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# Mathematical Modeling and a Multiswarm Collaborative Optimization Algorithm for Fuzzy Integrated Process Planning and Scheduling Problem

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Abstract: Considering both process planning and shop scheduling in manufacturing can fully utilize their complementarities, resulting in improved rationality of process routes and high-quality and efficient production. Hence, the study of Integrated Process Planning and Scheduling (IPPS) has become a hot topic in the current production field. However, when performing this integrated optimization, the uncertainty of processing time is a realistic key point that cannot be neglected. Thus, this paper investigates a Fuzzy IPPS (FIPPS) problem to minimize the maximum fuzzy completion time. Compared with the conventional IPPS problem, FIPPS considers the fuzzy process time in the uncertain production environment, which is more practical and realistic. However, it is difficult to solve the FIPPS problem due to the complicated fuzzy calculating rules. To solve this problem, this paper formulates a novel fuzzy mathematical model based on the process network graph and proposes a MultiSwarm Collaborative Optimization Algorithm (MSCOA) with an integrated encoding method to improve the optimization. Different swarms evolve in various directions and collaborate in a certain number of iterations. Moreover, the critical path searching method is introduced according to the triangular fuzzy number, allowing for the calculation of rules to enhance the local searching ability of MSCOA. The numerical experiments extended from the well-known Kim benchmark are conducted to test the performance of the proposed MSCOA. Compared with other competitive algorithms, the results obtained by MSCOA show significant advantages, thus proving its effectiveness in solving the FIPPS problem.

Key words: Integrated Process Planning and Scheduling (IPPS); fuzzy processing time; fuzzy completion time; MultiSwarm Collaborative Optimization Algorithm (MSCOA)

# **1** Introduction

Process planning and shop scheduling are two independent subsystems that are carried out sequentially in a conventional manufacturing system<sup>[1–3]</sup>. However, minimal coordination between the two subsystems may result in process and production fragmentation, which in turn, may lead to the generation of bottleneck resources, conflicting optimization objectives, and unbalanced machine loads<sup>[4]</sup>. Integrated Process Planning and Scheduling (IPPS) could achieve better process routes than process planning alone; it can also provide more efficient scheduling solutions than shop scheduling alone<sup>[5]</sup>. Therefore, due to the urgent need to investigate the IPPS problem, this has become a hot research topic in the field of industrial engineering.

Khoshnevis and Chen<sup>[6]</sup> are the first to propose

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the concept of integrating process planning and shop scheduling by demonstrating a methodology and the potential benefits of the integration. Jain et al.<sup>[7]</sup> tried to solve the Job-shop Scheduling Problem (JSP) with multiple process routes in a Flexible Manufacturing System (FMS). Chan et al.<sup>[8]</sup> proposed a genetic algorithm with dominant genes to solve the scheduling problem in FMS with optional process routes. Li et al.<sup>[9]</sup> proposed a hybrid algorithm for solving the Kim dataset, which refreshes most of the results. Zhang and Wong<sup>[10]</sup> proposed an enhanced ant colony optimization algorithm to solve the Kim benchmark and refreshed 17 out of 24 problems, including four new solutions reaching the lower bound. Li et al.<sup>[11]</sup> subsequently developed a GA and variable neighborhood search algorithm and, once again, updated nine solutions of this dataset, of which three reaches the lower bound. Liu et al.<sup>[12]</sup> first built a new MILP model based on the network graph's OR-nodes, and established four submodels with less flexibility through a flexibility decomposition strategy. They then combined this with the model collaboration framework to achieve the rapid search over a large-scale IPPS problem, thus creating one of the most effective methods for solving the IPPS problem.

In the above studies, the production parameters are treated as deterministic values<sup>[13–17]</sup>. However, a variety of uncertain factors may lead to the uncertainty of production parameters<sup>[18–20]</sup>. Specifically, one of the characteristics that cannot be neglected in actual production is the uncertainty of the processing time of the workpiece<sup>[21, 22]</sup>. Therefore, investigating the shop scheduling problem considering fuzzy or uncertain processing time is of great practical significance<sup>[23, 24]</sup>.

Related to the abovementioned problem, Behnamian<sup>[25]</sup> presented an extensive review of the fuzzy shop scheduling problem according to shop types, including single machines, parallel machines, flowshop, job shop, and open shop. Among them, fuzzy job shop scheduling and flow shop scheduling problems were studied the most, while parallel machines and open shop fuzzy scheduling were studied the least. Abdullah and Abdolrazzagh-Nezhad<sup>[26]</sup> conducted a review of the fuzzy JSP, focusing on the meta-heuristics of the solution methods and found that 63% of the approaches for solving fuzzy JSP are GA related. Wang et al.<sup>[27]</sup> proposed a hybrid artificial bee colony algorithm for solving the fuzzy flexible job-shop scheduling problem. Wang et al.<sup>[28]</sup> further designed a hybrid adaptive differential evolution algorithm to solve a multiobjective Tsinghua Science and Technology, April 2024, 29(2): 285–304

#### fuzzy JSP.

Gao et al.<sup>[29]</sup> enhanced the classic differential evolution algorithm through a selection mechanism to solve JSP with fuzzy execution time and fuzzy completion time. Cai and Lei<sup>[30]</sup> proposed a solution to the distributed fuzzy hybrid flow shop scheduling problems with fuzzy processing time by designing a cooperated shuffled frog-leaping algorithm to optimize the objective of fuzzy makespan, total agreement index, and fuzzy total energy consumption simultaneously. For the FJSP in the type-2 fuzzy logic system, Li et al.<sup>[31]</sup> proposed an improved artificial immune system algorithm to solve a special case of FJSP with the processing time, which is the nonsymmetric triangular Interval T2FS (IT2FS) value.

Although there are relatively more studies on fuzzy scheduling problems for flow shop and job shop types<sup>[32]</sup>, very few studies have investigated the FIPPS problem<sup>[33]</sup>. In fact, Wen et al.<sup>[34]</sup> used the Triangular Fuzzy Number (TFN) to represent uncertain processing time and proposed a multiobjective GA to solve the multiobjective IPPS problem by minimizing makespan, maximal machine workload, total machine workload, and total flow time. Zhang et al.<sup>[35]</sup> also used the TFN to represent the processing time and the transportation time in the distributed manufacturing environment, after which they proposed an extended GA with a three-class encoding method, improved crossover, and mutation to optimize the FIPPS problems. Wen et al.<sup>[36]</sup> further improved the encoding based on previous research and proposed a multilayer collaborative optimization method to solve the fuzzy multiobjective IPPS problem while including customer satisfaction as one of the optimizing objectives.

In accordance with the research status on the fuzzy shop scheduling problem presented above, we can conclude that, first, the current mainstream solving approach is still the intelligent optimization algorithm<sup>[37–41]</sup>. Among them, the swarm intelligence algorithm is one of the most common methods<sup>[42, 43]</sup>. Second, fewer studies have been conducted on the FIPPS problem, of which modeling research is urgently required. Third, the encoding methods based on the problem characteristics have not yet been significantly improved. Furthermore, most studies fail to effectively increase the integration of process planning and shop scheduling.

Therefore, the current paper establishes a new mathematical model for the FIPPS problem based on

the process network graph, and designs a MultiSwarm Collaborative Optimization Algorithm (MSCOA) with a novel integrated encoding approach. The integrated encoding approach can improve the integration degree of process planning and shop scheduling, extend the original solution space, and provide the possibility of finding more candidate solutions with better quality. The collaboration and interaction among multiple swarms enable the algorithm to achieve a global and local search balance.

The remainder of this paper is organized as follows: Section 2 presents the problem description and the mathematical model; Section 3 introduces the proposed MSCOA; Section 4 shows the comparative experiments and discussions; and finally, Section 5 discusses the conclusions and future research directions.

# 2 **Problem Formulation**

# 2.1 Problem description

The FIPPS problem can be defined as follows: N jobs with pending process routes are assigned to M machines for processing with fuzzy processing time. The FIPPS aims to plan proper process routes for each job and to assign reasonable machines to meet specific constraints and optimize the objective, such as maximum fuzzy completion time.

The process information of the jobs is usually represented by network graphs, as shown in Fig. 1. The corresponding values 9 of processing time of the operations are in the form of TFNs, as shown in Table 1. The three numbers in the bracket are the minimum, most probable, and maximum values of the processing time, respectively. Every job of the IPPS has its process network graph from which the process routes are planned. The network graph comprises five types of nodes: (1) the start node, which is virtual and represents the start of a job's production process;



Fig. 1 Network graphs of three jobs.

Job No.	Operation	Alternative machine	Processing time
	1	1, 2	(1, 2, 4), (2, 3, 5)
1	2	2, 4	(1, 1, 2), (1, 2, 2)
	3	2, 3	(2, 4, 5), (2, 3, 5)
	1	3, 5	(2, 3, 5), (2, 4, 6)
2	2	3, 4	(1, 4, 5), (1, 2, 4)
2	3	1, 5	(1, 2, 3), (2, 3, 5)
	4	1, 4, 5	(1, 2, 3), (1, 2, 4), (2, 3, 5)
	1	2, 3	(2, 3, 5), (1, 2, 4)
	2	2,4	(1, 3, 5), (1, 2, 4)
2	3	1, 5	(2, 4, 6), (5, 6, 9)
3	4	1, 3	(1, 2, 4), (2, 3, 6)
	5	1, 3, 4	(2, 3, 4), (1, 2, 4), (1, 2, 3)
	6	2,4	(3, 4, 6), (2, 3, 5)

(2) the end node, which is also virtual and indicates the end of the production process; (3) the intermediate nodes, which imply operations; and (4) the OR-nodes, which are combined with (5) the JOIN-node to represent the process flexibilities. The arrows connecting the operation nodes represent the precedence relationships between operations. For example, the arrow connecting operation nodes 3 and 4 in Job 2's network indicates that Operation 3 of Job 2 should be processed before Operation 4. Furthermore, only one of the two links between a pair of OR-node and JOIN-node would be selected, and the operations inside the selected link will be chosen as well. The time unit in the paper is omitted.

The operating rules on the TFN values of fuzzy processing time,  $\tilde{P} = (\tilde{p}_1, \tilde{p}_2, \tilde{p}_3)$  and  $\tilde{P}' = (\tilde{p}'_1, \tilde{p}'_2, \tilde{p}'_3)$ , are introduced below.

# (1) Addition rule:

- $\tilde{P} + \tilde{P}' = (\tilde{p}_1 + \tilde{p}_1', \tilde{p}_2 + \tilde{p}_2', \tilde{p}_3 + \tilde{p}_3').$
- (2) Subtraction rule:  $\tilde{z}$

$$P - P' = (p_1 - p'_1, p_2 - p'_2, p_3 - p'_3).$$

# (3) Comparing rule:

Situation 1:

If 
$$\frac{\tilde{p}_1 + 2\tilde{p}_2 + \tilde{p}_3}{4} > (<)\frac{\tilde{p}_1' + 2\tilde{p}_2' + \tilde{p}_3'}{4}$$
, then  $\tilde{P} > (<)\tilde{P}'$ .

Situation 2:

If 
$$\frac{\tilde{p}_1 + 2\tilde{p}_2 + \tilde{p}_3}{4} = \frac{\tilde{p}'_1 + 2\tilde{p}'_2 + \tilde{p}'_3}{4}$$
, then compare  $\tilde{p}_2$  and  $\tilde{p}'_2$ . If  $\tilde{p}_2 > (<)\tilde{p}'_2$ , then  $\tilde{P} > (<)\tilde{P}'$ .

Situation 3:

If 
$$\tilde{p}_2 = \tilde{p}'_2$$
, then compare  $(\tilde{p}_3 - \tilde{p}_1)$  and  $(\tilde{p}'_3 - \tilde{p}'_1)$ .

 Table 1
 Processing time of the three jobs in Fig. 1.

If  $(\tilde{p}_3 - \tilde{p}_1) > (<)(\tilde{p}'_3 - \tilde{p}'_1)$ , then  $\tilde{P} > (<)\tilde{P}'$ .

### (4) Max rule:

If  $\tilde{P} > \tilde{P}'$ , then  $\tilde{P} \vee \tilde{P}' = \tilde{P}$ ; otherwise,  $\tilde{P} \vee \tilde{P}' = \tilde{P}'$ .

# 2.2 Mathematical model of FIPPS

The FIPPS mathematical model is introduced below<sup>[12, 44, 45]</sup>, and the notations used in this model are shown in Table 2.

Objective:

$$\min C_{\max} \tag{1}$$

Constraints:

$$\sum_{l} R_{irl} = 1, \forall i, r \tag{2}$$

$$X_{ij} \leq M \times (1 - W_{ijrl}) + M \times R_{irl}, \forall i, j, r, l \quad (3)$$

$$X_{ij} \ge 1 - M \times \sum_{r} \sum_{l} W_{ijrl} \times (1 - R_{irl}), \forall i, j \quad (4)$$

$$\sum_{k} Z_{ijk} = X_{ij}, \forall i, j$$
(5)

$$M \times (2 - X_{ij} - X_{ij'}) + Y_{ijij'} \ge U_{ijj'}, \forall i, j, j', j \neq j'$$
(6)

Table 2Definition of notations in the FIPPS model.

Notation	Definition					
N	Set of jobs					
$J_i$	Operation set of Job <i>i</i>					
$R_i$	OR-node set of Job i					
М	Set of machines					
i, i'	Job, $1 \leq i, i' \leq  N $					
j, j'	Operation, $1 \leq j, j' \leq  J_i $					
$O_{ij}$	Operation <i>j</i> of Job <i>i</i>					
k, k'	Machine, $1 \leq k \leq  M $					
r, r'	OR-node, $1 \leq r \leq  R_i $					
l, l'	Link					
$arPsi_i^{ ext{Job}}$	Operation set of Job <i>i</i>					
$\Phi_i^{\mathrm{Last}}$	Possible last operation set of Job <i>i</i>					
	1, $O_{ij}$ is processed before $O_{ij'}$ according to the					
$U_{ijj'}$	precedence relationship represented by the network					
	of Job <i>i</i> ; 0, otherwise					
$\tilde{P}_{ijk}$	Fuzzy processing time of $O_{ij}$ on Machine k					
Week	1, $O_{ij}$ is located in the <i>l</i> -th link of the <i>r</i> -th OR-node;					
<i>vijri</i>	0, otherwise					
Rint	1, if the $l$ -th link of the $r$ -th OR-node of Job $i$ is					
$\mathbf{n}_{l}\mathbf{r}_{l}$	selected; 0, otherwise					
$X_{ij}$	1, if $O_{ij}$ is selected; 0, otherwise					
$Z_{ijk}$	1, if $O_{ij}$ is processed on Machine $k$ ; 0, otherwise					
$Y_{iji'j'}$	1, if $O_{ij}$ is processed before $O_{i'j'}$ ; 0, otherwise					
$\tilde{S}_{ij}$	Fuzzy start time of $O_{ij}$					
$\tilde{C}_{ij}$	Fuzzy completion time of $O_{ij}$					
$ ilde{C}_{\max}$	Maximum fuzzy completion time					

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$$\tilde{S}_{ij'} \ge \tilde{S}_{ij} + \sum_{k} \tilde{P}_{ijk} \times Z_{ijk} - M \times (3 - X_{ij} - X_{ij'} - Y_{ijij'}),$$

$$\forall i, j, j', j \neq j'$$
(7)

$$\tilde{S}_{ij} \geq \sum_{k} \tilde{P}_{ijk} \times Z_{ij'k} - M \times (2 - X_{ij} - X_{ij'} + Y_{ijij'}),$$

$$\forall i, j, j', j \neq j'$$
(8)

$$\tilde{S}_{i'j'} \ge \tilde{S}_{ij} + \tilde{P}_{ijk} - M \times \left(3 - Z_{ijk} - Z_{i'j'k} - Y_{iji'j'}\right), \\
\forall i, i', j, j', i \neq i' \text{ or } j \neq j', k$$
(9)

$$\tilde{S}_{ij} \ge \tilde{S}_{i'j'} + \tilde{P}_{i'j'k} - M \times \left(2 - Z_{ijk} - Z_{i'j'k} + Y_{iji'j'}\right),$$
  
$$\forall i, i', j, j', i \neq i' \text{ or } j \neq j', k$$
(10)

$$C_{\max} \ge S_{ij} + \sum_{k} \tilde{P}_{i'j'k} \times Z_{ijk} - M \times (1 - X_{ij}),$$
  
$$\forall i, j$$
(11)

Constraint in Eq. (2) means that only one link of an OR node can be selected. Constraints in Formulas (3) and (4) are the OR-node's link-selecting conditions, while constraint in Eq. (5) indicates that one operation would only be assigned to one machine for processing. Constraint in Formula (6) corresponds to the precedence relationships in the process network. Meanwhile, constraints in Formulas (7) and (8) denote that the two operations of the same job are supposed to be processed sequentially according to the precedence relationships of the corresponding networks. Similarly, the operations assigned to the same machines should also be arranged according to the precedence constraints in the process networks that are formulating constraints in Formulas (9) and (10). Finally, constraint in Formula (11) is the maximum fuzzy completion time constraint. The TFN calculation rules make the model nonlinear, which means it cannot be solved and verified by a general mathematical solver. Therefore, the above model is only used for the problem description function.

# 3 Multiswarm Collaborative Optimization Algorithm for FIPPS

The proposed MSCOA framework is shown in Fig. 2. The main steps of MSCOA are described below.

**Step 1:** Randomly generate three swarms named swarm\_1, swarm\_2, and swarm\_3, and set the generation number Gen=1.

**Step 2:** Perform crossover and mutation operator on swarm\_2 and update its individuals.

**Step 3:** Perform crossover and mutation operator on swarm\_3 along with the modified N8 searching on it to update its individuals.

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Fig. 2 MSCOA framework.

**Step 4:** Repeat Steps 2 and 3 20 times and set Gen = Gen + 20.

**Step 5:** Evaluate the three swarms and determine whether the termination criteria are satisfied. If "Yes", go to Step 7; if "No", go to Step 6.

**Step 6:** Execute crossover on individuals between swarm\_3 and swarm\_1, and update the individuals of swarm\_3. Execute crossover on individuals between swarm\_2 and swarm\_1 and update the individuals of swarm\_1. Gen = Gen + 1. Go to Step 2.

Step 7: Output the best solution.

Once the three swarms are generated randomly, swarm\_1 does not undergo any evolution during the 20 iterations of Steps 2 and 3 to preserve the original status and to provide a basis for swarm backtracking. Instead, swarm\_2 and swarm\_3 are evolved differentially: crossover and mutation are executed for swarm\_2, and a modified critical path-based N8 searching is additionally performed for swarm\_3<sup>[46]</sup>, which can enhance the local search capability of the MSCOA. Next, inter-swarm interactions are executed once sufficient differentiation between the various swarms has been generated (20 generations of evolution in each). These interactions are as follows: individuals from swarm\_1 and swarm\_3 are selected separately to perform crossover operators to update the two swarms; then, the same is done to individuals from swarm\_1 and swarm\_2 to update the

two swarms.

The purpose of differential evolution is to increase the diversity of individuals and improve the overall global search ability of the algorithm. At present, the N8 neighborhood searching in swarm\_3 is the most effective critical path structure for solving JSP<sup>[46]</sup>, as it can effectively enhance the local searching ability. Finally, we can achieve a balance between the global and local searching of the MSCOA after following the above two mechanisms.

#### 3.1 Encoding and decoding method

The paper adopts the integrated encoding approach proposed by Liu et al.<sup>[47]</sup> There are three coding strings for the corresponding three subproblems of the FIPPS problem. The first string is the OR-node string, which represents the link-selecting state of each OR-node. This is arranged according to the job number order. The inside numbers "1" and "2" represent the options of selecting either the left or the right link. As for the operation string, it is arranged according to the processing sequence. A pair of numbers "i - j" in the block represents the operation  $O_{ii}$ , as shown in Fig. 3. The third string is the machine string, which represents the processing machine assignment plan. This string is also arranged according to the job numbers from left to right. For example, the number "2" in the second block means that operation 2 of Job 1 selects Machine 2 for processing. We adopt the semiactive schedule decoding method to obtain the production scheme. Here, the operations are arranged one by one on the corresponding machines according to the processing sequence and the machine assignment plans.

#### 3.2 Initialization and selection

In accordance with the abovementioned encoding rules, we adopt the random initialization method to initialize different strings of individuals. Thus, the ORnode strings are initialized by randomly selecting the corresponding links. Machine strings are initialized



Fig. 3 Example of an encoding individual.

by randomly selecting the processing machine for each operation. As for the operation string, this is initialized by randomly sequencing the operations under the constraint of precedence relationships in the networks.

Next, we adopt the tournament selection method to select individuals that can carry evolutionary operators. Two individuals in the swarm are chosen randomly, and the one with better fitness will be selected if the randomly generated value (between 0 and 1) is greater than the predefined reproduction probability value  $P_r$ . Otherwise, the individual with the worse fitness will be selected. In the paper, the reproduction probability  $P_r$  is set to 0.8.

#### 3.3 Crossover and mutation

The crossover for the OR-node string is shown in Fig. 4. One crossover point a is randomly selected. Parents 1 and 2 pass the left part of the blocks to the corresponding Offspring 1 and 2, after which they pass the right half of the blocks to the other one. The mutation operator for the OR-node string is a single-point shifting action, in which a mutation point is randomly selected, and the state of that point is shifted, as shown in Fig. 5.

The crossover operator for the operation string is shown in Fig. 6. Two crossover points, a and b, are randomly selected. The blocks outside the two points of Parent 1 are transferred to offspring individual 1, and the blocks outside the two points of Parent 2 are transferred to offspring individual 2. Then, Parent 2 passes the different blocks to Offspring 1, and Parent 1 passes the different blocks to Offspring 2. There are three steps











Fig. 6 Crossover for the operation string.

involved in conducting mutation on an operation string, as shown in Fig. 7. First, a mutation point is randomly chosen, after which the mutation range of the chosen point is determined according to the corresponding network. Finally, one of the available positions inside the mutation range is randomly selected, and the chosen operation is inserted.

Given that the position orders of their blocks in the machine string contain the corresponding job and operation information, the crossover operator transfers the blocks at the same positions from the parent to offspring individuals, as shown in Fig. 8. The mutation operator of the machine string is similar to the action of the OR-node string. Then, one mutation point is randomly selected, and the current state is shifted to another one, as shown in Fig. 9.

#### 3.4 TFN-based critical path searching

The critical path-related approach was first applied in a study on JSP in Ref. [48]. Previous studies have proven



Fig. 8 Crossover for the machine string.



that movements on noncritical paths do not change the makespan. In common shop scheduling problems, the critical path is defined as the longest path from the first operation of the first job to the last operation of the final job<sup>[49]</sup>. The maximum and comparing calculations are achieved according to the normal number rather than the TFN. Therefore, the calculation method of the critical path in the TFN way must be introduced.

A scheduling scheme of the FIPPS can be represented in the form of a disjunctive graph, as shown in Fig. 10. One disjunctive graph consists of a set of nodes and a set of arcs. The nodes correspond to the operations, while the arcs connect pairs of nodes, each representing one precedence relationship. The two nodes, Ostart and  $O_{\rm end}$ , represent the start and the end of the processing, respectively. The two consecutive operation nodes of the same job are connected by the conjunctive arcs indicated by the solid arrows, while the operation nodes assigned successively on the same machines are connected by the disjunctive arcs indicated by the dashed arrows<sup>[50]</sup>. The length of the arc between  $O_{ij}$  and  $O_{i_0j_0}$  is denoted as  $D(O_{ii}, O_{i_0,i_0})$ , and the maximum fuzzy completion time of the scheme is equal to  $D(O_{\text{start}}, O_{\text{end}})$ . In explaining the critical path calculation, we introduce the head length  $\tilde{H}_{ij}$  and the tail length  $\tilde{T}_{ij}$  of  $O_{ij}$  in the form of TFN. The critical path is the longest path from node  $O_{\text{start}}$  to node  $O_{\text{end}}$ , and the nodes within this range are considered critical operations<sup>[50]</sup>. According to this definition, the critical operation in fuzzy scheduling has the following attribution:  $\tilde{H}_{ij} + \tilde{P}_{ij} + \tilde{T}_{ij} =$  $L(O_{\text{start}}, O_{\text{end}})$ , where  $P_{ij}$  is the fuzzy processing time



Fig. 10 Disjunctive graph of an FIPPS example.

of  $O_{ij}$  on its assigned machine, and the head and tail lengths can be obtained,

$$\tilde{H}_{\text{start}} = \tilde{T}_{\text{end}} = (0, 0.0) \tag{12}$$

$$H_{ij} = \max \{ H_{JP_{ij}} + P_{JP_{ij}}, R_{MP_{ij}} + P_{MP_{ij}} \}$$
(13)

$$T_{ij} = \max\{T_{JS_{ij}} + P_{JS_{ij}}, T_{MS_{ij}} + P_{MS_{ij}}\}$$
(14)

In Eqs. (13) and (14),  $JP_{ij}$  is the predecessor operation of  $O_{ij}$  of the same job, and  $JS_{ij}$  is the successor operation of  $O_{ij}$  of the same job. In addition,  $MP_{ij}$ is the predecessor operation of  $O_{ij}$  on the same machine, and  $MS_{ij}$  is the successor operation of  $O_{ij}$  on the same machine. The critical path of one FIPPS scheme can be determined based on the abovementioned definition. Similarly, the effective movements on the operation nodes can also be applied to this fuzzy critical path. The most advanced neighborhood structure is N8, proposed by Xie et al.<sup>[46]</sup>, which mainly has three types of movement: (1) moving operation in the critical block out to the noncritical path parts, (2) inserting the first or the last operation of the critical block into the inner segment, and (3) moving operation inside the critical block to the head or the tail of the block (see Fig. 11).

To increase the disturbance extent on the critical path, this paper modifies the three movements of the N8 neighborhood by (1) swapping the critical operation in the critical block with the operation on the noncritical path and by (2) swapping the head or tail operation with the inner operation of the critical path, as shown in Fig. 12.

# 4 Experimental Study and Discussion

In testing the effectiveness of the MSCOA, we performed comparative experiments with the classical GA and Particle Swarm Optimization (PSO) algorithm. The testing instances are extended from the most wellknown Kim benchmark, which has been widely used in IPPS-related research since its introduction in 2003<sup>[51]</sup>. The benchmark contains 24 instances comprising 18 jobs and 15 machines. Each instance consists of a different combination of jobs. For follow-up researchers using a new benchmark, the fuzzy processing time data of the 18 jobs are presented in Tables A1-A6 in Appendix A. The minimum and maximum values are obtained from the processing time in Kim's original data with a random deviation of  $\pm (10-40)\%$ . The original processing time is treated as the most probable value of the fuzzy processing time, and the iteration number is set to 500. The parameters are the Swarm\_Size = 200, the reproduction probability  $P_r = 0.8$ , the crossover



Moving operation in the critical block out to the non-critical path parts
 Inserting the first or last operation of the block into the inner segment
 Moving operation inside the critical block to the head or tail of the block

#### Fig. 11 N8 neighborhood structure schematic.



(1) Swapping the critical operation with the operation on the non-critical path(2) Swapping the head or tail operation with the inner operation of the critical path

Fig. 12 New neighborhood structure adopted in MSCOA.

probability  $P_c = 0.8$ , and the mutation probability  $P_m=0.2$ . The experiment platform is a PC with an i7-8700 CPU and 16 GB RAM. The Gantt chart of the best solution (391, 522, 663) of Kim 24 is shown in Fig. 13. The results are obtained by running the algorithms independently 20 times in Tables 3–6.

Table 3 shows the best results that can be obtained upon using the three algorithms. The proposed MSCOA can find better results on all 24 problems than the best results of GA and PSO. From a statistical perspective, we analyze the results of 20 independent runs. As shown by the results, the proposed MSCOA has a greater advantage in the minimum, most probable, and maximum values of  $\tilde{C}_{max}$  with better average and standard deviation values, thus indicating the superiority and stability of the results. In Table 5, all 24 results of the most probable value are found to be superior to those of GA and PSO, thereby proving that the MOSCA outperforms the other two algorithms.

Observing the results in Tables 3-6, we can see



Fig. 13 Gantt chart of the best solution pertaining to Kim\_24 (time unit is omitted according to the original data).

Number	MSCOA		GA		PSO				
Number	Minimum	Most probable	Maximum	Minimum	Most probable	Maximam	Minimum	Most probable	Maximum
1	306	428	529	319	429	533	313	428	526
2	250	343	411	257	348	425	255	345	414
3	249	346	430	270	352	428	265	349	426
4	226	306	373	234	306	368	217	305	388
5	229	314	385	226	318	378	227	314	397
6	313	427	527	331	451	559	323	437	545
7	262	372	468	266	372	467	265	375	460
8	250	339	423	264	339	429	257	345	415
9	315	428	520	318	430	527	313	427	527
10	317	424	534	358	464	577	340	437	541
11	273	350	424	285	371	454	275	354	424
12	234	322	388	247	337	408	236	336	404
13	317	431	530	355	471	575	351	467	585
14	273	378	474	298	403	510	295	392	473
15	313	427	527	341	445	555	322	436	537
16	328	432	537	363	494	613	363	483	576
17	291	384	449	348	466	569	316	432	538
18	246	349	419	284	390	496	283	386	484
19	331	445	546	399	505	624	376	490	601
20	301	394	501	342	464	570	320	448	572
21	317	436	539	369	482	602	363	462	583
22	358	481	616	423	568	702	388	543	674
23	326	440	558	378	528	635	348	490	626
24	391	522	663	445	619	776	445	590	726

 Table 3
 Best results (maximum fuzzy completion time) of 20 independent runs of the three algorithms.

that, on the one hand, the PSO can obtain better fuzzy values than the conventional GA algorithm, mainly due to the adoption of a more advanced evolutionary framework. The proposed MSCOA, on the other hand, can achieve better results than the two compared algorithms, benefiting from the enhanced swarm collaborative framework and the critical pathbased neighborhood search. The experimental results also demonstrate the effectiveness of the improvements at the level of the algorithmic framework and the operators.

# 5 Conclusions and Future Research Directions

The current study focuses on the FIPPS problem considering fuzzy processing time and builds a process network based mathematical model. Then, we propose the MOSCA based on an integrated encoding method with critical path neighborhood searching. In this work, we sufficiently consider the advantages of the critical path-based N8 neighborhood searching and apply it to the FIPPS problem by modifying the calculation procedure according to the TFN rules. The experimental results on the extended famous benchmark indicate the outstanding performance of the proposed MSCOA in terms of its searching capability and stability.

Furthermore, by guiding the searching and evolving of each swarm in different directions, we ensure the diversity of all individuals and prevent the premature maturity of the algorithm. The proposed integrated encoding can also effectively promote the integration of the two systems of process and production, thus providing more candidate solutions with better quality. Combined with the modified critical path N8 searching, the proposed approach effectively reduces the invalid searching actions, enabling the proposed MSCOA to demonstrate significant superiority on famous open problems.

In the comparison results on the minimum values of  $\tilde{C}_{max}$ , MSCOA had three small-scale instances that did not perform best. Therefore, exploring more adaptive swarm collaboration strategies in future works might improve the above situation. Furthermore, developing new critical path structures according to the TFN calculation rules is another important research direction

	MS	SCOA	(	GA	P	SO	Signit	ficant?
Number	Average	Standard deviation	Average	Standard deviation	Average	Standard deviation	MSCOA better than GA	MSCOA better than PSO
1	314.30	5.17	330.70	10.08	327.95	9.36	Yes	Yes
2	253.45	4.41	264.40	7.38	262.60	8.42	Yes	Yes
3	263.25	6.96	278.70	12.24	277.90	9.69	Yes	Yes
4	231.75	5.64	235.40	5.61	234.85	7.11	No	No
5	230.85	5.23	238.10	7.22	235.95	7.19	Yes	No
6	326.70	8.35	352.65	13.76	337.45	8.70	Yes	Yes
7	269.50	5.69	274.55	6.35	271.50	6.80	No	No
8	253.60	4.08	265.00	8.01	265.15	11.22	Yes	Yes
9	320.05	7.69	330.10	11.62	324.10	7.35	Yes	No
10	325.80	7.04	364.20	14.45	348.50	12.44	Yes	Yes
11	275.60	9.52	304.15	7.00	296.60	14.35	Yes	Yes
12	239.40	8.69	272.60	13.73	262.15	13.22	Yes	Yes
13	335.90	10.94	376.35	16.05	363.65	12.66	Yes	Yes
14	280.90	9.69	316.80	14.15	306.90	14.48	Yes	Yes
15	324.50	9.43	354.05	8.86	344.40	16.07	Yes	Yes
16	343.85	15.38	395.20	21.52	374.80	15.02	Yes	Yes
17	303.60	12.92	363.05	17.56	352.70	15.31	Yes	Yes
18	264.35	12.29	322.50	20.37	300.60	8.15	Yes	Yes
19	354.90	15.18	413.30	19.21	395.25	22.42	Yes	Yes
20	311.30	15.66	367.30	17.09	344.80	17.58	Yes	Yes
21	342.65	9.65	395.40	15.76	374.35	15.34	Yes	Yes
22	382.30	13.00	444.60	16.80	428.30	19.11	Yes	Yes
23	342.65	11.63	421.05	17.19	395.30	20.29	Yes	Yes
24	411.15	11.94	494.05	25.55	467.85	22.22	Ves	Yes

Table 4 Minimum values in  $\tilde{C}_{max}$  by T-test ( $p \leq 0.05$ ).

that should be explored.

#### Appendix A

- Table A1 Fuzzy processing time of Jobs 1–3.
- Table A2Fuzzy processing time of Jobs 4–6.
- Table A3Fuzzy processing time of Jobs 7–9.
- Table A4Fuzzy processing time of Jobs 10–12.
- Table A5Fuzzy processing time of Jobs 13–15.
- Table A6Fuzzy processing time of jobs 16–18.

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	MS	SCOA	(	GA	Р	SO	Signif	icant?
Number	Average	Standard	Average	Standard	Average	Standard	MSCOA better	MSCOA better
		deviation		deviation		deviation	than GA	than PSO
1	427.20	0.52	443.75	0.73	438.50	0.70	Yes	Yes
2	344.80	0.48	358.50	0.62	357.90	0.67	Yes	Yes
3	346.60	0.61	366.45	0.80	358.95	0.71	Yes	Yes
4	307.25	0.54	315.05	0.54	314.90	0.61	Yes	Yes
5	321.55	0.52	328.55	0.62	327.25	0.61	Yes	Yes
6	433.80	0.66	467.45	0.85	447.60	0.68	Yes	Yes
7	371.35	0.55	379.75	0.58	377.30	0.60	Yes	Yes
8	345.00	0.46	355.65	0.65	359.10	0.77	Yes	Yes
9	424.40	0.64	441.30	0.78	435.85	0.62	Yes	Yes
10	438.85	0.61	482.75	0.87	463.65	0.81	Yes	Yes
11	359.75	0.71	401.45	0.61	386.95	0.87	Yes	Yes
12	331.05	0.68	370.10	0.85	358.75	0.83	Yes	Yes
13	447.90	0.76	504.80	0.92	484.25	0.82	Yes	Yes
14	385.85	0.71	435.35	0.86	418.60	0.87	Yes	Yes
15	433.45	0.70	466.40	0.68	454.20	0.92	Yes	Yes
16	456.65	0.90	529.10	1.06	500.10	0.89	Yes	Yes
17	409.35	0.82	493.15	0.96	465.30	0.90	Yes	Yes
18	367.25	0.80	439.05	1.04	407.55	0.65	Yes	Yes
19	475.95	0.89	550.30	1.01	530.50	1.09	Yes	Yes
20	426.20	0.91	499.90	0.95	472.10	0.96	Yes	Yes
21	453.80	0.71	520.45	0.91	489.75	0.90	Yes	Yes
22	505.45	0.83	593.95	0.94	571.35	1.00	Yes	Yes
23	468.85	0.78	572.25	0.95	531.10	1.03	Yes	Yes
24	550.40	0.79	663.95	1.16	624.25	1.08	Yes	Yes

Table 5 Most probable values in  $\tilde{C}_{max}$  by T-test ( $p \leq 0.05$ ).

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	MS	SCOA	(	GA	P	SO	Signif	icant?
Number	Average	Standard deviation	Average	Standard deviation	Average	Standard deviation	MSCOA better than GA	MSCOA better than PSO
1	530.15	6.24	551.35	14.00	546.15	14.21	Yes	Yes
2	420.40	8.65	435.80	14.37	435.95	17.47	Yes	Yes
3	428.65	7.80	448.15	15.30	441.95	12.09	Yes	Yes
4	372.70	7.14	383.40	10.32	386.90	11.90	Yes	Yes
5	392.35	8.27	404.70	10.25	401.60	9.07	Yes	Yes
6	538.80	9.45	582.25	18.77	558.30	13.27	Yes	Yes
7	466.50	6.52	478.35	9.39	472.60	11.54	Yes	No
8	420.95	8.29	434.75	12.35	440.55	24.74	Yes	Yes
9	527.55	6.07	546.95	13.36	543.10	12.11	Yes	Yes
10	545.70	13.52	608.35	21.48	577.35	21.37	Yes	Yes
11	446.00	10.46	496.55	15.02	482.15	19.63	Yes	Yes
12	404.25	14.27	455.75	21.10	441.45	22.01	Yes	Yes
13	557.45	16.16	628.65	24.53	603.25	13.75	Yes	Yes
14	485.20	10.10	548.10	18.71	519.75	26.34	Yes	Yes
15	539.40	11.49	582.05	14.90	563.40	19.54	Yes	Yes
16	571.20	21.90	663.70	34.52	625.80	22.86	Yes	Yes
17	503.25	19.66	605.60	25.95	575.80	28.72	Yes	Yes
18	456.10	16.32	543.20	22.90	506.60	19.27	Yes	Yes
19	595.55	18.76	685.35	25.83	666.45	33.07	Yes	Yes
20	530.35	22.02	623.70	21.69	591.25	23.09	Yes	Yes
21	565.50	15.78	649.90	26.27	610.20	20.75	Yes	Yes
22	634.75	14.89	743.70	22.79	711.25	25.03	Yes	Yes
23	582.60	20.87	710.45	33.51	666.75	24.58	Yes	Yes
24	684.60	18.59	826.40	28.09	773.05	32.74	Yes	Yes

Table 6 Maximum values in  $\tilde{C}_{max}$  by T-test ( $p \leq 0.05$ ).

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		Table A1 Fuzzy	processing time of Jobs 1–5.
Job No.	Operation	Alternative machine	Processing time
	1	9, 14	(12, 13, 16), (8, 10, 13)
	2	11, 15	(20, 24, 32), (15, 18, 25)
	3	15	(27, 43, 59)
1	4	12	(29, 43, 47)
1	5	13	(22, 30, 38)
	6	4, 12	(28, 32, 38), (16, 25, 31)
	7	1, 5, 11	(26, 40, 48), (32, 49, 68), (34, 39, 45)
	8	8	(35, 47, 57)
	1	5, 8, 14	(8, 10, 14), (13, 16, 22), (11, 13, 18)
	2	8, 9, 15	(5, 6, 7), (6, 8, 10), (4, 7, 9)
	3	4	(29, 40, 49)
	4	6, 7, 9, 12	(12, 14, 17), (12, 20, 28), (8, 10, 12), (10, 13, 15)
	5	1, 7, 11	(20, 33, 41), (31, 40, 48), (29, 43, 49)
	6	1, 5	(30, 42, 48), (25, 38, 42)
2	7	6, 11, 15	(19, 25, 31), (23, 33, 42), (22, 30, 35)
2	8	10, 15	(36, 41, 52), (37, 44, 53)
	9	2, 13	(9, 10, 12), (8, 12, 14)
	10	11, 14, 15	(25, 34, 38), (15, 24, 30), (22, 30, 38)
	11	6, 11	(26, 38, 52), (31, 42, 56)
	12	4, 8, 12	(20, 25, 33), (17, 26, 34), (18, 30, 35)
	13	7	(30, 39, 50)
	14	10, 12	(23, 37, 50), (27, 40, 51)
	1	4, 7, 11	(21, 29, 37), (31, 36, 50), (27, 34, 47)
	2	2, 8, 11, 12, 15	(28, 35, 43), (21, 29, 37), (17, 27, 37), (26, 30, 35), (26, 33, 40)
	3	1, 6, 7, 10, 14	(9, 11, 14), (7, 9, 11), (7, 8, 9), (15, 19, 22), (9, 12, 15)
	4	2, 3, 6, 7	(14, 18, 23), (13, 20, 23), (20, 27, 36), (11, 13, 16)
	5	2, 3, 4, 5, 7	(13, 19, 22), (16, 24, 30), (17, 22, 26), (23, 31, 41), (27, 37, 41)
	6	1, 7, 10, 12, 13	(9, 13, 16), (5, 8, 11), (10, 12, 16), (5, 9, 12), (4, 5, 6)
	7	1, 5, 6, 13	(36, 50, 56), (28, 39, 46), (30, 44, 61), (36, 48, 54)
	8	1, 12	(5, 6, 7), (7, 9, 12)
	9	3, 4, 5, 13, 14	(29, 44, 57), (32, 36, 50), (19, 30, 34), (26, 39, 45), (28, 33, 41)
3	10	3, 4, 7, 12, 13	(28, 39, 50), (29, 45, 59), (33, 41, 56), (42, 50, 67), (34, 40, 54)
	11	2, 4, 9	(23, 29, 40), (22, 36, 40), (25, 33, 38)
	12	1, 4, 9, 12, 14	(16, 19, 22), (13, 20, 25), (11, 17, 21), (11, 16, 22), (14, 21, 29)
	13	9, 14, 15	(33, 40, 45), (23, 33, 45), (25, 35, 48)
	14	2, 4, 6, 7	(9, 11, 14), (9, 12, 16), (12, 14, 18), (11, 15, 19)
	15	1, 8, 12, 13, 14	(7, 10, 13), (16, 19, 26), (16, 20, 23), (14, 17, 21), (14, 16, 21)
	16	10, 15	(44, 49, 65), (34, 44, 51)
	17	10, 12, 13	(18, 20, 23), (28, 33, 40), (28, 39, 44)
	18	6, 7, 9, 10, 13	(24, 30, 38), (21, 29, 33), (31, 40, 46), (24, 39, 48), (26, 33, 45)
	19	3, 8, 10, 14, 15	(14, 20, 22), (23, 29, 35), (27, 40, 50), (27, 34, 47), (28, 31, 40)

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 Table A1
 Fuzzy processing time of Jobs 1–3.

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Job No.	Operation	Alternative machine	Processing time
	1	2	(13, 18, 23)
	2	8, 12	(33, 38, 44), (27, 30, 35)
	3	9	(16, 20, 26)
	4	5, 11	(6, 8, 9), (7, 9, 13)
	5	3	(21, 29, 40)
	6	2, 4, 5	(23, 36, 44), (27, 33, 41), (33, 39, 50)
	7	1.10	(15, 23, 31), (17, 20, 25)
	8	6	(32, 45, 50)
4	9	3.12	(4, 5, 7), (5, 9, 12)
	10	5.9	$(30 \ 39 \ 47) \ (27 \ 33 \ 39)$
	11	7 11	(24, 36, 40) $(30, 41, 57)$
	12	14 15	(19, 31, 35), (18, 29, 39)
	12	1 2 6	(19, 28, 33) $(18, 22, 27)$ $(16, 21, 24)$
	13	1, 2, 0	(19, 20, 33), (10, 22, 27), (10, 21, 24) (16, 18, 24), (20, 28, 30)
	14	10, 14	(10, 10, 24), (20, 20, 59)
	15	12	(20, 24, 27) (18, 23, 27) $(17, 25, 20)$
	10	3 4 5 11	(10, 25, 27), (17, 25, 29) $(20, 35, 48), (21, 29, 33), (23, 36, 45), (27, 31, 38)$
	1	1, 3, 5, 10, 15	(25, 55, 46), (21, 25, 55), (25, 56, 45), (21, 51, 56) (35, 40, 45), (30, 34, 30), (36, 44, 51), (36, 41, 48), (34, 30, 45)
	2	1, 5, 5, 10, 15	(33, 40, 43), (30, 34, 39), (30, 44, 31), (30, 41, 40), (34, 39, 43) (12, 15, 17), (0, 12, 15)
	3	2671014	(12, 13, 17), (2, 13, 13) (27, 21, 42), (26, 23, 43), (21, 20, 40), (22, 27, 37), (10, 25, 34)
	4	6 12 13 14	(27, 51, 42), (20, 53, 43), (21, 27, 40), (22, 27, 57), (17, 23, 54) (0, 12, 16), (6, 0, 12), (6, 9, 11), (11, 14, 17)
	5	0, 12, 13, 14 11 12	(9, 15, 10), (0, 9, 12), (0, 0, 11), (11, 14, 17) (24, 28, 34), (21, 20, 37)
	0	11, 12	(24, 20, 34), (21, 27, 37) (24, 21, 27), (20, 24, 27), (25, 28, 24), (21, 26, 26), (28, 22, 28)
	7 0	6 10 11 12	(24, 51, 57), (20, 24, 27), (25, 26, 54), (21, 20, 30), (20, 52, 56) (27, 24, 28), (25, 22, 40), (22, 20, 28), (24, 20, 40)
	0	0, 10, 11, 15	(27, 54, 56), (23, 55, 40), (25, 50, 56), (24, 29, 40)
5	9	<i>3</i> , 4, 1 <i>3</i>	(34, 41, 33), (24, 37, 43), (25, 40, 43) (22, 28, 51), (21, 20, 28), (21, 25, 45), (26, 20, 25), (22, 21, 20)
	10	1, 4, 5, 6, 9	(35, 56, 51), (21, 29, 56), (21, 55, 45), (20, 50, 55), (25, 51, 59)
	11	6, 10, 11, 14, 13	(41, 46, 03), (41, 30, 06), (52, 44, 39), (29, 41, 37), (53, 47, 53) (18, 26, 22), (26, 22, 20), (28, 28, 42), (10, 20, 26), (22, 20, 20)
	12	0, 7, 9, 12, 13	(16, 20, 55), (20, 52, 59), (26, 56, 42), (19, 29, 50), (25, 50, 59)
	15	6, 10, 13, 14, 15	(18, 25, 25), (15, 20, 25), (19, 25, 54), (12, 18, 24), (17, 22, 25)
	14	2, 8, 10, 12	(11, 14, 17), (9, 11, 14), (14, 17, 20), (10, 15, 17)
	15	3, 5, 15	(20, 27, 38), (17, 24, 31), (17, 20, 30)
	10	4, 8, 15	(14, 20, 23), (15, 19, 23), (12, 14, 10)
	17	2, 3, 4, 7, 14	(17, 27, 34), (14, 21, 28), (24, 28, 37), (20, 30, 36), (23, 29, 39)
	18	3, 7, 10, 14	(34, 39, 53), (29, 34, 40), (27, 40, 50), (20, 35, 41)
	1	3, 4, 11	(31, 38, 49), (27, 33, 40), (32, 36, 44)
	2	2, 3, 6	(14, 22, 26), (14, 21, 26), (12, 19, 23)
	3	5	(11, 14, 16)
	4	5, 8	(15, 17, 22), (15, 20, 22)
	5	3, 6, 11	(22, 36, 40), (27, 33, 38), (31, 39, 54)
	6	2, 3, 10	(16, 24, 31), (16, 20, 25), (12, 18, 21)
	7	3, 9, 15	(16, 21, 23), (15, 17, 19), (16, 24, 32)
	8	8	(33, 38, 44)
	9	4,11	(15, 19, 22), (9, 15, 20)
6	10	8, 10, 15	(11, 14, 18), (15, 19, 23), (11, 17, 23)
	11	1, 4, 9, 13	(22, 25, 30), (14, 21, 27), (15, 19, 23), (21, 28, 36)
	12	7,12	(29, 42, 52), (30, 43, 58)
	13	1, 2, 6	(41, 48, 67), (27, 42, 58), (29, 46, 54)
	14	7, 10	(7, 10, 13), (11, 14, 18)
	15	2, 13, 14	(10, 14, 19), (14, 16, 20), (11, 13, 15)
	16	1, 4, 5, 8	(22, 36, 48), (24, 33, 44), (20, 31, 37), (26, 34, 45)
	17	1, 12	(30, 47, 56), (27, 44, 48)
	18	9, 10, 13	(25, 30, 36), (21, 26, 34), (18, 29, 33)
	19	1, 5, 12	(11, 18, 24), (14, 19, 25), (11, 15, 17)
	20	9, 11	(17, 24, 30), (22, 25, 28)

Table A2Fuzzy processing time of Jobs 4–6.

Job No.	Operation	Alternative machine	Processing time
	1	7, 8	(9, 12, 15), (12, 17, 20)
	2	2, 5, 8, 12	(5, 7, 9), (5, 6, 8), (6, 8, 11), (8, 11, 13)
	3	11, 15	(21, 30, 34), (24, 27, 30)
	4	7	(24, 27, 35)
	5	2, 3, 9	(7, 10, 12), (8, 11, 14), (13, 16, 19)
	6	4, 11, 13	(32, 46, 58), (45, 50, 67), (44, 49, 67)
	7	3	(18, 22, 28)
	8	5, 10	(7, 9, 12), (7, 10, 11)
	9	1, 2, 6	(20, 27, 37), (17, 28, 36), (19, 24, 31)
	10	6, 12, 13	(7, 8, 10), (4, 5, 7), (8, 11, 14)
7	11	2,9	(31, 47, 65), (40, 48, 56)
	12	1, 8	(22, 27, 37), (24, 30, 40)
	13	9, 13, 14	(12, 18, 24), (15, 19, 22), (16, 20, 28)
	14	7, 11	(14, 22, 29), (15, 20, 25)
	15	7, 8, 13	(10, 13, 16), (8, 11, 13), (9, 14, 19)
	16	6, 14	(12, 15, 19), (7, 10, 12)
	17	4, 5, 7	(15, 21, 26), (20, 26, 35), (17, 20, 25)
	18	4, 5, 13	(19, 29, 34), (26, 30, 35), (19, 26, 30)
	19	1, 2	(23, 35, 46), (24, 31, 41)
	20	5, 6, 11	(15, 22, 25), (15, 18, 22), (14, 23, 32)
	21	1, 7, 10	(22, 32, 40), (25, 33, 43), (24, 28, 39)
	1	4	(33, 50, 57)
	2	8, 15	(15, 23, 28), (17, 21, 24)
	3	12	(22, 35, 42)
	4	4, 7, 8	(7, 11, 13), (11, 13, 16), (13, 16, 22)
	5	2, 11	(14, 18, 22), (15, 20, 25)
	6	4, 13	(32, 36, 49), (20, 33, 44)
	7	1, 3	(24, 38, 49), (25, 35, 44)
	8	6, 15	(12, 16, 20), (15, 17, 21)
	9	9	(20, 24, 32)
8	10	5, 14	(17, 23, 27), (20, 26, 29)
	11	1, 5, 8	(11, 15, 19), (12, 16, 18), (10, 17, 20)
	12	3, 10	(38, 43, 47), (32, 49, 58)
	13	10	(29, 44, 59)
	14	15, 14	(20, 32, 44), (27, 31, 41)
	15	10, 15	(28, 30, 43), (27, 38, 42)
	10	14	(21, 20, 50) (26, 20, 54) $(24, 24, 20)$
	17	2, 10	(20, 39, 34), (24, 34, 39) (12, 18, 22), (11, 15, 18)
	10	<i>J</i> , <i>1</i>	(12, 16, 22), (11, 15, 16)
	19	3 12	(14, 10, 20) (33, 45, 53), (36, 48, 53)
	1	5 0 12 13 14	(33, 43, 53), (30, 40, 55)
	2	<i>J</i> , <i>J</i> , 12, 1 <i>J</i> , 14 7	(20, 51, 55), (27, 27, 57), (17, 21, 27), (25, 20, 57), (17, 25, 20) $(18, 21, 27)$
	2	7 9 10 13	(10, 21, 27) (10, 21, 26) $(16, 22, 27)$ $(17, 25, 33)$ $(17, 20, 24)$
9	5 4	3 15	(17, 21, 20), (10, 22, 27), (17, 25, 55), (17, 20, 24) (12, 13, 16), (12, 15, 18)
,	5	3 6 9 12 13	$(4 \ 6 \ 7)$ $(4 \ 5 \ 6)$ $(6 \ 7 \ 8)$ $(7 \ 10 \ 12)$ $(7 \ 9 \ 12)$
	6	1, 2, 5, 6, 10	(3, 3, 7, 50), $(20, 33, 43)$ , $(32, 39, 53)$ , $(22, 29, 38)$ , $(20, 32, 36)$
	7	4, 10, 11, 13, 15	(22, 57, 50, (20, 52, 50), (52, 55), (52, 55), (22, 25, 50), (20, 52, 50) (6, 7, 8), (6, 8, 9), (6, 9, 12), (4, 5, 7), (4, 6, 7)
	1	т, 10, 11, 13, 13	(0, 7, 0), (0, 0, 2), (0, 2, 12), (7, 3, 7), (7, 0, 7)

Table A3Fuzzy processing time of Jobs 7–9.

(To be continued)

		Table A5 Fuzzy pr	(Continued)
Job No.	Operation	Alternative machine	Processing time
	8	1, 4, 8, 9, 12	(35, 42, 49), (34, 41, 49), (35, 39, 47), (31, 45, 61), (34, 44, 57)
	9	4, 7, 14	(8, 10, 14), (8, 9, 12), (10, 14, 19)
	10	4, 5, 7, 8, 14	(13, 19, 22), (12, 14, 17), (11, 15, 18), (11, 12, 16), (8, 10, 14)
	11	3, 9, 10, 11, 15	(23, 28, 39), (15, 24, 32), (18, 27, 32), (14, 22, 31), (17, 20, 25)
	12	4, 6	(34, 45, 58), (29, 41, 51)
	13	2	(40, 44, 51)
9	14	4, 5, 14, 15	(34, 47, 52), (33, 43, 59), (29, 44, 57), (33, 42, 48)
	15	2, 5, 6, 9, 14	(12, 17, 19), (11, 14, 15), (15, 20, 25), (18, 21, 26), (14, 18, 24)
	16	2, 4, 13	(39, 45, 55), (28, 42, 48), (34, 46, 58)
	17	1, 3, 5, 6, 15	(20, 27, 30), (20, 25, 28), (17, 23, 30), (24, 28, 35), (17, 20, 24)
	18	2, 7, 9, 10, 11	(8, 10, 11), (8, 12, 16), (14, 17, 21), (13, 16, 21), (7, 11, 13)
	19	2, 13	(6, 9, 12), (8, 10, 13)
	20	4, 6, 8, 12	(12, 18, 21), (15, 23, 31), (18, 21, 28), (22, 25, 30)

Table A3Fuzzy processing time of Jobs 7–9.

Table A4Fuzzy processing time of Jobs 10–12.

Jobs No.	Operations	Alternative machines	Processing time
	1	1, 2, 3, 4	(22, 34, 47), (27, 39, 45), (34, 40, 54), (24, 33, 42)
	2	6, 15	(19, 27, 33), (18, 20, 24)
	3	1, 13	(13, 22, 27), (16, 24, 27)
	4	10, 13	(15, 22, 29), (16, 20, 23)
	5	4, 7	(33, 37, 41), (31, 35, 48)
10	6	5, 9	(8, 10, 12), (8, 12, 15)
	7	8, 12, 14	(32, 39, 44), (24, 32, 40), (23, 36, 50)
	8	12	(32, 44, 53)
	9	2, 3, 6, 9	(19, 23, 27), (15, 24, 30), (15, 21, 29), (15, 19, 26)
	10	3, 12	(42, 48, 66), (30, 45, 53)
	11	14	(11, 17, 22)
	1	1, 6	(30, 38, 47), (20, 30, 41)
	2	5, 8, 14, 15	(33, 39, 48), (33, 40, 51), (22, 36, 45), (27, 44, 59)
	3	3, 5, 11, 12, 13	(8, 11, 13), (9, 13, 15), (8, 9, 10), (10, 12, 13), (6, 8, 11)
	4	5, 6, 8, 13, 14	(16, 21, 26), (16, 23, 29), (18, 29, 37), (20, 27, 35), (18, 25, 28)
11	5	3, 4, 6	(22, 33, 39), (28, 31, 37), (23, 29, 36)
	6	2, 10	(21, 28, 33), (19, 27, 38)
	7	1, 14, 15	(30, 40, 49), (28, 42, 53), (35, 46, 63)
	8	2, 7, 9, 11, 14	(4, 6, 7), (5, 8, 10), (7, 10, 13), (9, 11, 12), (6, 7, 8)
	9	5, 9, 13	(30, 40, 45), (27, 39, 50), (31, 36, 40)
	1	1, 11	(22, 31, 36), (24, 29, 36)
	2	8, 15	(33, 46, 63), (27, 44, 57)
	3	5, 11	(4, 5, 6), (7, 11, 13)
	4	12	(29, 41, 53)
	5	15	(19, 24, 32)
	6	2, 13	(27, 42, 52), (39, 45, 62)
	7	8, 11	(14, 19, 22), (11, 15, 21)
	8	3, 12	(14, 18, 22), (17, 20, 27)
12	9	6, 14	(4, 5, 6), (5, 7, 9)
12	10	4	(15, 18, 22)
	11	7	(25, 39, 53)
	12	6, 10	(12, 13, 16), (5, 7, 9)
	13	2, 3	(22, 26, 35), (16, 22, 25)
	14	1, 8, 13	(4, 5, 6), (5, 8, 10), (7, 9, 11)
	15	9	(27, 39, 44)
	16	7, 10	(6, 10, 13), (11, 13, 14)
	17	4, 12	(25, 41, 46), (25, 38, 48)
	18	5, 9, 13	(17, 21, 29), (16, 22, 29), (15, 19, 23)

Job No.	Operation	Alternative machine	Processing time
	1	3, 4, 10, 11, 13	(35, 46, 61), (31, 47, 57), (31, 44, 57), (34, 41, 54), (35, 50, 59)
	2	5, 7, 10, 12, 15	(5, 8, 9), (3, 5, 7), (8, 10, 14), (10, 11, 13), (6, 9, 11)
	3	2,7	(13, 16, 19), (8, 12, 15)
	4	1, 3, 4, 6, 15	(5, 7, 8), (4, 5, 7), (8, 13, 16), (11, 12, 15), (7, 8, 10)
	5	6, 8, 10	(21, 26, 36), (17, 24, 29), (24, 28, 37)
	6	1, 9, 12, 15	(4, 5, 6), (3, 4, 5), (5, 7, 8), (8, 9, 11)
	7	2, 8, 10, 12	(23, 27, 34), (21, 30, 36), (28, 33, 42), (23, 29, 35)
	8	2	(32, 40, 47)
13	9	3, 6, 7, 9	(19, 23, 26), (16, 24, 33), (18, 29, 33), (16, 21, 29)
15	10	1, 2, 5, 9, 10	(9, 12, 15), (12, 14, 17), (14, 19, 26), (11, 18, 21), (13, 17, 20)
	11	10, 11, 12	(30, 47, 53), (43, 49, 61), (42, 50, 60)
	12	4, 6, 13	(39, 44, 57), (29, 38, 47), (27, 41, 48)
	13	1, 9, 10, 12	(16, 22, 26), (15, 21, 29), (14, 16, 21), (14, 18, 22)
	14	1, 4, 6, 13, 14	(11, 15, 18), (15, 18, 22), (10, 13, 15), (12, 14, 17), (13, 19, 21)
	15	3, 4, 6, 13	(5, 6, 7), (3, 4, 5), (4, 5, 6), (8, 9, 12)
	16	5, 11, 14	(11, 15, 19), (13, 18, 21), (9, 13, 15)
	17	8, 9, 11	(13, 15, 21), (11, 16, 22), (12, 19, 22)
	18	2, 15	(32, 44, 59), (41, 50, 64)
	1	3, 9	(38, 46, 63), (28, 43, 59)
	2	1, 2, 7, 12	(7, 10, 12), (13, 17, 23), (9, 11, 14), (9, 13, 18)
	3	4, 7, 8	(6, 8, 9), (6, 9, 11), (8, 10, 11)
	4	3, 6	(11, 18, 20), (18, 25, 30)
	5	4	(6, 9, 12)
	6	3, 4, 13	(25, 29, 35), (22, 27, 37), (20, 33, 37)
14	7	5, 9	(22, 30, 33), (23, 29, 32)
	8	2, 3	(7, 9, 10), (5, 8, 11)
	9	5, 12, 15	(12, 18, 23), (8, 10, 14), (14, 19, 26)
	10	9, 15	(22, 28, 32), (16, 25, 32)
	11	4, 14, 15	(26, 42, 54), (31, 43, 49), (37, 47, 55)
	12	9, 10, 14	(22, 35, 47), (26, 31, 38), (26, 29, 38)
	13	6, 10	(8, 9, 10), (5, 7, 10)
	1	1, 11	(17, 20, 25), (14, 18, 23)
	2	12, 13	(27, 41, 50), (27, 43, 58)
	3	6	(12, 17, 23)
	4	/	(5, 8, 11)
	5	2, 15	(8, 12, 10), (13, 15, 18)
	6	4, 5	(29, 48, 63), (35, 43, 52)
15	/	9	(35, 47, 58)
15	8	8, 12	(17, 28, 34), (22, 30, 36)
	9	2,8	(16, 18, 23), (19, 22, 30)
	10	14	(37, 50, 55)
	11	8, 13	$(5, 6, \delta), (4, 7, 9)$
	12	Э, / 1 <i>5 с</i>	(35, 48, 67), (35, 45, 56)
	13	1, 5, 6	(7, 9, 12), (8, 10, 14), (10, 11, 14)
	14	3, 0, 9	(17, 22, 20), (17, 24, 28), (19, 21, 28)
	15	1, 12	(31, 42, 50), (31, 47, 65)

Table A5Fuzzy processing time of Jobs 13–15.

Job No.	Operation	Alternative machine Processing time		
	1	1, 2, 11	(35, 43, 52), (31, 45, 53), (35, 41, 52)	
	2	7	(26, 32, 43)	
	3	3, 6, 13	(25, 33, 46), (27, 39, 50), (26, 35, 48)	
	4	4, 5, 7, 9, 15	(32, 40, 51), (28, 43, 49), (30, 41, 48), (37, 44, 51), (42, 49, 61)	
	5	1, 2, 6	(16, 25, 30), (24, 30, 39), (19, 31, 35)	
	6	6, 12	(4, 5, 6), (7, 10, 12)	
	7	3, 5, 6, 11, 14	(6, 7, 9), (4, 5, 6), (8, 9, 11), (6, 8, 9), (7, 10, 13)	
	8	5, 7, 9	(10, 16, 22), (14, 19, 24), (14, 18, 20)	
	9	1, 4, 9, 10, 13	(10, 12, 15), (11, 18, 25), (12, 14, 16), (6, 9, 12), (9, 15, 18)	
	10	6, 11, 12, 13, 14 (17, 19, 25), (8, 11, 13), (14, 16, 20), (14, 20, 27), (19, 21, 25)		
16	11	4, 7, 8, 10, 15	(24, 31, 35), (32, 39, 47), (19, 30, 39), (26, 33, 39), (32, 40, 48)	
	12	5, 8, 14	(22, 28, 35), (20, 27, 36), (17, 24, 27)	
	13	1, 4, 7, 8, 15	(42, 50, 60), (37, 44, 50), (30, 47, 64), (40, 49, 67), (35, 48, 62)	
	14	1, 2, 8, 9	(11, 17, 19), (13, 19, 21), (12, 20, 25), (16, 21, 28)	
	15	1, 2, 4, 5, 7	(6, 8, 10), (6, 7, 8), (4, 6, 7), (8, 9, 12), (8, 10, 12)	
	16	9.12	(3, 5, 6), (4, 6, 7)	
	17	1, 3, 8, 13, 15	(18, 21, 26), (15, 20, 26), (16, 24, 30), (17, 19, 21), (18, 23, 26)	
	18	4, 5, 10, 11	(21, 26, 30), (23, 27, 30), (15, 24, 28), (15, 18, 21)	
	19	3. 11	(16, 19, 21), (16, 20, 23)	
	20	7. 8. 14. 15	(16, 19, 26), (13, 20, 28), (11, 15, 20), (20, 23, 27)	
	21	2. 12. 14	(10, 12, 20), (10, 20, 20), (11, 10, 20), (20, 20, 21) (23, 27, 33), (23, 29, 35), (12, 14, 16)	
	1	10 11	(29, 46, 59) (30, 44, 53)	
	2	5 10	(12, 16, 20), (30, 11, 33)	
	3	4	(12, 10, 20), (10, 10, 17) (9, 11, 13)	
	4	12 13	$(11 \ 13 \ 16) \ (10 \ 14 \ 18)$	
	5	7 12	(10, 11, 13), (15, 17, 21)	
	6	13	(10, 11, 10), (10, 17, 21) (41, 46, 64)	
	7	2 11	(17, 23, 31) (16, 19, 22)	
	8	1 5 9	(17, 25, 51), (10, 19, 22) (14, 20, 22), (16, 18, 24), (13, 17, 23)	
	9	2 12	(14, 20, 22), (10, 10, 24), (13, 17, 23) (20, 29, 40), (19, 30, 39)	
	10	13 15	(12, 16, 20), (19, 30, 59)	
	10	15, 15	(12, 10, 20), (2, 13, 17)	
17	11	1 8	(17, 24, 55) (13, 21, 20) $(13, 17, 20)$	
	12	4, 0 6	(13, 21, 2), (13, 17, 20) (28, 23, 28)	
	13	136	(26, 55, 56) (13, 17, 21) $(0, 12, 15)$ $(12, 14, 15)$	
	14	1, 5, 0	(13, 17, 21), (7, 12, 13), (12, 14, 13) (5, 9, 11), (4, 7, 0)	
	15	J, 8	(3, 6, 11), (4, 7, 9)	
	10	1, 0	(3, 5, 0), (0, 8, 9)	
	17	$\frac{2}{2}$	(34, 42, 50) (12, 15, 10) $(15, 18, 21)$	
	18	2, 10	(15, 15, 19), (15, 18, 21)	
	19	1, 0	(4, 0, 7), (5, 7, 8)	
	20	9, 10, 11	(3, 5, 6), (8, 10, 12), (6, 9, 10)	
	21	3, 12	(10, 15, 18), (13, 18, 23)	
	22	0, 11, 14	(13, 19, 26), (11, 17, 22), (14, 20, 25)	
	1	3, 8	(5, 8, 11), (11, 13, 17)	
	2	5, 6, 8	(14, 16, 20), (8, 12, 16), (10, 13, 15)	
18	3	2	(19, 21, 28)	
	4	1, 5, 10	(8, 13, 16), (12, 16, 21), (14, 18, 22)	
	5	9	(13, 17, 24)	
	6	5, 8	(40, 46, 64), (42, 47, 65)	

Table A6Fuzzy processing times of Jobs 16–18.

(To be continued)

	Table	The Fully processing times of J	(Continued)
Job No.	Operation	Alternative machine	Processing time
	7	3, 7, 10	(39, 44, 55), (36, 48, 64), (38, 49, 64)
	8	5, 6, 13	(14, 17, 22), (10, 14, 16), (7, 10, 12)
	9	8, 15	(11, 16, 19), (9, 13, 15)
	10	3, 11, 15	(20, 28, 34), (17, 27, 34), (20, 30, 42)
	11	10, 13	(40, 48, 54), (43, 50, 56)
18	12	5, 13, 15	(25, 31, 40), (25, 32, 43), (25, 36, 41)
	13	3, 6, 9	(27, 30, 38), (23, 28, 36), (20, 26, 35)
	14	2	(8, 11, 12)
	15	1, 14	(10, 16, 18), (12, 18, 22)
	16	4, 15	(14, 18, 24), (16, 19, 24)
	17	3, 10, 14	(28, 36, 40), (22, 32, 41), (23, 35, 48)

Table A6Fuzzy processing times of Jobs 16–18.

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