IEEE Access
Multidisciplinary : Rapid Review : Open Access Journal

Received April 27, 2020, accepted June 4, 2020, date of publication June 11, 2020, date of current version June 24, 2020.

Digital Object Identifier 10.1109/ACCESS.2020.3001724

Design of Cooperative MIMO Wireless Sensor Networks With Partial Channel State Information

DONATELLA DARSENA[1,2], (Senior Member, IEEE), GIACINTO GELLI[2,3], (Senior Member, IEEE), AND FRANCESCO VERDE[2,3], (Senior Member, IEEE)

[1]Department of Engineering, University of Napoli Parthenope, I-80143 Naples, Italy
[2]National Inter-University Consortium for Telecommunications (CNIT), 80125 Naples, Italy
[3]Department of Electrical Engineering and Information Technology, University Federico II, I-80125 Naples, Italy

Corresponding author: Francesco Verde (f.verde@unina.it)

ABSTRACT Wireless sensor networks (WSNs) play a key role in automation and consumer electronics applications. This paper deals with joint design of the source precoder, relaying matrices, and destination equalizer in a multiple-relay amplify-and-forward (AF) cooperative multiple-input multiple-output (MIMO) WSN, when partial channel-state information (CSI) is available in the network. In particular, the considered approach assumes knowledge of instantaneous CSI of the first-hop channels and statistical CSI of the second-hop channels. In such a scenario, compared to the case when instantaneous CSI of both the first- and second-hop channels is exploited, existing network designs exhibit a significant performance degradation. Relying on a relaxed minimum-mean-square-error (MMSE) criterion, we show that strategies based on potential activation of all antennas belonging to all relays lead to mathematically intractable optimization problems. Therefore, we develop a new joint relay-and-antenna selection procedure, which determines the best subset of the available antennas possibly belonging to different relays. Monte Carlo simulations show that, compared to conventional relay selection strategies, the proposed design offers a significant performance gain, outperforming also other recently proposed relay/antenna selection schemes.

INDEX TERMS Amplify-and-forward relays, multiple-input multiple-output (MIMO), partial channel state information, wireless sensor networks.

I. INTRODUCTION AND SYSTEM MODEL

With the advent of massive Internet of Things and massive machine-type communications, especially in the domain of 5G consumer electronics, there is a need to further enhance physical-layer performance of *wireless sensor networks (WSNs)*. In this respect, *amplify-and-forward (AF)* relaying is an effective way to improve transmission reliability over fading channels, by taking advantage of the broadcast nature of wireless communications [1], [2], especially when the network nodes are equipped with multiple-input multiple-output (MIMO) transceivers [3]–[5].

We consider a one-way cooperative MIMO WSN aimed at transmitting a symbol block $\boldsymbol{b} \in \mathbb{C}^{N_\mathrm{B}}$ from a source to a destination, with the assistance of N_C half-duplex relays.[1]

The associate editor coordinating the review of this manuscript and approving it for publication was Liang Yang.

[1]The fields of complex and numbers are denoted with $\mathbb{C}$ and $\mathbb{R}$, respectively; matrices [vectors] are denoted with upper [lower] case boldface letters (e.g., $\boldsymbol{A}$ or $\boldsymbol{a}$); the field of $m \times n$ complex [real] matrices is denoted as $\mathbb{C}^{m\times n}$ [$\mathbb{R}^{m\times n}$], with $\mathbb{C}^m$ [$\mathbb{R}^m$] used as a shorthand for $\mathbb{C}^{m\times 1}$ [$\mathbb{R}^{m\times 1}$]; the superscripts $*$, T, H, -1, and $\dagger$ denote the conjugate, the transpose, the conjugate transpose, the inverse, and the Moore-Penrose generalized inverse of a matrix, respectively; $\{\boldsymbol{A}\}_{ij}$ indicates the $(i+1, j+1)$th element of $\boldsymbol{A} \in \mathbb{C}^{m\times n}$, with $i \in \{0, 1, \ldots, m-1\}$ and $j \in \{0, 1, \ldots, n-1\}$; $\boldsymbol{0}_m \in \mathbb{R}^m$, $\mathbf{O} \in \mathbb{R}^{m\times n}$, and $\boldsymbol{I}_m \in \mathbb{R}^{m\times m}$ denote the null vector, the null matrix, and the identity matrix, respectively; $\mathrm{tr}(\boldsymbol{A})$ denotes the trace of $\boldsymbol{A} \in \mathbb{C}^{n\times n}$; $\mathrm{rank}(\boldsymbol{A})$ is the rank of $\boldsymbol{A} \in \mathbb{C}^{m\times n}$; finally, the operator $\mathbb{E}[\,\cdot\,]$ denotes ensemble averaging.

We assume that there is no direct link between the source and the destination, due to high path loss values or obstructions, and we denote with N_S, N_R, and N_D, respectively, the numbers of antennas at the source, relays, and destination. The received signal at the destination can be expressed as

$$\boldsymbol{r} = \boldsymbol{C}\,\boldsymbol{b} + \boldsymbol{v} \tag{1}$$

where $\boldsymbol{C} \triangleq \boldsymbol{G}\,\boldsymbol{F}\,\boldsymbol{H}\,\boldsymbol{F}_0 \in \mathbb{C}^{N_\mathrm{D}\times N_\mathrm{B}}$ is the *dual-hop* channel matrix and $\boldsymbol{v} \triangleq \boldsymbol{G}\,\boldsymbol{F}\,\boldsymbol{w} + \boldsymbol{n}$ is the equivalent noise vector at the destination. The composite matrices

$$\boldsymbol{H} \triangleq [\boldsymbol{H}_1^\mathrm{T}, \boldsymbol{H}_2^\mathrm{T}, \ldots, \boldsymbol{H}_{N_\mathrm{C}}^\mathrm{T}]^\mathrm{T} \in \mathbb{C}^{(N_\mathrm{C}N_\mathrm{R})\times N_\mathrm{S}} \tag{2}$$

$$\boldsymbol{G} \triangleq [\boldsymbol{G}_1, \boldsymbol{G}_2, \ldots, \boldsymbol{G}_{N_c}] \in \mathbb{C}^{N_\mathrm{D}\times(N_\mathrm{C}N_\mathrm{R})} \tag{3}$$

collect the *first-* (backward) and *second-hop* (forward) MIMO channel coefficients of all the relays, respectively, whereas the diagonal blocks $\boldsymbol{F}_i \in \mathbb{C}^{N_\mathrm{R}\times N_\mathrm{R}}$ of

$$\boldsymbol{F} \triangleq \mathrm{diag}(\boldsymbol{F}_1, \boldsymbol{F}_2, \ldots, \boldsymbol{F}_{N_\mathrm{C}}) \tag{4}$$

denote the relaying matrices, and $\boldsymbol{F}_0 \in \mathbb{C}^{N_\mathrm{S}\times N_\mathrm{B}}$ represents the source precoding matrix. Finally, $\boldsymbol{w} \in \mathbb{C}^{N_\mathrm{C}N_\mathrm{R}}$ and $\boldsymbol{n} \in \mathbb{C}^{N_\mathrm{D}}$ gather the noise samples at all the relays and at the destination, respectively. The vector $\boldsymbol{r}$ is subject to linear equalization at the destination through the equalizing matrix $\boldsymbol{D} \in \mathbb{C}^{N_\mathrm{B}\times N_\mathrm{D}}$, hence yielding an estimate $\hat{\boldsymbol{b}} \triangleq \boldsymbol{D}\boldsymbol{r}$ of the source block $\boldsymbol{b}$, whose entries are then subject to minimum-distance (in the Euclidean sense) detection. Increase in spectral efficiency can be obtained by considering *two-way relaying* [6], which is based on establishing bidirectional connections between two or more terminals using one or several half-duplex relays.

To achieve the expected gains, channel state information (CSI) is required at the network nodes, i.e, source, AF relays, and destination. *Full CSI (F-CSI)* is invoked in many papers dealing with optimization of one-way (see, e.g., [7]–[15]) and two-way (see, e.g., [16], [17]) cooperative MIMO networks. Specifically, with reference to the system model (1), F-CSI is tantamount to requiring: (i) instantaneous knowledge of the first-hop channel matrix $\boldsymbol{H}$; (ii) instantaneous knowledge of the second-hop channel matrix $\boldsymbol{G}$; (iii) instantaneous knowledge of the dual-hop channel matrix $\boldsymbol{C}$. While the dual-hop channel matrix $\boldsymbol{C}$ can be directly estimated at the destination by training, separate acquisition of the first- and second-hop matrices $\boldsymbol{H}$ and $\boldsymbol{G}$ is more complicated to achieve, both in terms of communication resources and signal overhead, especially in multiple-relay WSNs. Moreover, since channel estimation errors occur in practical situations, robust optimization designs are needed [18], [19], which further complicate system deployment. In resource-constrained WSNs, the use of *partial* CSI (P-CSI) can extend network lifetime and reduce the complexity burden.

Relay selection is a common strategy to reduce signaling overhead and system design complexity in single-input single-output (SISO) cooperative WSNs [20]–[23]. Design of SISO relay selection procedures providing diversity gains – even when F-CSI is not available – has been addressed in [24]–[26], [26], [27]. Such methods rely on P-CSI, since selection of the best relay is based only on instantaneous knowledge of the source-to-relay channels. However, the diversity order of the methods developed in these papers does not scale in the number of relays $N_{\rm C}$. For SISO nodes, a P-CSI relay selection scheme has been proposed in [28], yielding full diversity order $N_{\rm C}$. However, besides the instantaneous knowledge of the source-to-relay channels, such a method requires that the selected relay sends instantaneous CSI of the corresponding source-to-relay channel to the destination for optimal decoding. Moreover, the optimization problem in [28] does not admit a closed-form solution and is solved by using a line search algorithm.

It has been shown in [29] that P-CSI relay selection approaches for MIMO nodes, based only on the instantaneous knowledge of $\boldsymbol{H}$, do not fully exploit the diversity arising from the presence of multiple relays. Besides instantaneous knowledge of $\boldsymbol{H}$, statistical CSI of the second-hop matrix $\boldsymbol{G}$ is used in [14], [30] to perform relay/antenna selection for a MIMO AF cooperative network. However, the solutions developed in [14], [30] still exhibit a significant performance degradation compared to designs based on F-CSI.

In this paper, we present new optimization methods for multiple-relay cooperative MIMO WSNs with P-CSI, i.e., knowledge of the instantaneous value of $\boldsymbol{H}$ and the statistical properties of $\boldsymbol{G}$. Our design does not rely on F-CSI as in [7]–[13], [15]–[17], and needs the same amount of P-CSI exploited in [24]–[26], [26], [27], [30]. In this scenario, we consider a relaxed joint minimum-mean-square-error (MMSE) optimization of the source precoder $\boldsymbol{F}_0$, the AF relaying matrices in $\boldsymbol{F}$, and the destination equalizer $\boldsymbol{D}$, with a power constraint at the source [31] and *a sum-power constraint* at the relays [10]. Specifically, capitalizing on our preliminary results [14], the novel contributions can be summarized as follows:

1) We prove that the MMSE-based design attempting to activate all possible antennas of all relays leads to a mathematically intractable optimization problem.
2) We provide the proofs of the results reported in [14], by enlightening that single relay selection [14], [24]–[26], [26], [27] is suboptimal in the considered P-CSI scenario.
3) We develop a new joint antenna-and-relay selection algorithm, which is shown to significantly outperform the relay/antenna selection approaches [14], [26], [30] in terms of average symbol error probability (ASEP).

The paper is organized as follows. Section II introduces the basic assumptions and discusses their practical implications. The proposed designs are developed in Section III. Section IV reports simulation results in terms of ASEP, whereas Section V draws some conclusions.

II. BASIC ASSUMPTIONS AND PRELIMINARIES

The symbol block $\boldsymbol{b}$ in (1) is modeled as a circularly symmetric complex random vector, with $\mathbb{E}[\boldsymbol{b}\,\boldsymbol{b}^{\rm H}] = \boldsymbol{I}_{N_{\rm B}}$. The entries of $\boldsymbol{H}$ and $\boldsymbol{G}$ are assumed to be unit-variance circularly symmetric complex Gaussian (CSCG) random variables. The noise vectors $\boldsymbol{w}$ and $\boldsymbol{n}$ are modeled as mutually independent CSCG random vectors, statistically independent of $(\boldsymbol{b}, \boldsymbol{H}, \boldsymbol{G})$, with $\mathbb{E}[\boldsymbol{w}\,\boldsymbol{w}^{\rm H}] = \boldsymbol{I}_{N_{\rm C}N_{\rm R}}$ and $\mathbb{E}[\boldsymbol{n}\,\boldsymbol{n}^{\rm H}] = \boldsymbol{I}_{N_{\rm D}}$, respectively.

Hereinafter, we assume that $\boldsymbol{C}$ in (1) and the following conditional covariance matrix of $\boldsymbol{v}$, given $\boldsymbol{G}$,

$$\boldsymbol{K}_{\boldsymbol{vv}} \triangleq \mathbb{E}[\boldsymbol{v}\,\boldsymbol{v}^{\rm H} \,|\, \boldsymbol{G}] = \boldsymbol{G}\,\boldsymbol{F}\,\boldsymbol{F}^{\rm H}\boldsymbol{G}^{\rm H} + \boldsymbol{I}_{N_{\rm D}} \tag{5}$$

have been previously acquired at the destination during a training session. Under such assumptions, it is known (see, e.g., [31]) that, for fixed matrices $\boldsymbol{F}_0$ and $\boldsymbol{F}$, the matrix $\boldsymbol{D}$ minimizing the trace of the conditional mean square error (MSE) matrix $\boldsymbol{E}(\boldsymbol{F}_0, \boldsymbol{F}, \boldsymbol{D}) \triangleq \mathbb{E}[(\hat{\boldsymbol{b}} - \boldsymbol{b})\,(\hat{\boldsymbol{b}} - \boldsymbol{b})^{\rm H} \,|\, \boldsymbol{H}, \boldsymbol{G}]$, given $\boldsymbol{H}$ and $\boldsymbol{G}$, is the Wiener filter

$$\boldsymbol{D}_{\rm mmse} = \boldsymbol{C}^{\rm H}(\boldsymbol{C}\,\boldsymbol{C}^{\rm H} + \boldsymbol{K}_{\boldsymbol{vv}})^{-1}\,. \tag{6}$$

Optimization of $\boldsymbol{F}_0$ and $\boldsymbol{F}$ is carried out under the assumption that only P-CSI is available at the source and the relays. Specifically, the source and the relays perfectly know the first-hop channel matrix $\boldsymbol{H}$, but the ith relay has only knowledge of the second-order statistics (SOS) of its own second-hop channel matrix $\boldsymbol{G}_i$. These assumptions are justified since, in some systems, the relays may be able to exchange information among themselves before transmission [32]. In this case, knowledge of $\boldsymbol{H}$ at the relays is realistic [7]–[14], [20]–[22], [24]–[29]. Moreover, since the SOS of $\boldsymbol{G}_i$ vary much more slowly than the instantaneous values of $\boldsymbol{G}_i$, the feedback overhead from the destination to the relays is significantly reduced, compared to [28].

III. THE PROPOSED P-CSI-BASED DESIGN

To obtain $\boldsymbol{F}_0$ and $\boldsymbol{F}$, we minimize the statistical average (with respect to $\boldsymbol{G}$) of the trace of the following matrix:

$$\boldsymbol{E}(\boldsymbol{F}_0, \boldsymbol{F}) \triangleq \boldsymbol{E}(\boldsymbol{F}_0, \boldsymbol{F}, \boldsymbol{D}_{\text{mmse}}) = (\boldsymbol{I}_{N_\text{B}} + \boldsymbol{C}^\text{H} \boldsymbol{K}_{vv}^{-1} \boldsymbol{C})^{-1} \quad (7)$$

under suitable power constraints. To this aim, we assume that: **a1)** $\boldsymbol{F}_0$ is full-column rank, i.e., $\text{rank}(\boldsymbol{F}_0) = N_\text{B} \le N_\text{S}$; **a2)** $\boldsymbol{GFH}$ is full-column rank, i.e., $\text{rank}(\boldsymbol{GFH}) = N_\text{S} \le N_\text{D}$. It is noteworthy that assumption **a2)** necessarily requires that the matrices $\boldsymbol{FH}$ and $\boldsymbol{H}$ are full-column rank, i.e., $\text{rank}(\boldsymbol{FH}) = \text{rank}(\boldsymbol{H}) = N_\text{S} \le N_\text{C} N_\text{R}$. Such assumptions ensure that $\boldsymbol{C}$ is full-column rank as well. Specifically, we consider the following optimization problem:

$$\begin{aligned} &\min_{\boldsymbol{F}_0, \boldsymbol{F}} \mathbb{E}_{\boldsymbol{G}} \left\{ \text{tr}\left[\left(\boldsymbol{I}_{N_\text{B}} + \boldsymbol{C}^\text{H} \boldsymbol{K}_{vv}^{-1} \boldsymbol{C} \right)^{-1} \right] \mid \boldsymbol{H} \right\} \\ &\text{subject to (s.to) } \text{tr}(\boldsymbol{F}_0 \boldsymbol{F}_0^\text{H}) \le \mathscr{P}_\text{S} \\ &\qquad \text{and } \mathbb{E}_{\boldsymbol{G}} \left[\text{tr}(\boldsymbol{G}\boldsymbol{F}\boldsymbol{K}_{zz}\boldsymbol{F}^\text{H}\boldsymbol{G}^\text{H}) \mid \boldsymbol{H} \right] \le \mathscr{P}_\text{D} \end{aligned} \quad (8)$$

where $\boldsymbol{K}_{zz} \triangleq \mathbb{E}[\boldsymbol{z}\boldsymbol{z}^\text{H} \mid \boldsymbol{H}] = \boldsymbol{H}\boldsymbol{F}_0\boldsymbol{F}_0^\text{H}\boldsymbol{H}^\text{H} + \boldsymbol{I}_{N_\text{C}N_\text{R}}$ is the conditional (given $\boldsymbol{H}$) covariance matrix of the vector $\boldsymbol{z} \in \mathbb{C}^{N_\text{C}N_\text{R}}$ collecting the signals received by all the relays, with $\mathscr{P}_\text{S} > 0$ and $\mathscr{P}_\text{D} > 0$ denoting the power threshold at the source and at the destination, respectively. The constraint on the received power at the destination automatically limits the power expenditure at the relays. Since problem (8) is nonconvex, we consider its *relaxed* version:

$$\begin{aligned} &\min_{\boldsymbol{F}_0, \boldsymbol{F}} \mathbb{E}_{\boldsymbol{G}} \left\{ \text{tr}\left[\left(\boldsymbol{I}_{N_\text{B}} + \boldsymbol{F}_0^\text{H}\boldsymbol{H}^\text{H}\boldsymbol{F}^\text{H}\boldsymbol{G}^\text{H}\boldsymbol{G}\boldsymbol{F}\boldsymbol{H}\boldsymbol{F}_0 \right)^{-1} \right] \mid \boldsymbol{H} \right\} \\ &\text{s.to } \text{tr}(\boldsymbol{F}_0 \boldsymbol{F}_0^\text{H}) \le \mathscr{P}_\text{S} \quad \text{and} \\ &\qquad \mathbb{E}_{\boldsymbol{G}} \left[\text{tr}(\boldsymbol{G}\boldsymbol{F}\boldsymbol{F}^\text{H}\boldsymbol{G}^\text{H}) \mid \boldsymbol{H} \right] \le \mathscr{P}_\text{D} \end{aligned} \quad (9)$$

where we have used the expression of $\boldsymbol{C}$ and the inequalities $\text{tr}[(\boldsymbol{I}_{N_\text{B}} + \boldsymbol{C}^\text{H}\boldsymbol{K}_{vv}^{-1}\boldsymbol{C})^{-1}] \ge \text{tr}[(\boldsymbol{I}_{N_\text{B}} + \boldsymbol{C}^\text{H}\boldsymbol{C})^{-1}]$ and $\text{tr}(\boldsymbol{G}\boldsymbol{F}\boldsymbol{K}_{zz}\boldsymbol{F}^\text{H}\boldsymbol{G}^\text{H}) \le \text{tr}(\boldsymbol{G}\boldsymbol{F}\boldsymbol{F}^\text{H}\boldsymbol{G}^\text{H})\,\text{tr}(\boldsymbol{K}_{zz})$ [33], [34]. Closed-form evaluation of the cost function in (9) is cumbersome; however, under **a1)** and **a2)**, it can be observed that[2]

$$\begin{aligned} &\text{tr}\left[\left(\boldsymbol{I}_{N_\text{B}} + \boldsymbol{F}_0^\text{H}\boldsymbol{H}^\text{H}\boldsymbol{F}^\text{H}\boldsymbol{G}^\text{H}\boldsymbol{G}\boldsymbol{F}\boldsymbol{H}\boldsymbol{F}_0 \right)^{-1} \right] \\ &\qquad < \text{tr}\left[\left(\boldsymbol{F}_0^\text{H}\boldsymbol{H}^\text{H}\boldsymbol{F}^\text{H}\boldsymbol{G}^\text{H}\boldsymbol{G}\boldsymbol{F}\boldsymbol{H}\boldsymbol{F}_0 \right)^{-1} \right] \end{aligned} \quad (10)$$

where the difference between the left- and right-hand sides tends to zero as the minimum eigenvalue of $\boldsymbol{F}_0^\text{H}\boldsymbol{H}^\text{H}\boldsymbol{F}^\text{H}\boldsymbol{G}^\text{H}\boldsymbol{G}\boldsymbol{F}\boldsymbol{H}\boldsymbol{F}_0$ is significantly larger than one. This happens in the high signal-to-noise ratio (SNR) region, i.e., when $\mathscr{P}_\text{S}$ and $\mathscr{P}_\text{D}$ are sufficiently large. Relying on (10), we pursue a further relaxation of (8) by replacing $\mathbb{E}_{\boldsymbol{G}}\{\text{tr}[(\boldsymbol{I}_{N_\text{B}} + \boldsymbol{F}_0^\text{H}\boldsymbol{H}^\text{H}\boldsymbol{F}^\text{H}\boldsymbol{G}^\text{H}\boldsymbol{G}\boldsymbol{F}\boldsymbol{H}\boldsymbol{F}_0)^{-1}] \big| \boldsymbol{H}\}$ in (9) with its upper bound $\mathbb{E}_{\boldsymbol{G}}\{\text{tr}[(\boldsymbol{F}_0^\text{H}\boldsymbol{H}^\text{H}\boldsymbol{F}^\text{H}\boldsymbol{G}^\text{H}\boldsymbol{G}\boldsymbol{F}\boldsymbol{H}\boldsymbol{F}_0)^{-1}] \big| \boldsymbol{H}\}$, which can be evaluated in closed-form as stated by the following Lemma.

Lemma 1: Let us assume that: **a3)** $N_\text{D} > N_\text{B}$. Then, under **a1)**, **a2)**, and **a3)**, it results that

$$\mathbb{E}_{\boldsymbol{G}}\left\{ \text{tr}\left[\left(\boldsymbol{F}_0^\text{H}\boldsymbol{H}^\text{H}\boldsymbol{F}^\text{H}\boldsymbol{G}^\text{H}\boldsymbol{G}\boldsymbol{F}\boldsymbol{H}\boldsymbol{F}_0 \right)^{-1} \right] \big| \boldsymbol{H} \right\} = \frac{\text{tr}(\boldsymbol{R}^{-1})}{N_\text{D} - N_\text{B}} \quad (11)$$

where $\boldsymbol{R} \triangleq \boldsymbol{F}_0^\text{H}\boldsymbol{H}^\text{H}\boldsymbol{F}^\text{H}\boldsymbol{F}\boldsymbol{H}\boldsymbol{F}_0 \in \mathbb{C}^{N_\text{B} \times N_\text{B}}$.

Proof: See Appendix A. ■

At this point, evaluation of the expectation in the second constraint of (9) is in order. In this respect, one has

$$\mathbb{E}_{\boldsymbol{G}}\left[\text{tr}(\boldsymbol{F}^\text{H}\boldsymbol{G}^\text{H}\boldsymbol{G}\boldsymbol{F}) \mid \boldsymbol{H} \right] = \text{tr}\left[\mathbb{E}_{\boldsymbol{G}}\left(\boldsymbol{G}^\text{H}\boldsymbol{G} \right) \boldsymbol{F}\boldsymbol{F}^\text{H} \right] = \text{tr}\left(\boldsymbol{F}^\text{H}\boldsymbol{F} \right) \quad (12)$$

where we have also used the cyclic property [33] of the trace operator. Therefore, under **a1)**, **a2)**, and **a3)**, the optimization problem (9) can be simplified as follows

$$\begin{aligned} &\min_{\boldsymbol{F}_0, \boldsymbol{F}} \text{tr}\left[\left(\boldsymbol{F}_0^\text{H}\boldsymbol{H}^\text{H}\boldsymbol{F}^\text{H}\boldsymbol{F}\boldsymbol{H}\boldsymbol{F}_0 \right)^{-1} \right] \\ &\text{s.to } \text{tr}(\boldsymbol{F}_0\boldsymbol{F}_0^\text{H}) \le \mathscr{P}_\text{S} \quad \text{and } \text{tr}\left(\boldsymbol{F}^\text{H}\boldsymbol{F} \right) \le \mathscr{P}_\text{D}\,. \end{aligned} \quad (13)$$

At this point, a comment regarding the constraints in (8) and (13) is in order. The constraint $\text{tr}(\boldsymbol{F}_0\boldsymbol{F}_0^\text{H}) \le \mathscr{P}_\text{S}$ in (8) and (13) limits the average transmitted power of the source and it is standard in the design of linear MIMO transceivers [31]. Regarding the second constraint in (8), we observe that, given $\boldsymbol{H}$ and $\boldsymbol{G}$, $\mathscr{P}(\boldsymbol{H}, \boldsymbol{G}) \triangleq \text{tr}(\boldsymbol{G}\boldsymbol{F}\boldsymbol{K}_{zz}\boldsymbol{F}^\text{H}\boldsymbol{G}^\text{H})$ represents the average received power at the destination. It is noteworthy that $\mathscr{P}(\boldsymbol{H}, \boldsymbol{G})$ is typically limited in those scenarios where a target performance has to be achieved and per-node fairness is not of concern [3], [7]. The constraint $\text{tr}\left(\boldsymbol{F}^\text{H}\boldsymbol{F}\right) \le \mathscr{P}_\text{D}$ in (13), which has been obtained by averaging a relaxed version of $\mathscr{P}(\boldsymbol{H}, \boldsymbol{G})$ with respect to the probability distribution of $\boldsymbol{G}$, fixes a limit on the total average power transmitted by the relays, so-called sum-power constraint [10].[3]

To solve (13), we use the following Lemma.

Lemma 2: For a positive definite matrix $\boldsymbol{A} \in \mathbb{C}^{n \times n}$, the following inequality holds:

$$\text{tr}(\boldsymbol{A}^{-1}) \ge \sum_{\ell=1}^{m} \frac{1}{\{\boldsymbol{A}\}_{\ell\ell}} \quad (14)$$

where $\{\boldsymbol{A}\}_{\ell\ell}$ is the ℓth diagonal entry of $\boldsymbol{A}$ and the inequality is achieved if $\boldsymbol{A}$ is diagonal.

Proof: See [37, p. 65]. ■

As a consequence of Lemma 2, the minimum value of the cost function in (13) is achieved if $\boldsymbol{F}_0^\text{H}\boldsymbol{A}\boldsymbol{F}_0$ is diagonal, with

[2] The proof follows easily from the facts [33] that the trace of $\boldsymbol{A}$ is equal to the sum of its eigenvalues and, if λ is an eigenvalue of a nonsingular matrix $\boldsymbol{A}$, then λ^{-1} is an eigenvalue of $\boldsymbol{A}^{-1}$.

[3] Design with per-relay power constraints can be solved by properly reformulating the problem into an equivalent optimization with a sum-power constraint [35], [36].

$\boldsymbol{A} \triangleq \boldsymbol{H}^{\mathrm{H}}\boldsymbol{F}^{\mathrm{H}}\boldsymbol{F}\boldsymbol{H} \in \mathbb{C}^{N_{\mathrm{S}}\times N_{\mathrm{S}}}$. In what follows, we consider three different approaches to achieve the desired diagonalization of $\boldsymbol{F}_0^{\mathrm{H}}\boldsymbol{A}\,\boldsymbol{F}_0$: the first one is based on the SVD of the composite matrix $\boldsymbol{H} = [\boldsymbol{H}_1^{\mathrm{T}}, \boldsymbol{H}_2^{\mathrm{T}}, \ldots, \boldsymbol{H}_{N_{\mathrm{C}}}^{\mathrm{T}}]^{\mathrm{T}}$ and it results in a (possible) selection of all the relays; the second one relies on the SVDs of the individual matrices $\boldsymbol{H}_1, \boldsymbol{H}_2, \ldots, \boldsymbol{H}_{N_{\mathrm{C}}}$, thus leading to a single-relay selection; the last one exploits the SVDs of row-based partitions of $\boldsymbol{H}$ and it can be interpreted as a joint antenna-and-relay selection scheme.

A. DESIGN BASED ON THE SVD OF THE COMPOSITE FIRST-HOP CHANNEL MATRIX

One can attempt to recruit all the relays in the second hop of the cooperative scheme by diagonalizing $\boldsymbol{F}_0^{\mathrm{H}}\boldsymbol{A}\,\boldsymbol{F}_0$ through the SVD $\boldsymbol{H} = \boldsymbol{U}_{\mathrm{h}}\,[\boldsymbol{O}_{N_{\mathrm{S}}\times(N_{\mathrm{C}}N_{\mathrm{R}}-N_{\mathrm{S}})}, \boldsymbol{\Lambda}_{\mathrm{h}}]^{\mathrm{T}}\,\boldsymbol{V}_{\mathrm{h}}^{\mathrm{H}}$ of $\boldsymbol{H}$, where the matrices $\boldsymbol{U}_{\mathrm{h}} \in \mathbb{C}^{(N_{\mathrm{C}}N_{\mathrm{R}})\times(N_{\mathrm{C}}N_{\mathrm{R}})}$ and $\boldsymbol{V}_{\mathrm{h}} \in \mathbb{C}^{N_{\mathrm{S}}\times N_{\mathrm{S}}}$ are unitary, and $\boldsymbol{\Lambda}_{\mathrm{h}} \triangleq \mathrm{diag}[\lambda_{\mathrm{h}}(1), \lambda_{\mathrm{h}}(2), \ldots, \lambda_{\mathrm{h}}(N_{\mathrm{S}})]$ gathers the corresponding nonzero singular values arranged in increasing order. By substituting the SVD of $\boldsymbol{H}$ in $\boldsymbol{A}$, it follows by direct inspection that $\boldsymbol{F}_0^{\mathrm{H}}\boldsymbol{A}\,\boldsymbol{F}_0$ is diagonal if (see, e.g., [38])

$$\boldsymbol{F}_0 = \boldsymbol{V}_{\mathrm{h,right}}\,\boldsymbol{\Omega}^{1/2} \tag{15}$$

$$\boldsymbol{F}_i = \boldsymbol{Q}_i\,\boldsymbol{\Delta}_i^{1/2}\,\boldsymbol{U}_{\mathrm{h,right},i}^{\dagger} \tag{16}$$

where $\boldsymbol{V}_{\mathrm{h,right}} \in \mathbb{C}^{N_{\mathrm{S}}\times N_{\mathrm{B}}}$ contains the N_{B} rightmost columns from $\boldsymbol{V}_{\mathrm{h}}$, the matrices $\boldsymbol{\Omega} \triangleq \mathrm{diag}[\omega(1), \omega(2), \ldots, \omega(N_{\mathrm{B}})]$ and $\boldsymbol{\Delta}_i \triangleq \mathrm{diag}[\delta_i(1), \delta_i(2), \ldots, \delta_i(N_{\mathrm{S}})]$ are determined in a second step, for $i \in \{1, 2, \ldots, N_{\mathrm{C}}\}$, the arbitrary matrix $\boldsymbol{Q}_i \in \mathbb{C}^{N_{\mathrm{R}}\times N_{\mathrm{S}}}$ obeys $\boldsymbol{Q}_i^{\mathrm{H}}\boldsymbol{Q}_i = \boldsymbol{I}_{N_{\mathrm{S}}}$, provided that $N_{\mathrm{S}} \leq N_{\mathrm{R}}$, $\boldsymbol{U}_{\mathrm{h,right}} \triangleq [\boldsymbol{U}_{\mathrm{h,right},1}^{\mathrm{T}}, \boldsymbol{U}_{\mathrm{h,right},2}^{\mathrm{T}}, \ldots, \boldsymbol{U}_{\mathrm{h,right},N_C}^{\mathrm{T}}]^{\mathrm{T}} \in \mathbb{C}^{(N_{\mathrm{C}}N_{\mathrm{R}})\times N_{\mathrm{S}}}$ contains the N_{S} rightmost columns from $\boldsymbol{U}_{\mathrm{h}}$, with the matrix $\boldsymbol{U}_{\mathrm{h,right},i} \in \mathbb{C}^{N_{\mathrm{R}}\times N_{\mathrm{S}}}$ being full-column rank.

Using (15) and (16), problem (13) ends up to

$$\begin{aligned}
&\min_{\boldsymbol{\omega}, \{\boldsymbol{\delta}_i\}_{i=1}^{N_{\mathrm{C}}}} f_0\left(\boldsymbol{\omega}, \{\boldsymbol{\delta}_i\}_{i=1}^{N_{\mathrm{C}}}\right) \\
&\text{s.to} \sum_{\ell=1}^{N_{\mathrm{B}}} \omega(\ell) \leq \mathscr{P}_{\mathrm{S}}, \quad \omega(\ell) > 0, \\
&\quad\text{and} \sum_{i=1}^{N_{\mathrm{C}}}\sum_{\ell=1}^{N_{\mathrm{S}}} \delta_i(\ell)\left(\boldsymbol{U}_{\mathrm{h,right},i}^{\mathrm{H}}\,\boldsymbol{U}_{\mathrm{h,right},i}\right)^{-1}_{\ell\ell} \\
&\quad\leq \mathscr{P}_{\mathrm{D}},\ \delta_i(\ell) > 0
\end{aligned} \tag{17}$$

where we have defined $\boldsymbol{\omega} \triangleq [\omega(1), \omega(2), \ldots, \omega(N_{\mathrm{B}})]^T \in \mathbb{R}^{N_{\mathrm{B}}}$, $\boldsymbol{\delta}_i \triangleq [\delta_i(1), \delta_i(2), \ldots, \delta_i(N_{\mathrm{S}})] \in \mathbb{R}^{N_{\mathrm{S}}}$, for $i \in \{1, 2, \ldots, N_{\mathrm{C}}\}$,

$$f_0\left(\boldsymbol{\omega}, \{\boldsymbol{\delta}_i\}_{i=1}^{N_{\mathrm{C}}}\right) \triangleq \sum_{\ell=1}^{N_{\mathrm{B}}} \frac{1}{\omega(\ell)\,\lambda_{\mathrm{h}}^2(\Delta N + \ell)\displaystyle\sum_{i=1}^{N_{\mathrm{C}}} \delta_i(\Delta N + \ell)} \tag{18}$$

with $\Delta N \triangleq N_{\mathrm{S}} - N_{\mathrm{B}} \geq 0$. All the inequality constraints in (17) are linear. However, it is shown in Appendix B that, when $N_{\mathrm{C}} > 1$, the cost function (18) is the sum of N_{B} functions that are neither strictly convex nor strictly concave on $\mathbb{R}_+^{n+1}$. Hence, trying to solve (17) with the available optimization tools leads to poor performance in multiple-relay WSNs.

B. DESIGN BASED ON THE SVD OF THE INDIVIDUAL FIRST-HOP CHANNEL MATRICES

A simple design can be developed by setting $\boldsymbol{F}_i = \boldsymbol{O}_{N_{\mathrm{R}}\times N_{\mathrm{R}}}$, for each $i \in \{1, 2, \ldots, N_{\mathrm{C}}\} - i^\star$. Basically, such a choice leads to a *single-relay selection* scheme [14], which imposes that only one relay (i.e., that for $i = i^\star$) is recruited to transmit and all the remaining ones keep silent in the second hop.

Herein, we assume that $\boldsymbol{H}_i$ is full-column rank, i.e., $\mathrm{rank}(\boldsymbol{H}_i) = N_{\mathrm{S}} \leq N_{\mathrm{R}}$, for each $i \in \{1, 2, \ldots, N_{\mathrm{C}}\}$. Let

$$\boldsymbol{U}_{\mathrm{h},i}\,[\boldsymbol{O}_{N_{\mathrm{S}}\times(N_{\mathrm{R}}-N_{\mathrm{S}})}, \boldsymbol{\Lambda}_{\mathrm{h},i}]^{\mathrm{T}}\,\boldsymbol{V}_{\mathrm{h},i}^{\mathrm{H}} \tag{19}$$

be the SVD of $\boldsymbol{H}_i$, where

$$\boldsymbol{\Lambda}_{\mathrm{h},i} \triangleq \mathrm{diag}[\lambda_{\mathrm{h},i}(1), \lambda_{\mathrm{h},i}(2), \ldots, \lambda_{\mathrm{h},i}(N_{\mathrm{S}})] \tag{20}$$

contains the singular values of $\boldsymbol{H}_i$, arranged in increasing order, and the unitary matrices $\boldsymbol{U}_{\mathrm{h},i} \in \mathbb{C}^{N_{\mathrm{R}}\times N_{\mathrm{R}}}$ and $\boldsymbol{V}_{\mathrm{h},i} \in \mathbb{C}^{N_{\mathrm{S}}\times N_{\mathrm{S}}}$ collect the corresponding left and right singular vectors, respectively. In this case, one has $\boldsymbol{A} = \boldsymbol{H}_{i^\star}^{\mathrm{H}}\boldsymbol{F}_{i^\star}^{\mathrm{H}}\boldsymbol{F}_{i^\star}\,\boldsymbol{H}_{i^\star}$ and, by substituting the SVD of $\boldsymbol{H}_{i^\star}$ in this matrix equation, one has that the diagonalization of $\boldsymbol{F}_0^{\mathrm{H}}\boldsymbol{A}\,\boldsymbol{F}_0$ is ensured by

$$\boldsymbol{F}_0 = \boldsymbol{V}_{\mathrm{h},i^\star,\mathrm{right}}\,\boldsymbol{\Omega}^{1/2} \tag{21}$$

$$\boldsymbol{F}_{i^\star} = \boldsymbol{Q}_{i^\star}\boldsymbol{\Delta}^{1/2}\,\boldsymbol{U}_{\mathrm{h},i^\star,\mathrm{right}}^{\mathrm{H}} \tag{22}$$

where $\boldsymbol{U}_{\mathrm{h},i^\star,\mathrm{right}} \in \mathbb{C}^{N_{\mathrm{R}}\times N_{\mathrm{S}}}$ and $\boldsymbol{V}_{\mathrm{h},i^\star,\mathrm{right}} \in \mathbb{C}^{N_{\mathrm{S}}\times N_{\mathrm{B}}}$ contain the N_{S} and N_{B} rightmost columns from $\boldsymbol{U}_{\mathrm{h},i^\star}$ and $\boldsymbol{V}_{\mathrm{h},i^\star}$, respectively, $\boldsymbol{Q}_{i^\star} \in \mathbb{C}^{N_{\mathrm{R}}\times N_{\mathrm{S}}}$ is an arbitrary matrix obeying $\boldsymbol{Q}_{i^\star}^{\mathrm{H}}\boldsymbol{Q}_{i^\star} = \boldsymbol{I}_{N_{\mathrm{S}}}$, $\boldsymbol{\Omega}$ has been defined in Subsection III-A, and $\boldsymbol{\Delta} \triangleq \mathrm{diag}[\delta(1), \delta(2), \ldots, \delta(N_{\mathrm{S}})]$. To fully specify the solution of (13) in the case of single-relay selection, optimization of $\boldsymbol{\Omega}$, $\boldsymbol{\Delta}$, and $i^\star$ is accomplished in two steps.

First, for a given $i^\star \in \{1, 2, \ldots, N_{\mathrm{C}}\}$, by substituting (21) and (22) in (13), one obtains the scalar optimization problem with linear inequality constraints:

$$\begin{aligned}
&\min_{\boldsymbol{\omega},\boldsymbol{\delta}} f_1\left(i^\star, \boldsymbol{\omega}, \boldsymbol{\delta}\right) \\
&\text{s.to} \sum_{\ell=1}^{N_{\mathrm{B}}} \omega(\ell) \leq \mathscr{P}_{\mathrm{S}}, \quad \omega(\ell) > 0, \\
&\quad\text{and} \sum_{\ell=1}^{N_{\mathrm{S}}} \delta(\ell) \leq \mathscr{P}_{\mathrm{D}}, \quad \delta(\ell) > 0
\end{aligned} \tag{23}$$

with

$$f_1\left(i^\star, \boldsymbol{\omega}, \boldsymbol{\delta}\right) \triangleq \sum_{\ell=N_{\mathrm{S}}-N_{\mathrm{B}}+1}^{N_{\mathrm{S}}} \frac{1}{\lambda_{\mathrm{h},i^\star}^2(\ell)\,\omega(\ell)\,\delta(\ell)} \tag{24}$$

where $\boldsymbol{\omega}$ has been previously defined in Subsection III-A and $\boldsymbol{\delta} \triangleq [\delta(1), \delta(2), \ldots, \delta(N_{\mathrm{S}})]^T \in \mathbb{R}^{N_{\mathrm{S}}}$. Since $f_1\,(i^\star, \boldsymbol{\omega}, \boldsymbol{\delta})$ is a convex function (see Appendix B), the optimization problem (23) is convex and, thus, its solution $\boldsymbol{\omega}_{\mathrm{opt}}(i^\star)$ and $\boldsymbol{\delta}_{\mathrm{opt}}(i^\star)$

can be found by using efficient numerical techniques [39]. For instance, if one resorts to interior point methods, convergence arbitrarily close to the optimal solution is achieved in a number of iterations that is proportional to the logarithm of the problem dimension [40], with a complexity per iteration dictated by the cost M of computing a Newton direction [41].

Second, the optimal value i_{opt} of $i^\star$ is obtained as

$$i_{\text{opt}} \triangleq \arg \min_{i^\star \in \{1,2,\ldots,N_{\text{C}}\}} f_1\left(i^\star, \boldsymbol{\omega}_{\text{opt}}(i^\star), \boldsymbol{\delta}_{\text{opt}}(i^\star)\right) \tag{25}$$

which allows one to single out the best relay among the N_{C} available ones. The solution of (25) can be obtained by solving (23) for each $i^\star \in \{1, 2, \ldots, N_{\text{C}}\}$, with an overall complexity $\mathscr{O}\left[N_{\text{C}}\, M \log(N_{\text{B}} + N_{\text{S}})\right]$.

In the SISO configuration, i.e., when $N_{\text{B}} = N_{\text{S}} = N_{\text{R}} = N_{\text{D}} = 1$, and when there is no precoding at the source, i.e., $\boldsymbol{F}_0 \equiv f_0 = \sqrt{\mathscr{P}_{\text{S}}}$, one gets $\boldsymbol{F}_{i^\star} = \text{diag}(0, \ldots, 0, f_{i^\star}, 0, \ldots, 0)$ and (25) boils down to $i_{\text{opt}} \triangleq \arg\max_{i^\star \in \{1,2,\ldots,N_{\text{C}}\}}\{|h_i|^2\}$, with h_i denoting the channel coefficients between the source and the ith relay. According to [42], such a scheme has a full diversity order equal to N_{C}. However, as we will see in Section IV, such a design suffers from a diversity loss in a MIMO setting, i.e., when $N_{\text{S}}, N_{\text{R}}, N_{\text{D}} > 1$.

C. DESIGN BASED ON THE SVD OF ROW-BASED PARTITIONS OF THE COMPOSITE FIRST-HOP CHANNEL MATRIX

In the considered cooperative MIMO WSN, there are N_{C} relays equipped with N_{R} antennas, which amounts to a total number of $N_{\text{C}}\, N_{\text{R}}$ distributed antennas. Here, we propose to choose the best $N_{\text{B}} = N_{\text{S}}$ antennas out of the $N_{\text{C}}\, N_{\text{R}}$ ones.[4] Such antennas can either be physically located on a single relay, or be spatially distributed over different relays, thus accomplishing a joint antenna-and-relay selection scheme.

Let $\mathscr{S} \triangleq \{(n, i), \forall n \in \{1, 2, \ldots, N_{\text{R}}\}, \forall i \in \{1, 2, \ldots, N_{\text{C}}\}\}$ collect all the $N_{\text{C}}\, N_{\text{R}}$ antenna elements in the network, with the generic (ordered) pair (n, i) uniquely identifying the nth antenna located on the ith relay. The number of distinct subsets of $\mathscr{S}$ that have exactly N_{B} elements is given by the binomial coefficient $Q \triangleq \binom{N_{\text{C}} N_{\text{R}}}{N_{\text{B}}}$.[5] By excluding the trivial choice $\emptyset$ and the degenerate case $\mathscr{S}$ (discussed in Subsection III-A), we denote with

$$\mathscr{S}^{(q)} \triangleq \left\{(n_1^{(q)}, i_1^{(q)}), (n_2^{(q)}, i_2^{(q)}), \ldots, (n_{N_{\text{B}}}^{(q)}, i_{N_{\text{B}}}^{(q)})\right\} \tag{26}$$

the selected subset of $\mathscr{S}$, with $q \in \{1, 2, \ldots, Q-2\}$, obeying $\mathscr{S}^{(q_1)} \neq \mathscr{S}^{(q_2)}$ for $q_1 \neq q_2$. Additionally, we use the notation $N_i^{(q)} \in \{0, 1, \ldots, N_{\text{B}}\}$ to indicate the number of pairs of $\mathscr{S}^{(q)}$ having the same second entry: in other words, $N_i^{(q)}$ represents the number of antennas activated on the ith relay according to the qth selection. It results that $\sum_{i=1}^{N_{\text{C}}} N_i^{(q)} = N_{\text{B}}$.

The selected antennas generate a first-hop channel matrix

$$\boldsymbol{H}^{(q)} \triangleq [(\boldsymbol{H}_1^{(q)})^{\text{T}}, (\boldsymbol{H}_2^{(q)})^{\text{T}}, \ldots, (\boldsymbol{H}_{N_{\text{C}}}^{(q)})^{\text{T}}]^{\text{T}} \in \mathbb{C}^{N_{\text{B}} \times N_{\text{B}}} \tag{27}$$

and a relaying matrix

$$\boldsymbol{F}^{(q)} \triangleq \text{diag}(\boldsymbol{F}_1^{(q)}, \boldsymbol{F}_2^{(q)}, \ldots, \boldsymbol{F}_{N_{\text{C}}}^{(q)}) \in \mathbb{C}^{N_{\text{B}} \times N_{\text{B}}} \tag{28}$$

with $\boldsymbol{H}_i^{(q)} \in \mathbb{C}^{N_i^{(q)} \times N_{\text{B}}}$ and $\boldsymbol{F}_i^{(q)} \in \mathbb{C}^{N_i^{(q)} \times N_i^{(q)}}$. By convention, if $N_i^{(q)} = 0$, then $\boldsymbol{H}_i^{(q)}$ and $\boldsymbol{F}_i^{(q)}$ are empty matrices.

With reference to the qth selection, we formulate a new optimization problem, for $q \in \{1, 2, \ldots, Q-2\}$, which is formally obtained from (13) by replacing $\boldsymbol{H}$ and $\boldsymbol{F}$ with $\boldsymbol{H}^{(q)}$ and $\boldsymbol{F}^{(q)}$, respectively, whose cost function achieves its minimum value if $\boldsymbol{F}_0^{\text{H}}\, \boldsymbol{A}^{(q)}\, \boldsymbol{F}_0$ is diagonal (see Lemma 2), with $\boldsymbol{A}^{(q)} \triangleq (\boldsymbol{H}^{(q)})^{\text{H}}\, (\boldsymbol{F}^{(q)})^{\text{H}}\, \boldsymbol{F}^{(q)}\, \boldsymbol{H}^{(q)}$. For $i \in \{1, 2, \ldots, N_{\text{C}}\}$, let $\boldsymbol{H}_i^{(q)} = \boldsymbol{U}_{\text{h},i}^{(q)}\, [\boldsymbol{O}_{N_i^{(q)} \times (N_{\text{B}} - N_i^{(q)})}, \boldsymbol{\Lambda}_{\text{h},i}^{(q)}]\, (\boldsymbol{V}_{\text{h},i}^{(q)})^{\text{H}}$ be the SVD of the (nonempty) matrix $\boldsymbol{H}_i^{(q)}$, which is assumed to be full-row rank, i.e., $\text{rank}(\boldsymbol{H}_i^{(q)}) = N_i^{(q)}$, where $\boldsymbol{U}_{\text{h},i}^{(q)} \in \mathbb{C}^{N_i^{(q)} \times N_i^{(q)}}$ and $\boldsymbol{V}_{\text{h},i}^{(q)} \in \mathbb{C}^{N_{\text{B}} \times N_{\text{B}}}$ are unitary, and the diagonal matrix $\boldsymbol{\Lambda}_{\text{h},i}^{(q)} \triangleq \text{diag}[\lambda_{\text{h},i}^{(q)}(1), \lambda_{\text{h},i}^{(q)}(2), \ldots, \lambda_{\text{h},i}^{(q)}(N_i^{(q)})]$ collects the corresponding nonzero singular values arranged in increasing order. In this case, the diagonalization of $\boldsymbol{F}_0^{\text{H}}\, \boldsymbol{A}^{(q)}\, \boldsymbol{F}_0$ can be obtained by resorting to the following structures

$$\boldsymbol{F}_0 = [(\boldsymbol{V}_{\text{h,right}}^{(q)})^{\text{H}}]^{-1}\, \boldsymbol{\Omega}^{1/2} \tag{29}$$

$$\boldsymbol{F}_i^{(q)} = \boldsymbol{Q}_i\, \boldsymbol{\Delta}_i^{1/2}\, (\boldsymbol{U}_{\text{h},i}^{(q)})^{\text{H}} \tag{30}$$

where

$$\boldsymbol{V}_{\text{h,right}}^{(q)} \triangleq [\boldsymbol{V}_{\text{h,right},1}^{(q)}, \boldsymbol{V}_{\text{h,right},2}^{(q)}, \ldots, \boldsymbol{V}_{\text{h,right},N_{\text{C}}}^{(q)}] \in \mathbb{C}^{N_{\text{B}} \times N_{\text{B}}} \tag{31}$$

with $\boldsymbol{V}_{\text{h,right},i}^{(q)} \in \mathbb{C}^{N_{\text{B}} \times N_i^{(q)}}$ gathering the $N_i^{(q)}$ rightmost columns from $\boldsymbol{V}_{\text{h},i}^{(q)}$, $\boldsymbol{Q}_i \in \mathbb{C}^{N_i^{(q)} \times N_i^{(q)}}$ is an arbitrary unitary matrix, $\boldsymbol{\Omega}$ and $\boldsymbol{\Delta}_i$ have been defined in Subsection III-A.

To optimize $\boldsymbol{\Omega}$, $\boldsymbol{\Delta}_i$, and q, we resort to a two-step procedure as in the previous subsection. By substituting (29)–(30) in (13) (with $\boldsymbol{H}^{(q)}$ and $\boldsymbol{F}^{(q)}$ in lieu of $\boldsymbol{H}$ and $\boldsymbol{F}$, respectively), for a given $q \in \{1, 2, \ldots, Q-2\}$, one gets the convex optimization problem (see Appendix B) with linear inequality constraints:

$$\begin{aligned} &\min_{\boldsymbol{\omega}, \{\boldsymbol{\delta}_i\}_{i=1}^{N_{\text{C}}}} f_2\left(q, \boldsymbol{\omega}, \{\boldsymbol{\delta}_i\}_{i=1}^{N_{\text{C}}}\right) \\ &\quad \text{s.to} \sum_{\ell=1}^{N_{\text{B}}} \omega(\ell) \left[(\boldsymbol{V}_{\text{h,right}}^{(q)})^{\text{H}}\, \boldsymbol{V}_{\text{h,right}}^{(q)}\right]_{\ell\ell}^{-1} \leq \mathscr{P}_{\text{S}}, \quad \omega(\ell) > 0, \\ &\qquad \text{and} \sum_{i=1}^{N_{\text{C}}} \sum_{\ell=1}^{N_{\text{S}}} \delta_i(\ell) \leq \mathscr{P}_{\text{D}}, \quad \delta_i(\ell) > 0 \end{aligned} \tag{32}$$

with

$$f_2\left(q, \boldsymbol{\omega}, \{\boldsymbol{\delta}_i\}_{i=1}^{N_{\text{C}}}\right) \triangleq \sum_{i=1}^{N_{\text{C}}} \sum_{\ell=1}^{N_{\text{B}}} \frac{1}{\left[\lambda_{\text{h},i}^{(q)}(\ell)\right]^2 \omega_i(\ell)\, \delta_i(\ell)} \tag{33}$$

where $\omega_i(\ell) \triangleq \omega\left(\sum_{m=1}^{i-1} N_m^{(q)} + \ell\right)$, whereas $\boldsymbol{\omega}$ and $\boldsymbol{\delta}_i$ have been defined in Subsection III-A. Similarly to problem (23), the solution $\boldsymbol{\omega}_{\text{opt}}(q)$ and $\{\boldsymbol{\delta}_{i,\text{opt}}(q)\}_{i=1}^{N_{\text{C}}}$ of (32) can be found by

[4]Our design can be simply extended to the case $N_{\text{S}} \geq N_{\text{B}}$.

[5]The empty set $\emptyset$ and the set $\mathscr{S}$ are considered as subsets of $\mathscr{S}$ as well.

using, e.g., interior point methods [39]. Finally, the best value $q_{\rm opt}$ of q is found by solving

$$q_{\rm opt} \triangleq \arg \min_{q \in \{1,2,\ldots,Q-2\}} f_2\left(q, \boldsymbol{\omega}_{\rm opt}(q), \{\boldsymbol{\delta}_{i,\rm opt}(q)\}_{i=1}^{N_{\rm C}}\right) \quad (34)$$

which determines the best $N_{\rm B}$-dimensional subset of the available $N_{\rm C}\,N_{\rm R}$ antennas. The solution of (34) can be obtained by solving (32) for each $q \in \{1, 2, \ldots, Q-2\}$, with an overall complexity $\mathscr{O}\left[(Q-2)\,M\,\log(N_{\rm B} + N_{\rm B}\,N_{\rm C})\right]$, which is larger than that required to select the best relay (see Subsection III-B), especially for large number of relays.

When $N_{\rm B} = N_{\rm S} = N_{\rm R} = 1$, the optimization problems (23)-(25) and (32)-(34) yield the same solution and, thus, the design (32)-(34) exhibits full diversity order $N_{\rm C}$, too. However, we will show in the next section that, when $N_{\rm B}, N_{\rm S}, N_{\rm R} > 1$ (MIMO WSN), the proposed joint antenna-and-relay selection scheme ensures a significant performance improvement with respect to single-relay selection, in terms of both diversity order and coding gain.

IV. NUMERICAL RESULTS

In this section, to assess the performance of the considered P-CSI designs, we present the results of Monte Carlo computer simulations, aimed at evaluating the ASEP of the corresponding cooperative systems, transmitting quadrature phase-shift-keying (QPSK) symbols. We set $N \triangleq N_{\rm B} = N_{\rm S} = N_{\rm R} = N_{\rm D}$ in all the forthcoming examples, with $N \in \{1, 2, 3\}$. We also assume that $\mathscr{P}_{\rm S} = \widetilde{\mathscr{P}}_{\rm D} = \mathscr{P}_{\rm tot}$. Consequently, the SNR is defined as SNR $\triangleq \mathscr{P}_{\rm tot}$, which measures the per-antenna link quality of both the first- and second-hop transmissions. Besides the single-relay selection method described in Subsection III-B, referred to as "1-R Selection", and the joint antenna-and-relay selection scheme developed in Subsection III-C, referred to as "JAR Selection", we also report the performance of [26, CSI Assumptions I and II] in the case of single-antenna nodes (i.e., $N = 1$) and that of [30] for both single- and multiple-antennas nodes (i.e., $N \in \{2, 3\}$). As a reference lower bound, we additionally include in all the plots the ASEP curves of the F-CSI design proposed in [15], whose design relies on the additional knowledge of the ith second-hop channel matrix $\boldsymbol{G}_i$ at the ith relay, for $i \in \{1, 2, \ldots, N_{\rm C}\}$. This F-CSI method exhibits a theoretical diversity order equal to $N_{\rm C}\,N_{\rm R} - N_{\rm B} + 1$ [15].

The ASEP has been evaluated by carrying out 10^3 independent Monte Carlo trials, with each run using independent sets of channel realizations and noise, and an independent record of 10^6 source symbols.

A. EXAMPLE 1: SINGLE-ANTENNA NODES

We report in Figs. 1 and 2 the ASEP performance of the considered schemes as a function of the SNR, for single-antenna nodes (i.e., $N = 1$) and two different values of the number of relays $N_{\rm C} \in \{2, 3\}$. We would like to remember that, in the case of $N = 1$, the two approaches "1-R Selection" and "JAR Selection" are equivalent and, thus, only the performance of the "1-R Selection" method are reported.

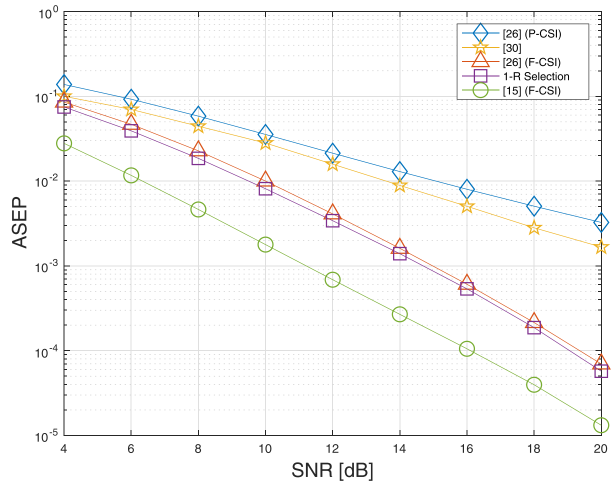


FIGURE 1. ASEP versus SNR (Example 1: $N = 1$ and $N_{\rm C} = 2$).

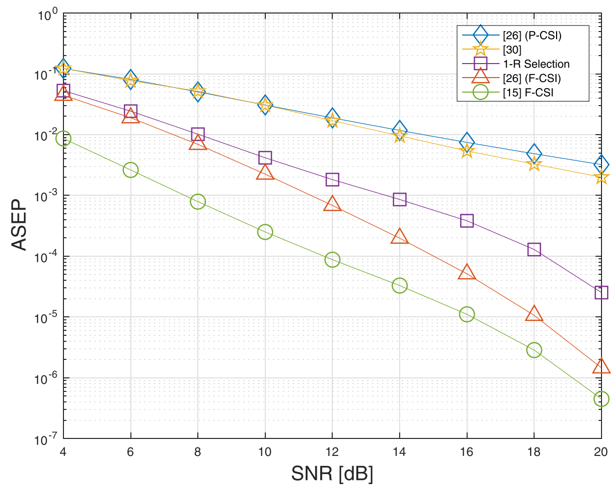


FIGURE 2. ASEP versus SNR (Example 1: $N = 1$ and $N_{\rm C} = 3$).

Results clearly show that no diversity is achieved by [26] (CSI Assumption II corresponding to P-CSI) and [30], irrespective of the number of relays. On the other hand, the "1-R Selection" scheme exhibits the same diversity order of the F-CSI methods proposed in [26] (CSI Assumption I) and [15], which linearly increases with $N_{\rm C}$. This fact allows the "1-R Selection" design to significantly outperform both [26] (P-CSI) and [30], which rely on the same amount of CSI. Remarkably, the "1-R Selection" scheme performs comparably to [26] (F-CSI) in the case of $N_{\rm C} = 2$ relays. Compared to single-relay selection, the performance improvement of the F-CSI – arising from the additional instantaneous knowledge of the second-hop matrix $\boldsymbol{G}$ – becomes more and more apparent when the number of relays $N_{\rm C}$ increases.

B. EXAMPLE 2: MULTIPLE-ANTENNA NODES

Figs. 3, 4, 5, and 6 show the ASEP performance of the considered designs as a function of the SNR, for two different

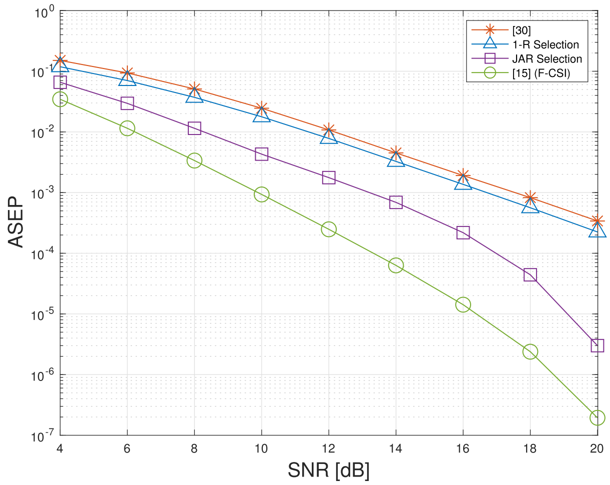


FIGURE 3. **ASEP versus SNR (Example 2: $N = 2$ and $N_C = 2$).**

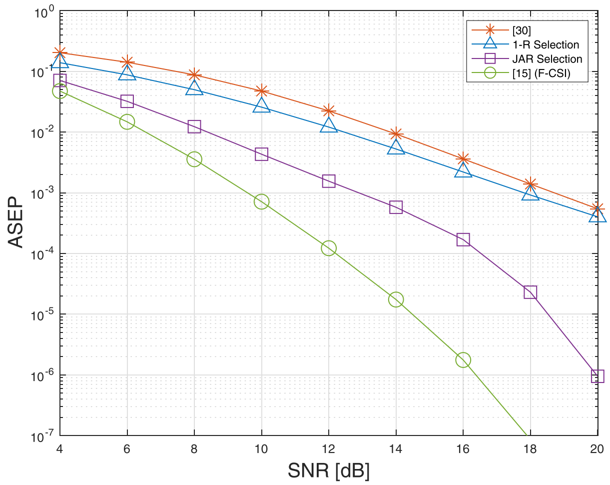


FIGURE 4. **ASEP versus SNR (Example 2: $N = 2$ and $N_C = 3$).**

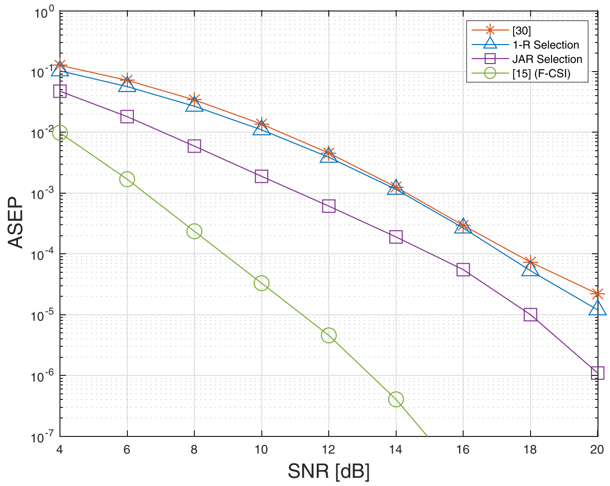


FIGURE 5. **ASEP versus SNR (Example 2: $N = 3$ and $N_C = 2$).**

multi-antenna configurations $N \in \{2, 3\}$ and two different values of the number of relays $N_C \in \{2, 3\}$, respectively.

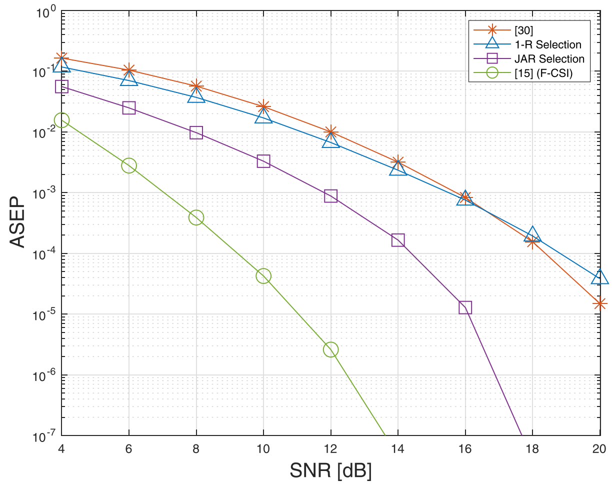


FIGURE 6. **ASEP versus SNR (Example 2: $N = 3$ and $N_C = 3$).**

It is apparent from these plots that, in a multi-antenna deployment, the "1-R Selection" approach and [30] perform comparably, both exhibiting a diversity loss with respect to the F-CSI design [15]. As claimed, especially in the high SNR regime, the proposed "JAR Selection" design significantly outperforms both the "1-R Selection" scheme and [30], under the same amount of P-CSI. Such a performance gap remarkably scales up as the number of antennas at the nodes increases from $N = 2$ to $N = 3$. Interestingly, the diversity order of the "JAR Selection" scheme increases with N_C, as in [15] which, however, requires F-CSI.

V. CONCLUSION

We studied the problem of designing multi-relay AF cooperative WSNs, based on the knowledge of the instantaneous values of the first-hop MIMO channel matrix and statistical characterization of the second-hop one (partial CSI scenario). In this case, antenna/relay selection schemes arise necessarily to formulate mathematically tractable design problems, which can be solved by using standard convex optimization tools. We have shown that, in a MIMO setting, the selection of the best relay is suboptimal and large performance improvements can be obtained by selecting the best antennas distributed over multiple relays. Numerical simulations have shown that the proposed joint antenna-and-relay selection approach significantly outperforms existing schemes, which exploit the same amount of P-CSI.

APPENDIX A
PROOF OF LEMMA 1

Preliminarily, we remember that $\boldsymbol{C} = \boldsymbol{G}\boldsymbol{F}\boldsymbol{H}\boldsymbol{F}_0$ is full-column rank if **a1)** and **a2)** hold. It can be shown (see, e.g., [43]) that, conditioned on $\boldsymbol{H}$, the kth diagonal entry $\{(\boldsymbol{C}^{\mathrm{H}}\boldsymbol{C})^{-1}\}_{kk}$ of the matrix $(\boldsymbol{C}^{\mathrm{H}}\boldsymbol{C})^{-1}$ follows an inverse-Gamma distribution, with shape parameter $\alpha \triangleq N_D - N_B + 1$ and scale parameter $\beta_k \triangleq 1/\{\boldsymbol{R}^{-1}\}_{kk}$, where $\boldsymbol{R}$ is defined in the lemma statement. Thus, the probability density function of the random variable $\{(\boldsymbol{C}^{\rm H}\boldsymbol{C})^{-1}\}_{kk}$, given $\boldsymbol{H}$, reads as

$$p_k(x) = \frac{1}{\Gamma(\alpha)\,\beta_k^{\alpha}}\, x^{-\alpha-1}\, e^{-\frac{1}{x\beta_k}} \tag{35}$$

where the gamma function $\Gamma(\alpha) = (\alpha-1)!$ since $N_{\rm D} - N_{\rm B}$ is a non-negative integer number [44]. Therefore, one has

$$\begin{aligned}
&\mathbb{E}_{\boldsymbol{G}}\left[\operatorname{tr}\left(\boldsymbol{C}^{\rm H}\boldsymbol{C}\right)^{-1} \,|\, \boldsymbol{H}\right] \\
&= \sum_{k=1}^{N_{\rm B}} \mathbb{E}_{\boldsymbol{G}}\left[\left\{\left(\boldsymbol{C}^{\rm H}\boldsymbol{C}\right)^{-1}\right\}_{kk} \,|\, \boldsymbol{H}\right] \\
&= \frac{1}{\Gamma(\alpha)} \sum_{k=1}^{N_{\rm B}} \frac{1}{\beta_k^{\alpha}} \left(\lim_{\delta\to 0}\int_{\delta}^{+\infty} x^{-\alpha} e^{-\frac{1}{x\beta_k}}\,{\rm d}x\right).
\end{aligned} \tag{36}$$

After some calculations, eq. (36) can be rewritten as

$$\begin{aligned}
&\mathbb{E}_{\boldsymbol{G}}\left[\operatorname{tr}\left(\boldsymbol{C}^{\rm H}\boldsymbol{C}\right)^{-1} \,|\, \boldsymbol{H}\right] \\
&= \sum_{k=1}^{N_{\rm B}} \frac{\beta_k^{-1}}{\Gamma(\alpha)} \lim_{\delta\to 0} \gamma(\alpha-1, (\delta\,\beta_k)^{-1}) \\
&= \frac{\Gamma(\alpha-1)}{\Gamma(\alpha)} \sum_{k=1}^{N_{\rm B}} \beta_k^{-1} = \frac{1}{\alpha-1}\sum_{k=1}^{N_{\rm B}} \{\boldsymbol{R}^{-1}\}_{kk} = \frac{\operatorname{tr}(\boldsymbol{R}^{-1})}{N_{\rm D}-N_{\rm B}}
\end{aligned} \tag{37}$$

where we have exploited the definition of the incomplete gamma function $\gamma(s,x) \triangleq \int_0^x t^{s-1}e^{-t}{\rm d}t$ [44] and its asymptotic property $\Gamma(s) = \lim_{x\to+\infty}\gamma(s,x)$.

APPENDIX B
HESSIAN OF THE COST FUNCTION (15)

Let us check convexity of a generic summand of the cost function (18). To this end, it is sufficient to study the multivariate function

$$f(x, y_1, y_2, \ldots, y_n) \triangleq \frac{1}{A\,x\,(y_1 + y_2 + \cdots + y_n)} \tag{38}$$

with $A > 0$, $x > 0$, and $y_i > 0$, for each $i \in \{1, 2, \ldots, n\}$. The domain of f is therefore given by $\mathbb{R}_+^{n+1}$, which is a convex set. The function f is twice differentiable over its domain. It is noteworthy that, when $n = 1$, the function (38) ends up to a generic summand of (23) or (32).

Let us calculate the Hessian matrix $\nabla^2 f \in \mathbb{R}^{(n+1)\times(n+1)}$, whose entries are the second-order partial derivatives of f at $(x, y_1, y_2, \ldots, y_n) \in \mathbb{R}_+^{n+1}$, i.e.,

$$\{\nabla^2 f\}_{ij} = \begin{cases} \dfrac{\partial^2}{\partial x^2} f, & \text{for } i = j = 1\,; \\ \dfrac{\partial^2}{\partial x\,\partial y_j} f, & \text{for } i = 1 \text{ and } j \in \{2, 3, \ldots, n\} \\ & \text{for } j = 1 \text{ and } i \in \{2, 3, \ldots, n\}\,; \\ \dfrac{\partial^2}{\partial y_i\,\partial y_j} f, & \text{for } i, j \in \{2, 3, \ldots, n\}\,. \end{cases} \tag{39}$$

We recall that the function f is convex [concave] if and only if the Hessian matrix $\nabla^2 f$ is positive [negative] semidefinite for all the points belonging to its domain.

Using standard calculus concepts, it can be verified that

$$\frac{\partial^2}{\partial x^2} f = \frac{2}{A\,x^3\,(y_1 + y_2 + \cdots + y_n)} \tag{40}$$

$$\frac{\partial^2}{\partial x\,\partial y_j} f = \frac{1}{A\,x^2\,(y_1 + y_2 + \cdots + y_n)^2} \tag{41}$$

$$\frac{\partial^2}{\partial y_i\,\partial y_j} f = \frac{2}{A\,x\,(y_1 + y_2 + \cdots + y_n)^3}\,. \tag{42}$$

We note that all the entries of $\nabla^2 f$ are nonnegative on $\mathbb{R}_+^{n+1}$. In the particular case of $n = 1$, it is readily seen that the determinant of $\nabla^2 f \in \mathbb{R}^{2\times 2}$ is given by

$$\det(\nabla^2 f) = \frac{3}{A^2\,x^4\,y_1^4} > 0 \tag{43}$$

which shows that, when $n = 1$, f is a strictly convex function on $\mathbb{R}_+^2$. Therefore, since the sum of convex functions is convex [39], the cost functions (23) or (32) are convex.

On the other hand, when $n > 1$, by resorting to the Laplacian determinant expansion by minors, it results that

$$\det(\nabla^2 f) = \sum_{j=1}^{n+1} (-1)^{j+1}\,\{\nabla^2 f\}_{1j}\,\boldsymbol{M}_{1j} \tag{44}$$

where $\boldsymbol{M}_{1\,j} \in \mathbb{R}^{n\times n}$ is a so-called minor of $\nabla^2 f$, obtained by taking the determinant of $\nabla^2 f$ with row 1 and column j crossed out. It can be verified that $\boldsymbol{M}_{1\,j}$ is zero, for each $j \in \{1, 2, \ldots, n+1\}$. Thus, the determinant of $\nabla^2 f$ is zero at each point belonging to the domain of f if $n > 1$. This is sufficient to infer that $\nabla^2 f$ is neither positive nor negative definitive, which implies in its turn that, when $n > 1$, f is neither strictly convex nor strictly concave on $\mathbb{R}_+^{n+1}$.

REFERENCES

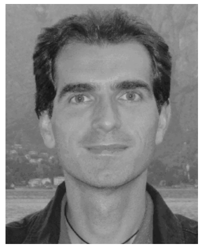
[1] J. N. Laneman, D. N. C. Tse, and G. W. Wornell, “Cooperative diversity in wireless networks: Efficient protocols and outage behavior,” *IEEE Trans. Inf. Theory*, vol. 50, no. 12, pp. 3062–3080, Dec. 2004.
[2] R. Yuan, T. Zhang, J. Huang, and L. Sun, “Opportunistic cooperation and optimal power allocation for wireless sensor networks,” *IEEE Trans. Consum. Electron.*, vol. 56, no. 3, pp. 1898–1904, Aug. 2010.
[3] S. Cui, A. J. Goldsmith, and A. Bahai, “Energy-efficiency of MIMO and cooperative MIMO techniques in sensor networks,” *IEEE J. Sel. Areas Commun.*, vol. 22, no. 6, pp. 1089–1098, Aug. 2004.
[4] W. Fang, F. Liu, F. Yang, L. Shu, and S. Nishio, “Energy-efficient cooperative communication for data transmission in wireless sensor networks,” *IEEE Trans. Consum. Electron.*, vol. 56, no. 4, pp. 2185–2192, Nov. 2010.
[5] D. N. Nguyen and M. Krunz, “Cooperative MIMO in wireless networks: Recent developments and challenges,” *IEEE Netw.*, vol. 27, no. 4, pp. 48–54, Jul./Aug. 2013.
[6] B. Rankov and A. Wittneben, “Spectral efficient protocols for half-duplex fading relay channels,” *IEEE J. Sel. Areas Commun.*, vol. 25, no. 2, pp. 379–389, Feb. 2007.
[7] A. S. Behbahani, R. Merched, and A. M. Eltawil, “Optimizations of a MIMO relay network,” *IEEE Trans. Signal Process.*, vol. 56, no. 10, pp. 5062–5073, Oct. 2008.
[8] A. Toding, M. Khandaker, and Y. Rong, “Joint source and relay optimization for parallel MIMO relay networks,” *EURASIP J. Adv. Signal Process.*, vol. 2012, no. 1, p. 174, 2012.

[9] K. T. Truong, P. Sartori, and R. W. Heath, Jr., "Cooperative algorithms for MIMO amplify-and-forward relay networks," *IEEE Trans. Signal Process.*, vol. 61, no. 5, pp. 1272–1287, Mar. 2013.
[10] C. Zhao and B. Champagne, "Joint design of multiple non-regenerative MIMO relaying matrices with power constraints," *IEEE Trans. Signal Process.*, vol. 61, no. 19, pp. 4861–4873, Oct. 2013.
[11] K. Lee, J. Yang, H. M. Kwon, H. Park, and Y. H. Lee, "Closed form of optimum cooperative distributed relay amplifying matrix," *IEEE Trans. Wireless Commun.*, vol. 13, no. 5, pp. 2741–2755, May 2014.
[12] A. Ikhlef and R. Schober, "Joint source-relay optimization for fixed receivers in multi-antenna multi-relay networks," *IEEE Trans. Wireless Commun.*, vol. 13, no. 1, pp. 62–74, Jan. 2014.
[13] C. Zhao and B. Champagne, "A unified approach to optimal transceiver design for nonregenerative MIMO relaying," *IEEE Trans. Veh. Technol.*, vol. 64, no. 7, pp. 2938–2951, Jul. 2015.
[14] D. Darsena, G. Gelli, and F. Verde, "Multiple-relay cooperative MIMO designs," in *Proc. Int. Symp. Wireless Commun. Syst. (ISWCS)*, Bologna, Italy, Aug. 2017, pp. 372–377.
[15] D. Darsena, G. Gelli, and F. Verde, "Design and performance analysis of multiple-relay cooperative MIMO networks," *J. Commun. Netw.*, vol. 21, no. 1, pp. 25–32, Feb. 2019.
[16] R. Zhang, Y.-C. Liang, C. C. Chai, and S. Cui, "Optimal beamforming for two-way multi-antenna relay channel with analogue network coding," *IEEE J. Sel. Areas Commun.*, vol. 27, no. 5, pp. 699–712, Jun. 2009.
[17] Q. Li and L. Yang, "Beamforming for cooperative secure transmission in cognitive two-way relay networks," *IEEE Trans. Inf. Forensics Security*, vol. 15, pp. 130–143, 2020.
[18] B. K. Chalise and L. Vandendorpe, "MIMO relay design for multipoint-to-multipoint communications with imperfect channel state information," *IEEE Trans. Signal Process.*, vol. 57, no. 7, pp. 2785–2796, Jul. 2009.
[19] Q. Li and L. Yang, "Robust optimization for energy efficiency in MIMO two-way relay networks with SWIPT," *IEEE Syst. J.*, vol. 14, no. 1, pp. 196–207, Mar. 2020.
[20] A. Bletsas, A. Khisti, D. P. Reed, and A. Lippman, "A simple cooperative diversity method based on network path selection," *IEEE J. Sel. Areas Commun.*, vol. 24, no. 3, pp. 659–672, Mar. 2006.
[21] Y. Zhao, R. Adve, and T. Lim, "Symbol error rate of selection amplify-and-forward relay systems," *IEEE Commun. Lett.*, vol. 10, no. 11, pp. 757–759, Nov. 2006.
[22] Y. Jing and H. Jafarkhani, "Single and multiple relay selection schemes and their achievable diversity orders," *IEEE Trans. Wireless Commun.*, vol. 8, no. 3, pp. 1414–1423, Mar. 2009.
[23] L. Sun, T. Zhang, L. Lu, and H. Niu, "Cooperative communications with relay selection in wireless sensor networks," *IEEE Trans. Consum. Electron.*, vol. 55, no. 2, pp. 513–517, May 2009.
[24] I. Krikidis, J. Thompson, S. Mclaughlin, and N. Goertz, "Amplify-and-forward with partial relay selection," *IEEE Commun. Lett.*, vol. 12, no. 4, pp. 235–237, Apr. 2008.
[25] J.-B. Kim and D. Kim, "Comparison of tightly power-constrained performances for opportunistic amplify-and-forward relaying with partial or full channel information," *IEEE Commun. Lett.*, vol. 13, no. 2, pp. 100–102, Feb. 2009.
[26] Z. Yi and I.-M. Kim, "Joint optimization of relay-precoders and decoders with partial channel side information in cooperative networks," *IEEE J. Sel. Areas Commun.*, vol. 25, no. 2, pp. 447–458, Feb. 2007.
[27] D. B. da Costa and S. Aissa, "End-to-end performance of dual-hop semi-blind relaying systems with partial relay selection," *IEEE Trans. Wireless Commun.*, vol. 8, no. 8, pp. 4306–4315, Aug. 2009.
[28] B. K. Chalise, Y. D. Zhang, and M. G. Amin, "Local CSI based full diversity achieving relay selection for amplify-and-forward cooperative systems," *IEEE Trans. Signal Process.*, vol. 61, no. 21, pp. 5165–5180, Nov. 2013.
[29] B. K. Chalise, L. Vandendorpe, Y. D. Zhang, and M. G. Amin, "Local CSI based selection beamforming for amplify-and-forward MIMO relay networks," *IEEE Trans. Signal Process.*, vol. 60, no. 5, pp. 2433–2446, May 2012.
[30] A. Danaee, M. Maleki, J. S. Kota, and H. R. Bahrami, "Relay and antenna selection in multi-antenna amplify-and-forward (AF) systems with partial channel state information," in *Proc. IEEE Int. Conf. Commun. (ICC)*, Ottawa, ON, Canada, Jun. 2012, pp. 4155–4159.
[31] D. P. Palomar, J. M. Cioffi, and M. A. Lagunas, "Joint Tx-Rx beamforming design for multicarrier MIMO channels: A unified framework for convex optimization," *IEEE Trans. Signal Process.*, vol. 51, no. 9, pp. 2381–2401, Sep. 2003.
[32] P. S. Elamvazhuthi, B. K. Dey, and S. Bhashyam, "An MMSE strategy at relays with partial CSI for a multi-layer relay network," *IEEE Trans. Signal Process.*, vol. 62, no. 2, pp. 271–282, Jan. 2014.
[33] R. Horn and C. Johnson, *Matrix Analysis*. New York, NY, USA: Cambridge Univ. Press, 1990.
[34] X. M. Yang, X. Q. Yang, and K. L. Teo, "A matrix trace inequality," *J. Math. Anal. Appl.*, vol. 263, no. 1, pp. 327–331, Nov. 2001.
[35] T. Sartenaer, J. Louveaux, and L. Vandendorpe, "Balanced capacity of wireline multiple access channels with individual power constraints," *IEEE Trans. Commun.*, vol. 56, no. 6, pp. 925–936, Jun. 2008.
[36] R. S. Cheng and S. Verdú, "Gaussian multiaccess channels with ISI: Capacity region and multiuser water-filling," *IEEE Trans. Inf. Theory*, vol. 39, no. 3, pp. 773–785, May 1993.
[37] S. M. Kay, *Fundamentals of Statistical Signal Processing: Estimation Theory*. Englewood Cliffs, NJ, USA: Prentice-Hall, 1993.
[38] A. Ben-Israel and T. N. E. Greville, *Generalized Inverses*. New York, NY, USA: Springer-Verlag, 2002.
[39] S. Boyd and L. Vandenberghe, *Convex Optimization*. New York, NY, USA: Cambridge Univ. Press, 2004.
[40] M. Colombo and J. Gondzio, "Further development of multiple centrality correctors for interior point methods," *Comput. Optim. Appl.*, vol. 41, pp. 277–305, Dec. 2008.
[41] F. Roger, *Practical Methods of Optimization*, 2nd ed. New York, NY, USA: Wiley, 1987.
[42] A. Ribeiro, X. Cai, and G. B. Giannakis, "Symbol error probabilities for general cooperative links," *IEEE Trans. Wireless Commun.*, vol. 4, no. 3, pp. 1264–1273, May 2005.
[43] P. Li, D. Paul, R. Narasimhan, and J. Cioffi, "On the distribution of SINR for the MMSE MIMO receiver and performance analysis," *IEEE Trans. Inf. Theory*, vol. 52, no. 1, pp. 271–286, Jan. 2006.
[44] M. Abramowitz and I. Stegun, *Handbook of Mathematical Functions*. New York, NY, USA: Dover, 1965.

DONATELLA DARSENA (Senior Member, IEEE) received the Dr.Eng. degree *(summa cum laude)* in telecommunications engineering and the Ph.D. degree in electronic and telecommunications engineering from the University of Napoli Federico II, Italy, in 2001 and 2005, respectively. From 2001 to 2002, she worked as an Embedded System Designer with the Telecommunications, Peripherals and Automotive Group, STMicroelectronics, Milano, Italy. Since 2005, she has been an Assistant Professor with the Department of Engineering, University of Napoli Parthenope, Italy. Her research interests include the broad area of signal processing for communications, with current emphasis on multicarrier modulation systems, space–time techniques for cooperative and cognitive communications, and the green communications for IoT. She has served as an Associate Editor for the IEEE Communications Letters from December 2016 to July 2019. Since August 2019, she has been a Senior Area Editor of the IEEE Communications Letters and an Associate Editor of IEEE Access since October 2018.

GIACINTO GELLI (Senior Member, IEEE) was born in Napoli, Italy, in July 1964. He received the Dr.Eng. degree *(summa cum laude)* in electronic engineering and the Ph.D. degree in computer science and electronic engineering from the University of Napoli Federico II, in 1990 and 1994, respectively.

From 1994 to 1998, he was an Assistant Professor with the Department of Information Engineering, Second University of Napoli. Since 1998, he has been with the Department of Electrical Engineering and Information Technology, University of Napoli Federico II, first as an Associate Professor, and since November 2006 as a Full Professor of telecommunications. He also held teaching positions at the University of Napoli Parthenope. His research interests include the broad area of signal and array processing for communications, with current emphasis on multicarrier modulation systems and space–time techniques for cooperative and cognitive communications systems.



FRANCESCO VERDE (Senior Member, IEEE) was born in Santa Maria Capua Vetere, Italy, in June 1974. He received the Dr.Eng. degree *(summa cum laude)* in electronic engineering from the Second University of Napoli, Italy, in 1998, and the Ph.D. degree in information engineering from the University of Napoli Federico II, in 2002. Since December 2002, he has been with the University of Napoli Federico II. He first served as an Assistant Professor of signal theory and mobile communications and, since December 2011, he has been an Associate Professor of telecommunications with the Department of Electrical Engineering and Information Technology. His research interests include orthogonal/non-orthogonal multiple-access techniques, space–time processing for cooperative/cognitive communications, wireless systems optimization, and software-defined networks.

Prof. Verde has been involved in several technical program committees of major IEEE conferences in signal processing and wireless communications. He has been an Associate Editor of the IEEE Transactions on Communications since 2017 and a Senior Area Editor of the IEEE Signal Processing Letters since 2018. He was an Associate Editor of the IEEE Transactions on Signal Processing from 2010 to 2014 and the IEEE Signal Processing Letters from 2014 to 2018, and the Guest Editor of the *EURASIP Journal on Advances in Signal Processing* in 2010 and *Sensors* (MDPI) in 2018.