232 kbps). This is due to the fact that there are high number of dedicated BW [UEs](#page-0-0) in *PCI* − 209 cell in addition to normal [UEs,](#page-0-0) which increase the amount of traffic in buffer of the scheduler. In comparison, *PCI* − 378 cell has less number of dedicated BW [UEs,](#page-0-0) so that dedicated BW [UEs](#page-0-0) can get higher dedicated bandwidth. On the other hand, the throughput of dedicated BW [UE](#page-0-0) diminishes to average value of 2386 kbps during handover where the dedicated [RB](#page-0-0) allocation requirement is violated significantly. These momentary changes in dedicated BW [UE](#page-0-0) throughput values have demonstrated how performance of mobile dedicated [UE](#page-0-0) between cells can differ as a result of the availability of [UE](#page-0-0) load with different [QoS](#page-0-0) deployment strategies inside the cell.

Fig. [5](#page-10-0) shows the experimental scenario where there are one normal [UE](#page-0-0) and $N_p^s = 10$ priority [UEs](#page-0-0) that obtain throughput values based on [\(8\)](#page-5-2). Similar to previous scenario,

while normal [UE](#page-0-0) is generating traffic, 10 higher priority [UEs](#page-0-0) are starting to generate traffic sequentially in time. Fig. [5a](#page-10-0) shows the scenario where all [UEs](#page-0-0) are in cell center whereas Fig. [5b](#page-10-0) shows the scenario when normal [UE](#page-0-0) is in cell center and 10 priority [UEs](#page-0-0) are in cell edge. From Fig. [5a](#page-10-0), it is observed that the expected theoretical 40% to 60% throughput ratio split between normal and priority [UEs](#page-0-0) respectively has been achieved in this experimental set-up. From Fig. [5b](#page-10-0), we can observe that after all priority [UEs](#page-0-0) are on cell-edge, normal [UEs'](#page-0-0) throughput values remain higher than the rest of the [UEs](#page-0-0) due to proximity to eNodeB. The throughput values of priority [UEs](#page-0-0) have diminished, hence the expected 40% to 60% throughput ratio split has not been achieved. These results indicate that even though theoretical throughput split can be achieved in cell-center scenarios, due to poor channel conditions in cell-edge, the designed scheduler performance cannot achieve a successful throughput split between normal and priority [UEs.](#page-0-0)

Fig. [6](#page-11-0) shows the amount of throughput generated by three [UE](#page-0-0) types namely normal, priority and dedicated BW [UE](#page-0-0)

FIGURE 6. UL performance of UEs (a) All UEs are in cell center (both $f = 1800$ MHz) (b) Normal UE is in cell center $(f = 1800$ MHz), 1 priority UE and 1 dedicated BW UE are in cell edge $(f = 800$ MHz) [Figures are best viewed on colors].

over the experiment duration. All [UEs](#page-0-0) are located in the cell center in Fig. [6a](#page-11-0) whereas normal [UE](#page-0-0) is in cell center and other [UEs](#page-0-0) (namely dedicated BW and priority ones) are located at cell edge in Fig. [6b](#page-11-0). First of all, it can be observed that once the priority [UE](#page-0-0) enters into the cell, the throughput values of the normal [UE](#page-0-0) drop accordingly as shown in both figures of Fig. [6.](#page-11-0) As expected in Fig. [6a](#page-11-0), the [UL](#page-0-0) throughput values of priority [UE](#page-0-0) are higher than normal [UE.](#page-0-0) The expected throughput split of 40% to 60% ratio among normal and prioritized [UE](#page-0-0) (UE-1) respectively has been partially achieved where deep fades in normal [UE](#page-0-0) also effects priority [UEs'](#page-0-0) throughput values in consecutive time intervals due to adaptation of scheduler allocations. Transmit buffer size of the different types of [UEs](#page-0-0) can also impact the obtained throughput values. Additionally, dedicated [UE](#page-0-0) (UE-2) has achieved less throughput compared to normal and priority [UE](#page-0-0) due to poor channel conditions even though it can obtain target of 20 RBs utilization.

In Fig. [6b](#page-11-0), both dedicated BW and priority [UEs](#page-0-0) have obtained approximately the same [UL](#page-0-0) throughput values which is below 4 Mbps. On the other hand, normal [UE](#page-0-0) has obtained the highest throughput value since priority [UEs](#page-0-0) are now located at cell edge and normal [UE](#page-0-0) is in cell center. Moreover, cell center [UE](#page-0-0) is using $f = 1800$ MHz whereas cell edge [UEs](#page-0-0) are now using $f = 800$ MHz (due to higher coverage potential at large distances) which also has an effect in throughput reductions. This scenario is planned by network operations experts to provide connectivity rather than higher data rates at far locations to [BSs.](#page-0-0) Even though the scheduler is configured to distribute resources 3 to 2 ratio among priority and normal [UEs](#page-0-0) respectively, the obtained throughput values in Fig. [6b](#page-11-0) differ due to low [RSRP](#page-0-0) values for priority [UE.](#page-0-0) Note that the eNodeB scheduler considers different metrics including [MCS](#page-0-0) index values. prior obtained throughput values, etc. during resource allocation in addition to statically assigned weight configuration defined in Section [III-B.](#page-5-3) Moreover, dedicated BW [UE](#page-0-0) has achieved the same minimum required [RBs](#page-0-0) utilization of $K_f = 20$. However, due to poor channel conditions Fig. [6b](#page-11-0)'s throughput values are low compared to Fig. [6a](#page-11-0) throughput values.

Fig. [7](#page-12-0) shows the experimental scenario when there are one normal [UE,](#page-0-0) $N_p^s = 5$ priority [UEs](#page-0-0) and $N_p^d = 5$ dedicated BW [UEs](#page-0-0) inside cell coverage. As before, 5 priority [UEs](#page-0-0) and 5 dedicated BW [UEs](#page-0-0) start to upload traffic inside the considered cell sequentially in time. Fig. [7a](#page-12-0) demonstrates the scenario when all [UEs](#page-0-0) are in cell center with good RF conditions whereas Fig. [7b](#page-12-0) shows the scenario where normal [UE](#page-0-0) is in cell center, 5 priority [UEs](#page-0-0) and 5 dedicated BW [UEs](#page-0-0)

FIGURE 7. UL performance of UEs (a) All UEs are in cell center (both f = 1800 MHz) (b) Normal UE is in cell center $(f = 1800$ MHz), 5 priority UEs and 5 dedicate BW UEs are in cell edge $(f = 800$ MHz) [Figures are best viewed on colors].

are at cell edge. From Fig. [7a](#page-12-0), we can observe that throughput values of dedicated BW [UEs](#page-0-0) are around 4 Mbps. We can also observe that as number of [UEs](#page-0-0) increases in time and more specifically after UE-6 (dedicated BW [UE\)](#page-0-0) enters into cell, normal and priority [UEs'](#page-0-0)s throughput values start to diminish and become zero value after dedicated BW UE-10 (fifth dedicated UE) enters into coverage area. This is due to lack of resources, i.e. unavailability of [RBs,](#page-0-0) which signifies that [LTE](#page-0-0) eNodeB cannot schedule resources for these normal and priority [UEs.](#page-0-0) Hence, the expected throughput split between normal and priority [UEs](#page-0-0) of 40% to 60% ratio respectively has not been achieved in this experimental scenario even though the channel quality is good for all the [UEs.](#page-0-0) Fig. [7b](#page-12-0) shows that after all priority and dedicated BW [UEs](#page-0-0) are located at cell-edge, dedicated BW [UEs](#page-0-0) start to obtain less throughput values (e.g. after UE-7 (dedicated BW) enters). After UE-10 (priority) enters, throughput of all dedicated BW [UEs](#page-0-0) start to diminish together. On the other hand, the performance of normal and priority [UEs](#page-0-0) becomes worse than the scenario in Fig. [7a](#page-12-0) where the data rates diminish significantly after UE-5 (priority) enters. After this point, normal and priority [UEs](#page-0-0) obtain no resources as UE-9 (dedicated BW) is connected to the same cell. In summary, the results in Fig. [7](#page-12-0) indicate that the diminishing effect of lack of radio resources can be observed directly on priority [UEs](#page-0-0) and normal [UEs](#page-0-0) rather than dedicated BW [UEs.](#page-0-0)

C. MAIN OBSERVATIONS AND TAKEAWAYS

In summary, we have tested three main types of experimental scenarios using the experimental set-up. First one is with priority and normal [UEs,](#page-0-0) second one is with dedicated BW and normal [UEs](#page-0-0) and final one is with priority, dedicated BW and normal [UEs.](#page-0-0) Our experimental observations have revealed that location of the [UE](#page-0-0) with respect to [BS](#page-0-0) and the availability of dedicated BW [UEs](#page-0-0) inside cell may have implications on the apriori defined resource allocation strategies of

the other [UEs.](#page-0-0) For example, the results in Fig. [5](#page-10-0) indicate that even though theoretical throughput split can be achieved in cell-center scenarios, due to poor channel conditions in celledge, the designed scheduler performance cannot achieve a successful throughput split between normal and priority [UEs](#page-0-0) even though dedicated BW [UE](#page-0-0) have obtained lower throughput values. This is also true in poor channel conditions of priority [UE](#page-0-0) as observed from experimental results of Fig. [6b](#page-11-0). These results signify that before deploying critical or non-critical services in a cellular network, network planning and optimization should consider the available number of diverse set of [UEs](#page-0-0) with different [QoS](#page-0-0) requirements in the surrounding area meticulously to avoid reaching limitations on eNodeB capacity.

Another important observation to consider is that dedicated BW [UEs](#page-0-0) should be placed in cell-center locations after careful network and coverage planning. In case dedicated BW [UEs](#page-0-0) are in cell-edge areas, they can have diminishing performances on dedicated BW [UEs,](#page-0-0) but will also have huge impact on normal and priority [UEs](#page-0-0) that are located in cell center regions by allocating their [RBs](#page-0-0) in exchange of poor throughput values. Dedicated BW [UEs](#page-0-0) is primarily designed for static [UEs](#page-0-0) such as [ATMs](#page-0-0) of a bank. In case, this feature is enabled for mobile dedicated BW [UEs,](#page-0-0) they can have detrimental impact on new cells in case normal and priority [UEs](#page-0-0) exist. Moreover, dedicated BW [UEs](#page-0-0) can experience severe throughput decrements during handover period due to not completed [RB](#page-0-0) allocation strategies. If a mission-critical network service is running over mobile dedicated BW [UEs,](#page-0-0) these handover interruptions can disrupt the service continuity. It has also been observed in [UL](#page-0-0) traffic tests that the [BS](#page-0-0) is scheduling rate for a [UE](#page-0-0) depends on the transmit buffer size of the [UE.](#page-0-0) Therefore, [UEs](#page-0-0) with high buffer size will be scheduled further. The difference between the theoretically calculated maximum [UE](#page-0-0) throughput amount and the throughput that the [UE](#page-0-0) can actually practically achieve is created by the transmit buffer size of the [UE.](#page-0-0) Furthermore, if the buffer sizes of the priority and normal [UEs](#page-0-0) are different, this would be an advantage for [UEs](#page-0-0) with high buffer size in terms of obtained throughput values and the targeted 60% to 40% ratios would not be provided.

If there were no dedicated BW [UEs](#page-0-0) in the system design and only [UEs](#page-0-0) with different priorities were present, the resource allocation problem would be solved by prioritizing these [UEs](#page-0-0) and their services simply based on priority ordering during scheduling interval. However, the problem arises on [MNO'](#page-0-0)s concern about [UEs](#page-0-0) with dedicated bandwidth requirements. To provide dedicated bandwidth, [MNO](#page-0-0) assigns the dedicated BW [UE](#page-0-0) to a higher priority [QoS,](#page-0-0) i.e. [QCI](#page-0-0) priority higher than the priority [UEs.](#page-0-0) Although this is positive in terms of providing dedicated bandwidth, the dedicated BW [UEs](#page-0-0) can deplete resources of the other [UEs.](#page-0-0) Therefore, an upper bound limitation on number of connected dedicated BW [UEs](#page-0-0) is needed by careful radio network capacity planning at each site. However, this situation cannot be prevented in case there is a single-tier scheduler

in the system. If a multi-tier scheduler was present, each [UE](#page-0-0) type could be scheduled on its own tier. However, in this case the complexity of the system would increase. The network slicing concept that comes with 5G networks can actually be a suitable solution for the problem presented in this paper. Dedicated slicing, a deployment method of network slicing, can be implemented according to different [QoS](#page-0-0) strategies since it can have a separate scheduler for each slice type.

V. CONCLUSION AND FUTURE WORK

In this paper, we have investigated a potential solution of providing better [QoS](#page-0-0) support for differentiated types of [UEs](#page-0-0) to improve the performance of the services provided by [MNOs.](#page-0-0) For this, we first formalized the optimization problem in a formal manner. Later, we described our experimental set-up where experimental evaluation of different [QoS](#page-0-0) deployment strategies are performed in a real operational network in Turkey. We have also analyzed the necessary network planning in detail for deploying the proposed QoS strategies. Our real-world experiments on a LTE network indicated that prioritized [UEs](#page-0-0) with dedicated bandwidth have higher precedence compared to prioritized [UEs](#page-0-0) with high scheduling rate and normal [UEs.](#page-0-0) Moreover, the theoretical 40% to 60% ratio of throughput split between prioritized [UEs](#page-0-0) with high scheduling rate and normal [UEs](#page-0-0) can only be achieved as long as the amount of dedicated resources allocated to dedicated BW [UEs](#page-0-0) is carefully planned during network capacity optimization stage. This signifies that before deployment of [QoS](#page-0-0) policies for any critical or non-critical services, [MNOs](#page-0-0) need to perform extensive experimental trials to find the best configuration and optimization parameters to serve all [UEs](#page-0-0) based on their assigned [QCI](#page-0-0) levels. As a result, [MNOs](#page-0-0) can only gain major benefits by differentiating priority [UEs](#page-0-0) and diversifying the network services provided for their [UEs](#page-0-0) via appropriate network planning. For future study, the field [KPIs](#page-0-0) and service use cases can be analyzed with [Machine](#page-0-0) [Learning \(ML\)](#page-0-0) and [QoS](#page-0-0) deployment strategy can be changed dynamically. Moreover, noting that we have only considered a single shared [S-GW](#page-0-0) in our both optimization problem and experimental scenario, a more general setting in large-scale deployments could be that multiple such [S-GWs](#page-0-0) can be utilized for sharing similar resources among differentiated UEs. This extension of the model would lead to a multi-layer problem and would be an interesting future work direction.

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