

Some Upper Bounds and Exact Values on Linear Complexities Over \mathbb{F}_M of Sidelnikov Sequences for $M = 2$ and 3

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Abstract—Sidelnikov sequences, a kind of cyclotomic sequences with many desired properties such as low correlation and variable alphabet sizes, can be employed to construct a polyphase sequence family that has many applications in high-speed data communications. Recently, cyclotomic numbers have been used to investigate the linear complexity of Sidelnikov sequences, mainly about binary ones, although the limitation on the orders of the available cyclotomic numbers makes it difficult. This paper continues to study the linear complexity over \mathbb{F}_M of M -ary Sidelnikov sequence of period $q - 1$ using Hasse derivative, which implies $q = p^m$, $m \geq 1$ and $M|(q - 1)$. The t th Hasse derivative formulas are presented in terms of cyclotomic numbers, and some upper bounds on the linear complexity for $M = 2$ and 3 are obtained only with some additional restrictions on q . Furthermore, concrete illustrations for several families of these sequences, such as $q \equiv 1 \pmod{2}$ and $q \equiv 1 \pmod{3}$, show these upper bounds are tight and reachable; especially for $q = 2 \times 3^\lambda + 1$ ($1 \leq \lambda \leq 20$), the exact linear complexities over \mathbb{F}_3 of the ternary Sidelnikov sequences are determined; and it turns out that all the linear complexities of the sequences considered are very close to their periods.

Index Terms—Array structure, cyclotomic numbers, Hasse derivative, linear complexity, Sidelnikov sequences.

I. INTRODUCTION

PSEUDO random sequences with certain properties are widely used in the communication engineering and cryptography [1]. The cyclotomic sequences have a number of

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attractive randomness properties [2]–[4]. Ding [5] studied their linear complexity, minimal polynomial, and autocorrelation function.

For $q = p^m$ where p is an odd prime and m is a positive integer, Sidelnikov [6] introduced a kind of cyclotomic sequence called the M -ary Sidelnikov sequence of period $q - 1$ where $M|(q - 1)$. Soon afterwards, Lempel, Cohn and Eastman [7] re-introduced its binary form called Sidelnikov-Lempel-Cohn-Eastman sequence. In the last two decades, a lot of attention has been devoted to this binary sequence. For example, using the cyclotomic numbers, Helleseeth and Yang [8] originally investigated the autocorrelation function and linear complexity over \mathbb{F}_2 of the binary Sidelnikov sequences. Later on, Kyureghan and Pott [9], and Meidl and Winterhof [10] determined the exact linear complexity over \mathbb{F}_2 of some of these sequences with well-known results on cyclotomic numbers; a lower bound on the linear complexity profile of these sequences was also introduced in [10], which is the desirable important property of applications. Then, Wang [11] and Su [12] studied the linear complexity of binary cyclotomic sequences of order 6 and Legendre-Sidelnikov sequences of period $p(q - 1)$, respectively. Ye *et al.* [13] further studied the linear complexity of a new kind of binary cyclotomic sequence, with length p^r , and Liang *et al.* [14] computed the linear complexity of Ding-Helleseeth generalized cyclotomic sequences by using cyclotomic numbers of order 8. Following the footsteps of these pioneers, Zeng *et al.* [15] discussed the \mathbb{F}_M -linear complexity of M -ary Sidelnikov sequences of period $p - 1 = f \times M^\lambda$ where M is not just equal to 2. On the other hand, using the discrete Fourier transform (i.e., DFT), Helleseeth *et al.* [16], [17] also determined the linear complexity over \mathbb{F}_p of binary Sidelnikov sequences, and Garaev *et al.* [18] derived the lower bound of the linear complexity over \mathbb{F}_p . For the k -error linear complexity over \mathbb{F}_p of the d -ary Sidelnikov sequence, Chung and Yang [19] presented many results of interest, then Aly and Meidl [20] further complemented these results.

In this paper, we continue to study the linear complexity over \mathbb{F}_M of M -ary Sidelnikov sequence of period $q - 1$ using Hasse derivative, where $q = p^m$, $m \geq 1$ and $M|(q - 1)$. Some upper bounds on the linear complexity for $M = 2$ and 3 are obtained, and some exact values of the linear complexity for several families of these sequences, such as $q \equiv 1 \pmod{2}$ and $q \equiv 1 \pmod{3}$, illustrate these upper bounds are tight and reachable. In particular, the exact linear complexities over

\mathbb{F}_3 of the ternary Sidelnikov sequences are determined for $q = 2 \times 3^\lambda + 1$ ($1 \leq \lambda \leq 20$), and it turns out that all the linear complexities of the sequences considered are very close to their periods. Some examples over \mathbb{F}_2 have been confirmed for binary Sidelnikov sequences by Helleseth and Yang.

The rest of this paper is organized as follows. In Section II, after reviewing some notations and definitions, we present formulas for the t th Hasse derivative of the generating function $S(x)$ of the M -ary Sidelnikov sequence $\{s_n\}_{n \geq 0}$ in terms of the cyclotomic numbers. In Section III, the multiplicities of some r th primitive roots of unity over \mathbb{F}_M as roots of $S(x)$ are determined using the Hasse derivative to estimate the \mathbb{F}_M -linear complexity of M -ary Sidelnikov sequences for the two cases of $q \equiv 1 \pmod{2}$ and $q \equiv 1 \pmod{3}$. Some special examples are listed in Table I for $q = 2 \times 3^\lambda + 1$ ($1 \leq \lambda \leq 20$). Note that this section extends our conference version [15] by adding Subsection III-A on the case $q \equiv 1 \pmod{2}$ which includes Theorem 1 and Examples 4 and 5, by supplementing a main result on the case $q \equiv 1 \pmod{3}$ in Theorem 2 of Subsection III-B, and by giving all proofs of the relevant results here. In Section IV, there are some concluding remarks. In addition, some known cyclotomic numbers of orders 2, 2r, 3, 6 and 9 are displayed in Chapter IV due to the need to prove the results of this paper.

II. PRELIMINARIES

In this section, after some notations are listed, the M -ary Sidelnikov sequence, the \mathbb{F}_M -linear complexity and the cyclotomic number are defined in Definitions 1, 3 and 4, respectively. Lemma 3 presents Hasse derivatives in terms of cyclotomic numbers, and will be used to determine the \mathbb{F}_M -linear complexity of the M -ary Sidelnikov sequence.

- p : an odd prime.
- q : an odd prime power p^m with $m \geq 1$.
- \mathbb{F}_p and \mathbb{F}_q : the finite fields with p and q elements, respectively.
- M : M is a prime with $M|(q-1)$.
- α : a fixed primitive element of \mathbb{F}_q .
- $\{s_n\}_{n \geq 0}$: the M -ary Sidelnikov sequence of period $q-1$.
- $S(x)$: the generating function of $\{s_n\}_{0 \leq n \leq q-2}$.
- $\mathbb{F}_M[x]$: the polynomial ring over finite field \mathbb{F}_M .
- $R(\gamma)$: the multiplicity of a primitive r th root γ of unity over \mathbb{F}_M as a root of $S(x)$, where $\gamma = e^{j \frac{2\pi}{r}}$ and $j = \sqrt{-1}$.
- $LC(\cdot)$: the \mathbb{F}_M -linear complexity of a sequence. It is written as LC for short if the context is clear.
- $\text{Ind } x$: the index of $x \in \mathbb{F}_q$ to the base g modulo q [21].

The M -ary Sidelnikov sequence is defined as follows.

Definition 1 ([22]): For a fixed primitive element α of \mathbb{F}_q and $M|(q-1)$, let $D_k^{(\alpha)}$, $k = 0, 1, \dots, M-1$, be the disjoint subsets of \mathbb{F}_q defined as

$$D_k^{(\alpha)} = \{\alpha^{Mi+k} - 1 \mid 0 \leq i \leq \frac{q-1}{M} - 1\}.$$

The M -ary Sidelnikov sequence $\{s_n\}_{n \geq 0}$ of period $q-1$ is defined as

$$s_n = \begin{cases} k & \text{if } \alpha^n \in D_k^{(\alpha)}, \\ 0 & \text{if } \alpha^n = -1. \end{cases} \quad (1)$$

Equivalently,

$$s_n \equiv \log_\alpha(\alpha^n + 1) \pmod{M}. \quad (2)$$

Note that $\bigcup_{k=0}^{M-1} D_k^{(\alpha)} = \mathbb{F}_q \setminus \{-1\}$, $0 \in D_0^{(\alpha)}$, and $\log_\alpha(0) = 0$.

Example 1: Let $q = 7$ and $M = 3$. For $\alpha = 3$, we have a ternary Sidelnikov sequence $\{s_n\}_{n \geq 0}$ of period 6, that is, $\{s_n\}_{0 \leq n \leq 5} = \{2, 1, 1, 0, 2, 0\}$.

Definition 2: A polynomial $C(x) = x^L + c_1x^{L-1} + \dots + c_{L-1}x + 1 \in \mathbb{F}_M[x]$ is the connection polynomial of the M -ary sequence $\{s_n\}_{n \geq 0}$ of period $T = q-1$ if there exist constants $c_0 = 1, c_1, \dots, c_{L-1}, c_L = 1 \in \mathbb{F}_M$, such that

$$s_j \equiv - \sum_{i=0}^{L-1} c_i s_{j-L+i} \pmod{M}, \text{ for all } j \geq L. \quad (3)$$

Definition 3: The linear complexity over \mathbb{F}_M of $\{s_n\}_{n \geq 0}$ is defined as

$$LC(\{s_n\}_{n \geq 0}) = \min\{\deg(C(x)) : C(x) \text{ is the connection polynomial of } \{s_n\}_{n \geq 0}\}.$$

Lemma 1 ([23]): Let $S(x) = s_0 + s_1x + \dots + s_{q-2}x^{q-2}$. Then $C(x)$ is the connection polynomial of $\{s_n\}_{n \geq 0}$ if and only if $S(x)C(x) \equiv 0 \pmod{(x^{q-1} - 1)}$.

Therefore, the \mathbb{F}_M -linear complexity of the M -ary Sidelnikov sequence $\{s_n\}_{n \geq 0}$ can be determined by

$$LC(\{s_n\}_{n \geq 0}) = q-1 - \deg[\gcd(x^{q-1} - 1, S(x))], \quad (4)$$

where $S(x)$ is by (1)

$$S(x) = \sum_{k=1}^{M-1} k \sum_{\substack{\alpha^n \in D_k^{(\alpha)} \\ 0 \leq n \leq q-2}} x^n \in \mathbb{F}_M[x]. \quad (5)$$

Example 2: The \mathbb{F}_3 -linear complexity of the ternary Sidelnikov sequence $\{s_n\}_{n \geq 0}$ of period 6 in Example 1 is 5 since $\gcd(x^6 - 1, 2x^4 + x^2 + x + 2) = x - 1$.

Similar to Example 2, in order to evaluate $LC(\{s_n\}_{n \geq 0})$ from (4), we will determine the multiplicity of γ as a root of $S(x)$, where γ is also a $(q-1)$ -th root of unity over \mathbb{F}_M or in an extension field of \mathbb{F}_M , by using the cyclotomic numbers defined as follows.

Definition 4: Let α be a primitive element in the finite field \mathbb{F}_q , and $e|(q-1)$. Then the cyclotomic classes $C_u^{(\alpha)}$, $0 \leq u \leq e-1$, are defined in \mathbb{F}_q as

$$C_u^{(\alpha)} = \{\alpha^{ed+u} \mid 0 \leq d \leq \frac{q-1}{e} - 1\}.$$

For fixed positive integers u and v , not necessarily distinct, the cyclotomic number $(u, v)_e$ is defined as the number of elements $z_u \in C_u^{(\alpha)}$ such that $z_u + 1 \in C_v^{(\alpha)}$, where e is called the order of the cyclotomic number.

Example 3: Let $\mathbb{F}_7 = \{0, 1, 2, 3, 4, 5, 6\} = \{0, \alpha, \alpha^2, \alpha^3, \alpha^4, \alpha^5, \alpha^6\}$ where $\alpha = 3$ is a primitive element in \mathbb{F}_7 . Let $e = 3$. Then $C_0^{(\alpha)} = \{1, \alpha^3\}$, $C_1^{(\alpha)} = \{\alpha, \alpha^4\}$, and $C_2^{(\alpha)} = \{\alpha^2, \alpha^5\}$. It is easy to see that $1+1 = 2 = \alpha^2$ and $\alpha^3+1 = 0$, thus $(0, 0)_3 = 0$, $(0, 1)_3 = 0$, and $(0, 2)_3 = 1$. Similarly, $(1, 0)_3 = 0$, $(1, 1)_3 = 1$, $(1, 2)_3 = 1$, $(2, 0)_3 = 1$, $(2, 1)_3 = 1$, and $(2, 2)_3 = 0$.

Let $q = f \times M^\lambda + 1$ where f and λ are two positive integers. Then $x^{q-1} - 1 = (x^f - 1)^{M^\lambda} \in \mathbb{F}_M[x]$. Let γ be a primitive r th root of unity over \mathbb{F}_M or in an extension field of \mathbb{F}_M where $r|f$. Then the multiplicity of γ as a root of $S(x)$ is i , i.e.,

$$R(\gamma) = i,$$

if $S(\gamma) = S(\gamma)^{(1)} = \dots = S(\gamma)^{(i-1)} = 0$ and $S(\gamma)^{(i)} \neq 0$, where $S(x)^{(t)}$ ($t = 0, 1, \dots, i$) is the t th Hasse derivative of $S(x)$ [24], and defined as

$$S(x)^{(t)} = \sum_{k=1}^{M-1} k \sum_{\substack{\alpha^n \in D_k^{(\alpha)} \\ t \leq n \leq q-2}} \binom{n}{t} x^{n-t} \in \mathbb{F}_M[x], \quad (6)$$

where the binomial coefficients $\binom{n}{t}$ modulo M can be evaluated with the following Corollary 1.

Lemma 2 (Lucas' Theorem [25]): Let $0 \leq b_j \leq a_j \leq M-1$ for $j = 0, 1, \dots, l-1$, where M is a prime. Then

$$\binom{a_0 + a_1M + \dots + a_{l-1}M^{l-1}}{b_0 + b_1M + \dots + b_{l-1}M^{l-1}} \equiv \binom{a_0}{b_0} \binom{a_1}{b_1} \dots \binom{a_{l-1}}{b_{l-1}} \pmod{M}.$$

It is clear from Lemma 2 that $b_j \leq a_j$ for $j = 0, 1, \dots, l-1$. However, since there exists a convention that if $x < y$ then $\binom{x}{y} = 0$, the Lucas' theorem can be extended to the following corollary.

Corollary 1 (Extension of Lucas' Theorem): Let M be a prime.

1) Let $0 \leq b_j, a_j \leq M-1$ for $j = 0, 1, \dots, l-1$. Then

$$\binom{a_0 + a_1M + \dots + a_{l-1}M^{l-1}}{b_0 + b_1M + \dots + b_{l-1}M^{l-1}} \equiv \binom{a_0}{b_0} \binom{a_1}{b_1} \dots \binom{a_{l-1}}{b_{l-1}} \pmod{M}. \quad (7)$$

2) Let $n \equiv i \pmod{M^l}$ where $l = \lfloor \log_M(t) \rfloor + 1$ if $t \geq 1$, and $l = 1$ if $t = 0$. Then

$$\binom{n}{t} \equiv \binom{i}{t} \pmod{M}, \quad (8)$$

where $\binom{i}{t} = 0$ if $i < t$.

Proof: First, we prove Lucas' theorem is still true if there exists j such that $b_j > a_j$. To the end, we need to compare the coefficients of binomial expansion of $(1+x)^a$, where $a = \sum_{j=0}^{l-1} a_j M^j$ and a_0, \dots, a_{l-1} are the digits in the M -ary representation of a .

Since M is a prime, it follows that

$$(1+x)^{M^j} \equiv 1 + x^{M^j} \pmod{M} \text{ for } j \geq 1. \quad (9)$$

Then we have

$$\begin{aligned} & (1+x)^a \\ &= (1+x)^{a_0} ((1+x)^M)^{a_1} \dots ((1+x)^{M^{l-1}})^{a_{l-1}} \\ &\equiv (1+x)^{a_0} (1+x^M)^{a_1} \dots (1+x^{M^{l-1}})^{a_{l-1}} \pmod{M}. \end{aligned} \quad (10)$$

Let $b = \sum_{j=0}^{l-1} b_j M^j$ where b_0, \dots, b_{l-1} are the digits in the M -ary representation of b . It is clear that the items of x^b on

the left and right sides of (10) should be equal by using the unique M -ary representation property,

$$\begin{aligned} \binom{a}{b} x^b &= \binom{a_0}{b_0} x^{b_0} \binom{a_1}{b_1} x^{b_1 M} \dots \binom{a_{l-1}}{b_{l-1}} x^{b_{l-1} M^{l-1}} \\ &= \binom{a_0}{b_0} \binom{a_1}{b_1} \dots \binom{a_{l-1}}{b_{l-1}} x^b. \end{aligned} \quad (11)$$

Thus, if $b_j \leq a_j$ for all $0 \leq j \leq l-1$, the coefficient on the left of (11) must be congruent modulo M to the coefficient on the right, which is exactly the result of Lucas' theorem. Otherwise, if there exists $0 \leq j \leq l-1$ such that $b_j > a_j$, then $\binom{a_j}{b_j} = 0$, which means there is no item of $x^{b_j M^j}$ on the right of (11), leading to there is no item of x^b . Then, the coefficients on both sides of (11) are equal to 0, that is to say, the Lucas' theorem is also true for $b_j > a_j$ ($0 \leq j \leq l-1$). So, (7) is true.

Secondly, we prove (8) is also true.

(i) If $t = 0$, the result is obvious.

(ii) If $t \geq 1$, let $l = \lfloor \log_M(t) \rfloor + 1$. Then $t < M^l$. Let $n_0, n_1, \dots, n_{l-1}, n_l, \dots, n_{l'}$ and t_0, \dots, t_{l-1} be the digits in the M -ary representations of n and t , respectively. Then $n = \sum_{j=0}^{l'} n_j M^j$ and $t = \sum_{j=0}^{l-1} t_j M^j$. Since $n \equiv i \pmod{M^l}$, it follows that $i = \sum_{j=0}^{l-1} n_j M^j$. From (7), it is easy to see that

$$\begin{aligned} \binom{n}{t} &= \binom{n_0 + \dots + n_{l-1} M^{l-1} + n_l M^l + \dots + n_{l'} M^{l'}}{t_0 + \dots + t_{l-1} M^{l-1} + 0 \times M^l + \dots + 0 \times M^{l'}} \\ &\equiv \binom{n_0}{t_0} \dots \binom{n_{l-1}}{t_{l-1}} \binom{n_l}{0} \dots \binom{n_{l'}}{0} \pmod{M} \\ &\equiv \binom{i}{t} \pmod{M}. \end{aligned}$$

Thus, the proof completes. \square

Next, the t th Hasse derivatives ($t = 0, 1, \dots$) in terms of cyclotomic numbers are listed in the following lemma that will be used to determine the multiplicities of all the f th roots of unity, as the roots of $S(x)$.

Lemma 3: [15] Let $q = p^m \equiv 1 \pmod{M}$ where p is an odd prime and M is prime. Let γ be a primitive r th root of unity over \mathbb{F}_M or in an extension field of \mathbb{F}_M . $S(x)$ is the generating function of an M -ary Sidelnikov sequence $\{s_n\}_{0 \leq n \leq q-2}$. Then the t th Hasse derivatives $S(x)^{(t)} \in \mathbb{F}_M[x]$ ($t = 0, 1, \dots$) satisfy the following identities.

1)

$$S(1) = \sum_{k=1}^{M-1} k \sum_{u=0}^{M-1} (u, k)_M, \quad (12)$$

where $n \equiv u \pmod{M}$;

2)

$$S(1)^{(t)} = \sum_{k=1}^{M-1} k \sum_{i=t}^{M^l-1} \binom{i}{t} \sum_{j=0}^{M^{l-1}-1} (i, Mj+k)_{M^l}, \quad (13)$$

where $n \equiv i \pmod{M^l}$, and $l = \lfloor \log_M(t) \rfloor + 1$ if $t \geq 1$;

3) (see (14a) and (14b), as shown at the bottom of the next page);

4) (see (15a) and (15b), as shown at the bottom of the next page).

Remark 1: 1) The Hasse derivative in Lemma 3 is a bridge across the cyclotomic number and the linear complexity. Using this technique, one can determine the exact \mathbb{F}_M -linear complexity of an M -ary Sidelnikov sequence according to certain cyclotomic numbers. However, the well-known results on cyclotomic numbers are now just limited to the orders $e \leq 24$. This limitation hinders our ability to calculate the multiplicity of γ if r is large. So, it seems difficult to determine the exact \mathbb{F}_M -linear complexity. 2) For the proof details of Lemma 3, please refer to [10], [15].

III. UPPER BOUNDS AND SOME EXACT VALUES

This section investigates the \mathbb{F}_M -linear complexities of the M -ary Sidelnikov sequences. In the case of $q \equiv 1 \pmod{2}$, Theorem 1 shows that the \mathbb{F}_2 -linear complexities of binary Sidelnikov sequences of period $q - 1$ are upper bounded by $q - 2r$ if r satisfies certain conditions. In the case of $q \equiv 1 \pmod{3}$, for the trivial root 1, the primitive 2nd root and the primitive 3rd root of unity over \mathbb{F}_3 or in an extension field of \mathbb{F}_3 , the multiplicities of them as the roots of $S(x)$ are determined in Propositions 1, 2 and 3, respectively. Furthermore, the \mathbb{F}_3 -linear complexities of the ternary Sidelnikov sequences are presented in Theorem 2 and Corollary 2.

Note that, for the detailed meanings of the capital letters such as “A”, “B”, “C”, etc. in this section, please refer to Appendices IV.

A. Binary Case

The linear complexity of binary Sidelnikov sequence was originally investigated by Helleseth and Yang, and later extended by Kyureghyan and Pott, and Meidl and Winterhof. In the following theorem, we continue to estimate the linear complexity of the Sidelnikov sequence for $q \equiv 1 \pmod{2}$ using the technique introduced by Meidl and Winterhof, and the result is an extension of that in [10].

Theorem 1: Let $q = p^m \equiv 1 \pmod{2r}$ for $m = uv$, where p and r are both odd primes, $u \geq 1$, v is the order of p

modulo r , and v is even. Let 2 be a primitive root modulo r and $\{s_n\}_{n \geq 0}$ be a binary Sidelnikov sequence of period $q - 1$. Then the linear complexity of $\{s_n\}_{n \geq 0}$ over \mathbb{F}_2 is less than or equal to $q - 2r$ if

- 1) u is even; or
- 2) u is odd, and $4 \nmid v$ with $p \equiv 3 \pmod{4}$.

Proof: Let $S(x)$ be the generating function of $\{s_n\}_{0 \leq n \leq q-2}$. The multiplicities of 1 as a root of $S(x)$ have been intensively discussed in [9] and [10]. Here we consider the multiplicity of γ as a root of $S(x)$ where $\gamma (\neq 1)$ is a primitive r th root of unity in an extension field of \mathbb{F}_2 . Note that $(x^{2r} - 1) | (x^{q-1} - 1)$ since $2r | (q - 1)$. Let $x^r - 1 = (x - 1)(x^{r-1} + x^{r-2} + \dots + 1) = (x - 1)\Phi_r(x)$ where $\Phi_r(x) = \prod_{\substack{1 \leq k \leq r \\ \gcd(k, r) = 1}} (x - e^{2\pi i k / r})$ is a cyclotomic polynomial. Then $\Phi_r(x)$ is irreducible over \mathbb{F}_2 since 2 is a primitive root modulo r [26], and $\Phi_r(\gamma) = 0$. Let $M = 2$.

First, consider γ is a single root of $S(x)$. Let $n \equiv h \pmod{r}$. Then $1 \leq M - \lceil \frac{h+1}{r} \rceil < M$. Thus, we get from (14a) that

$$\begin{aligned} S(\gamma) &= \sum_{k=1}^{2-1} k \sum_{h=0}^{r-1} \sum_{j=0}^{r-1} \sum_{i=0}^1 (ir + h, 2j + k)_{2r} \gamma^h \\ &= \sum_{h=0}^{r-1} \left(\sum_{j=0}^{r-1} ((h, 2j + 1)_{2r} + (r + h, 2j + 1)_{2r}) \right) \gamma^h = \sum_{h=1}^{r-1} T_{(0, h)} \gamma^h, \end{aligned}$$

where $T_{(0, h)} = \sum_{j=0}^{r-1} ((h, 2j + 1)_{2r} + (r + h, 2j + 1)_{2r} - (0, 2j + 1)_{2r} - (r, 2j + 1)_{2r})$ since $\Phi_r(\gamma) = 0$. From Appendix B, it follows that

$$\begin{aligned} T_{(0, h)} &= \sum_{j=0}^{r-1} ((h, 2j + 1)_{2r} + (r + h, 2j + 1)_{2r} - (0, 2j + 1)_{2r} \\ &\quad - (r, 2j + 1)_{2r}) \\ &\stackrel{*}{=} B + (2r - 1)C - rB - B - (r - 1)C \equiv C - B \pmod{2} \\ &= \frac{1 - (-1)^u p^{(uv)/2}}{2r}, \end{aligned}$$

where $*$ means that for a fixed h , one and only one of h and $h + r$ is odd, and is taken once by $2j + 1$ when j runs from 0 to $r - 1$.

$$S(\gamma) = \begin{cases} \sum_{k=1}^{M-1} k \sum_{h=0}^{r-1} \sum_{j=0}^{r-1} \sum_{i=0}^{M-1} (ir + h, Mj + k)_{rM} \cdot \gamma^h & \text{if } M \neq r, \\ \sum_{k=1}^{M-1} k \sum_{h=0}^{M-1} (h, k)_M \cdot \gamma^h & \text{if } M = r, \end{cases} \quad (14a)$$

$$\quad \quad \quad (14b)$$

where $n = h \pmod{r}$;

$$S(\gamma)^{(t)} = \begin{cases} \sum_{k=1}^{M-1} k \sum_{i=t}^{M^l-1} \binom{i}{t} \sum_{h=0}^{r-1} \sum_{j=0}^{M^{l-1}-1} (u(i, h), Mj + k)_{rM^l} \cdot \gamma^h & \text{if } M \neq r, \\ \sum_{k=1}^{M-1} k \sum_{i=t}^{M^l-1} \binom{i}{t} \sum_{h=0}^{M-1} \sum_{j=0}^{M^{l-1}-1} (u(i, h), Mj + k)_{M^l} \cdot \gamma^h & \text{if } M = r, \end{cases} \quad (15a)$$

$$\quad \quad \quad (15b)$$

where $l = \lceil \log_M(t) \rceil + 1$ if $t \geq 1$, and $u(i, h)$ is (by the Chinese-Remainder-Theorem) the unique integer u satisfying $u - t \equiv h \pmod{r}$ and $u \equiv i \pmod{M^l}$, with $0 \leq u \leq rM^l - 1$ or $0 \leq u \leq M^l - 1$.

Since $\gamma, \dots, \gamma^{r-1}$ are linear independent over \mathbb{F}_2 , it follows that $S(\gamma) = 0$ if and only if

$$T_{(0,h)} \equiv 0 \pmod{2} \text{ for } h = 1, 2, \dots, r-1,$$

which means that

$$p^{(uv)/2} \equiv (-1)^u \pmod{4} \quad (16)$$

$$\equiv \begin{cases} 1 \pmod{4} & \text{if } u \text{ is even,} \\ 3 \pmod{4} & \text{if } u \text{ is odd.} \end{cases} \quad (17)$$

Then, we can get that

- 1) if u is even and v is even, then $p^{(uv)/2} \equiv 1 \pmod{4}$,
- 2) if u is odd, and v is even and $4 \nmid v$, then $p^{(uv)/2} \equiv 3 \pmod{4}$ if $p \equiv 3 \pmod{4}$.

So, in these two cases, γ is a root of $S(x)$, and $\Phi_r(x) \mid \gcd(x^{q-1} - 1, S(x))$.

Second, we consider whether γ is a double root of $S(x)$ in the above cases. According to Lemma 3, let $l = 1$ since $t = 1 < M = 2$. From (15a),

$$\begin{aligned} S(\gamma)^{(1)} &= \sum_{k=1}^{M-1} k \sum_{i=1}^{M-1} \binom{i}{1} \sum_{h=0}^{r-1} \sum_{j=0}^{r-1} (u(i, h), Mj + k)_{Mr} \gamma^h \\ &= \sum_{h=0}^{r-1} \sum_{j=0}^{r-1} (u(1, h), 2j + 1)_{2r} \gamma^h, \end{aligned}$$

where $u(1, h)$ is the unique integer u with $0 \leq u \leq 2r - 1$, $u - 1 \equiv h \pmod{r}$, and $u \equiv 1 \pmod{2}$. Then we have $u(1, h) = h + 1$ if h is even, and $u(1, h) = r + h + 1$ if h is odd. So,

$$S(\gamma)^{(1)} = \sum_{h=1}^{r-1} T_{(1,h)} \gamma^h,$$

where

$$\begin{aligned} T_{(1,h)} &= \sum_{j=0}^{r-1} ((u(1, h), 2j + 1)_{2r} - (1, 2j + 1)_{2r}) \\ &= \begin{cases} \sum_{j=0}^{r-1} ((h + 1, 2j + 1)_{2r} - (1, 2j + 1)_{2r}) & \text{if } h \text{ is even,} \\ \sum_{j=0}^{r-1} ((r + h + 1, 2j + 1)_{2r} - (1, 2j + 1)_{2r}) & \text{if } h \text{ is odd.} \end{cases} \end{aligned}$$

For any fixed h , $u(1, h)$ is always odd, and can be equal to $2j + 1$ once with j running from 0 through $r - 1$. This means that $T_{(1,h)} = B + (r - 1)C - (B + (r - 1)C) = 0$ for $h = 1, 2, \dots, r - 1$ according to Appendix B. Thus, γ is a double root of $S(x)$ in the above cases. In addition, since $q = (p^{(uv)/2})^2 \equiv 1 \pmod{4}$ from (17), we have $S(1) = (0, 1)_2 + (1, 1)_2 = \frac{q-1}{2} \equiv 0 \pmod{2}$ according to Appendix A, that is to say, 1 is a root of $S(x)$. Thus, the upper bound immediately follows. \square

Next, there are two examples to illustrate Theorem 1.

Example 4: Let $p = 5$, $u = 2$, $v = 2$, and $r = 3$. Then we have a binary Sidelnikov sequence $\{s_n\}_{n \geq 0}$ of period 624. From the proof of Theorem 1, it is clear that 1 is a root of $S(x)$, and for any $\gamma (\neq 1)$ being a primitive 3rd root of unity, γ is a double root of $S(x)$, which means $(x^2 + x + 1)^2 \mid S(x)$. Thus, the linear complexity of $\{s_n\}_{n \geq 0}$ over \mathbb{F}_2 is $LC(\{s_n\}_{n \geq 0}) \leq 619$ from Theorem 1 1). In addition, from Proposition 2 in [10], it is easy to see that the multiplicity of 1 as a root

of $S(x)$ is 9, so, $LC(\{s_n\}_{n \geq 0}) \leq 611$. Indeed, we can get $\gcd(x^{624} - 1, S(x)) = (x - 1)^{12}(x^2 + x + 1)^{10}$, that is, $LC(\{s_n\}_{n \geq 0}) = 592$.

Example 5: Let $p = 19$, $u = 1$, $v = 2$, and $r = 5$. Then we have a binary Sidelnikov sequence $\{s_n\}_{n \geq 0}$ of period 360. Similar to example 4, it is clear that 1 is a root of $S(x)$, and for any $\gamma (\neq 1)$ being a primitive 5th root of unity, γ is a double root of $S(x)$, which means $(x^4 + x^3 + x^2 + x + 1)^2 \mid S(x)$. Thus, the linear complexity of $\{s_n\}_{n \geq 0}$ over \mathbb{F}_2 is $LC(\{s_n\}_{n \geq 0}) \leq 351$ from Theorem 1 2). In addition, from Proposition 1 in [10], it is clear that the multiplicity of 1 as a root of $S(x)$ is 2, so, $LC(\{s_n\}_{n \geq 0}) \leq 350$. In fact, it follows that $\gcd(x^{360} - 1, S(x)) = (x - 1)^2(x^2 + x + 1)^6(x^4 + x^3 + x^2 + x + 1)^4(x^6 + x^3 + 1)^2$, that is, $LC(\{s_n\}_{n \geq 0}) = 318$.

Remark 2: 1) In the proof of Theorem 1, we make full use of the formulas in Appendix B for the cyclotomic numbers of order $2r$ over \mathbb{F}_q with $q = p^{uv} \equiv 1 \pmod{2r}$, where the order v of p modulo r is only even. Unfortunately, when v is odd, the cyclotomic problem is more intricate [27]. 2) In general, the determination of cyclotomic numbers of order e is difficult if e is not small [10], meaning that we can only utilize these formulas for small r . Here we consider the cases $r = 3$ and 5 as examples.

B. Ternary Case

In this subsection, let $q \equiv 1 \pmod{3}$ where q is a prime. For the trivial root 1, the primitive 2nd and 3rd roots of unity, the multiplicities of them as the roots of $S(x)$ are determined in Propositions 1, 2 and 3, respectively, by using the cyclotomic numbers of orders e 's (e.g, 3, 6 and 9). However, if e is not small, it is very difficult to calculate the cyclotomic numbers, so the determined values of the multiplicity R are not very large in these propositions. For all that, Theorem 2 and Corollary 2 present the \mathbb{F}_3 -linear complexities of the ternary Sidelnikov sequences, especially in the case of $q = 2 \times 3^\lambda + 1$ where λ is a positive integer.

Firstly, we determine the multiplicities of the trivial root 1 of unity, as a root of $S(x)$, in the case of $\text{Ind } 3 \equiv 0 \pmod{3}$.

Proposition 1: Let $q \equiv 1 \pmod{3}$ be a prime where $4q = c^2 + 27d^2$ and $c \equiv 1 \pmod{3}$. Let $q = (\sum_{i=0}^5 c_i \xi^i) (\sum_{i=0}^5 c_i \xi^{-i})$ where ξ is a primitive 9th root of unity of \mathbb{F}_q and $c_i (i = 1, 2, \dots, 5)$ are integers. $S(x)$ is the generating function of a ternary Sidelnikov sequence $\{s_n\}_{0 \leq n \leq q-2}$. Then in the case of $\text{Ind } 3 \equiv 0 \pmod{3}$, the multiplicity (R) of 1 as a root of $S(x)$ can be determined by

- 1) $R(1) = 1$ is trivial;
- 2) $R(1) = 2$ if and only if $q \equiv 1 \pmod{9}$;
- 3) $R(1) = 3$ if and only if $q \equiv 1 \pmod{9}$ and $d \equiv 0 \pmod{3}$;
- 4) $R(1) = 4$ if and only if $q \equiv 1 \pmod{9}$, $d \equiv 0 \pmod{3}$, and $c_1 + c_5 \equiv 2(c_2 + c_4) \pmod{9}$;
- 5) $R(1) = 5$ or 6 if and only if $q \equiv 1 \pmod{9}$, $d \equiv 0 \pmod{3}$, $c_1 + c_5 \equiv 2(c_2 + c_4) \pmod{9}$, and $c_4 \equiv 2c_1 \pmod{9}$;
- 6) $R(1) = 7$ if and only if $q \equiv 1 \pmod{9}$, $d \equiv 0 \pmod{3}$, $c_1 + c_5 \equiv 2(c_2 + c_4) \pmod{9}$, $c_4 = 2c_1 \pmod{9}$, and $c_3 + c_4 + c_5 \equiv 0 \pmod{9}$;

7) $R(1) = 8$ if and only if $q \equiv 1 \pmod{9}$, $d \equiv 0 \pmod{3}$, $c_1 + c_5 \equiv 2(c_2 + c_4) \pmod{9}$, $c_4 \equiv 2c_1 \pmod{9}$, $c_3 + c_4 + c_5 \equiv 0 \pmod{9}$, and $1 + c_0 + c_2 + 2c_4 \equiv 0 \pmod{9}$;

8) $R(1) = 9$ if and only if $q \equiv 1 \pmod{9}$, $d \equiv 0 \pmod{3}$, and $c_0 \equiv -1 + 2c_5$, $c_1 \equiv -c_5$, $c_2 \equiv 2c_5$, $c_3 \equiv c_5$, $c_4 \equiv -2c_5 \pmod{9}$.

Proof: Let $M = 3$. According to Appendix C, we have from (12) in Lemma 3 that $S(1) = (0, 1)_3 + (1, 1)_3 + (2, 1)_3 + 2((0, 2)_3 + (1, 2)_3 + (2, 2)_3) = 3(B + C + D) \equiv 0 \pmod{3}$, that is to say, 1 is always a root of $S(x)$.

According to Lemma 3, let $l = 1$ if $t = 1$ or 2. From (13), $S(1)^{(1)} = (1, 1)_3 + 2(2, 1)_3 + 2(1, 2)_3 + 4(2, 2)_3 = B + C + D = \frac{q-1}{3}$ and $S(1)^{(2)} = (2, 1)_3 + 2(2, 2)_3 = 2B + D = \frac{q-1}{3} - d$. Thus, $S(1)^{(1)} = 0$ if and only if $q \equiv 1 \pmod{9}$, and $S(1)^{(2)} = 0$ if and only if $q \equiv 1 \pmod{9}$ and $d \equiv 0 \pmod{3}$.

Similarly, let $l = 2$ if $t = 3, 4, 5, 6, 7, 8$. From (13),

$$S(1)^{(t)} = \sum_{k=1}^2 k \sum_{i=t}^{9-1} \binom{i}{t} \sum_{j=0}^{3-1} (i, 3j+k)_9. \quad (18)$$

Let $q = (\sum_{i=0}^5 c_i \xi^i)(\sum_{i=0}^5 c_i \xi^{-i})$ where ξ is a primitive 9th root of unity of \mathbb{F}_q and $c_i (i = 1, 2, \dots, 5)$ are integers. In the case $\text{Ind } 3 \equiv 0 \pmod{3}$, according to Appendix E, it follows that

$$\begin{aligned} & S(1)^{(3)} \\ & \equiv (3, 1)_9 + (3, 4)_9 + (3, 7)_9 + (4, 1)_9 + (4, 4)_9 + (4, 7)_9 + (5, 1)_9 \\ & \quad + (5, 4)_9 + (5, 7)_9 + 2(6, 1)_9 + 2(6, 4)_9 + 2(6, 7)_9 + 2(7, 1)_9 \\ & \quad + 2(7, 4)_9 + 2(7, 7)_9 + 2(8, 1)_9 + 2(8, 4)_9 + 2(8, 7)_9 + 2(3, 2)_9 \\ & \quad + 2(3, 5)_9 + 2(3, 8)_9 + 2(4, 2)_9 + 2(4, 5)_9 + 2(4, 8)_9 + 2(5, 2)_9 \\ & \quad + 2(5, 5)_9 + 2(5, 8)_9 + (6, 2)_9 + (6, 5)_9 + (6, 8)_9 + (7, 2)_9 \\ & \quad + (7, 5)_9 + (7, 8)_9 + (8, 2)_9 + (8, 5)_9 + (8, 8)_9 \pmod{3} \\ & \equiv B + 2C + 2E + F + 2J + K + M + 2N + 2P + 2Q + 2R \\ & \quad \pmod{3} \\ & \equiv \frac{c_1 - 2c_2 - 2c_4 + c_5}{3} \pmod{3}, \end{aligned} \quad (19)$$

$$\begin{aligned} & S(1)^{(4)} \\ & \equiv (4, 1)_9 + (4, 4)_9 + (4, 7)_9 + 2(5, 1)_9 + 2(5, 4)_9 + 2(5, 7)_9 \\ & \quad + 2(7, 1)_9 + 2(7, 4)_9 + 2(7, 7)_9 + (8, 1)_9 + (8, 4)_9 + (8, 8)_9 \\ & \quad + 2(4, 2)_9 + 2(4, 5)_9 + 2(4, 8)_9 + (5, 2)_9 + (5, 5)_9 + (5, 8)_9 \\ & \quad + (7, 2)_9 + (7, 5)_9 + (7, 8)_9 + 2(8, 2)_9 + 2(8, 5)_9 + 2(8, 8)_9 \\ & \quad \pmod{3} \\ & \equiv 2B + 2C + E + F + 2K + M + 2P + R \pmod{3} \\ & \equiv \frac{-2c_1 + c_4}{3} \pmod{3}, \end{aligned} \quad (20)$$

$$\begin{aligned} & S(1)^{(5)} \\ & \equiv (5, 1)_9 + (5, 4)_9 + (5, 7)_9 + 2(8, 1)_9 + 2(8, 4)_9 + 2(8, 7)_9 \\ & \quad + 2(5, 2)_9 + 2(5, 5)_9 + 2(5, 8)_9 + (8, 2)_9 + (8, 5)_9 + (8, 8)_9 \\ & \quad \pmod{3} \\ & \equiv B + 2E + J + K + N + Q + 2R \pmod{3} \end{aligned}$$

$$\equiv \frac{3c_1 - 2c_2 - 3c_4 + c_5}{6} \pmod{3}, \quad (21)$$

$$\begin{aligned} & S(1)^{(6)} \\ & \equiv (6, 1)_9 + (6, 4)_9 + (6, 7)_9 + (7, 1)_9 + (7, 4)_9 + (7, 7)_9 + (8, 1)_9 \\ & \quad + (8, 4)_9 + (8, 7)_9 + 2(6, 2)_9 + 2(6, 5)_9 + 2(6, 8)_9 + 2(7, 2)_9 \\ & \quad + 2(7, 5)_9 + 2(7, 8)_9 + 2(8, 2)_9 + 2(8, 5)_9 + 2(8, 8)_9 \pmod{3} \\ & \equiv 2B + C + J + 2M + N + Q + R \pmod{3} \\ & \equiv \frac{c_1 + 2c_2 - 2c_3 - c_4 - 3c_5}{6} \pmod{3}, \end{aligned} \quad (22)$$

$$\begin{aligned} & S(1)^{(7)} \\ & \equiv (7, 1)_9 + (7, 4)_9 + (7, 7)_9 + 2(8, 1)_9 + 2(8, 4)_9 + 2(8, 7)_9 \\ & \quad + 2(7, 2)_9 + 2(7, 5)_9 + 2(7, 8)_9 + (8, 2)_9 + (8, 5)_9 + (8, 8)_9 \\ & \quad \pmod{3} \\ & \equiv B + C + K + 2N + O + P + Q + S \pmod{3} \\ & \equiv \frac{2 + 2c_0 - 4c_1 - 4c_2 - c_3 + 5c_4 + 2c_5}{18} \pmod{3}, \end{aligned} \quad (23)$$

$$\begin{aligned} & S(1)^{(8)} \\ & \equiv (8, 1)_9 + (8, 4)_9 + (8, 7)_9 + 2(8, 2)_9 + 2(8, 5)_9 + 2(8, 8)_9 \\ & \equiv 2B + 2J + 2K + N + 2O \pmod{3} \\ & \equiv \frac{-2 - 2c_0 + 7c_1 - 2c_2 + 4c_3 - 5c_4 + c_5}{9} \pmod{3}. \end{aligned} \quad (24)$$

So, $S(1)^{(3)} = 0$ if and only if $c_1 + c_5 \equiv 2(c_2 + c_4) \pmod{9}$. $S(1)^{(4)} = 0$ if and only if $2c_1 \equiv c_4 \pmod{9}$. $S(1)^{(5)} = 0$ if and only if $c_2 + c_5 \equiv 3(c_2 + c_4 - c_1) \pmod{18}$. Fortunately, $S(1)^{(3)} = 0$ and $S(1)^{(4)} = 0$ imply that $S(1)^{(5)} = 0$. $S(1)^{(6)} = 0$ if and only if $c_4 + c_5 - c_1 \equiv 2(c_2 - c_3 - c_5) \pmod{18}$. However, it follows from (19), (20), and (22) that $c_3 + c_4 + c_5 \equiv 0 \pmod{9}$. $S(1)^{(7)} = 0$ if and only if $2(1 + c_0 + c_5) + c_4 \equiv 4(c_1 + c_2 - c_4) + c_3 \pmod{54}$. From (19), (20), (22), and (23), it can be reduced to $1 + c_0 + c_2 + 2c_4 \equiv 0 \pmod{9}$. $S(1)^{(8)} = 0$ if and only if $2(1 + c_0 + c_2) + 5c_4 \equiv 7c_1 + 4c_3 + c_5 \pmod{27}$. Similarly, from (19), (20), (22), (23), and (24), we have $c_0 \equiv -1 + 2c_5$, $c_1 \equiv -c_5$, $c_2 \equiv 2c_5$, $c_3 \equiv c_5$, $c_4 \equiv -2c_5 \pmod{9}$. \square

Secondly, note that if $q \equiv 1 \pmod{2}$, then $2^{q-1} - 1 \equiv 0 \pmod{3}$, which implies that 2 is a root of $x^{q-1} - 1$ over $\mathbb{F}_3 = \{0, 1, 2\}$. In the following proposition, the multiplicity of 2 as a root of $S(x)$ is presented for the case that $q \equiv 1 \pmod{6}$ where $q = a^2 + 3b^2$ and $a \equiv 1 \pmod{3}$.

Proposition 2: Let $q = 6f + 1$ be a prime where $q = a^2 + 3b^2$ and $a \equiv 1 \pmod{3}$. Let $\gamma (\neq 1)$ be a primitive 2nd root of unity over \mathbb{F}_3 . $S(x)$ is the generating function of a ternary Sidelnikov sequence $\{s_n\}_{0 \leq n \leq q-2}$. Then in the case of f being even, the multiplicity (R) of γ as a root of $S(x)$ can be determined by

1) $R(\gamma) = 1$ or 2 if and only if $b \equiv 0 \pmod{3}$;

2) (see the bottom line of the next page.)

Proof: Let $q = 6f + 1$ where f is even. Since $\gamma \neq 1$ and $\gamma^2 = 1$, $1 + \gamma = 0$. The Appendix D lists the cyclotomic numbers of order 6, which distinguish among the following three cases: $\text{Ind } 2 \equiv 0 \pmod{6}$, $\text{Ind } 2 \equiv 2$ or $5 \pmod{6}$, and $\text{Ind } 2 \equiv 1$ or $4 \pmod{6}$. Let $M = 3$ and $r = 2$.

We will determine the multiplicities of γ as a root of $S(x)$ using Appendix D.

1) First, consider the case $R(\gamma) = 1$. From (14a), it follows that

$$\begin{aligned} S(\gamma) &= \sum_{k=1}^{3-1} k \sum_{h=0}^{2-1} \sum_{j=0}^{2-1} \sum_{i=0}^{3-1} (2i+h, 3j+k)_6 \gamma^h \\ &= \sum_{h=0}^1 \sum_{j=0}^1 \sum_{i=0}^2 ((2i+h, 3j+1)_6 + 2(2i+h, 3j+2)_6) \gamma^h \\ &= \sum_{h=1}^1 T_{(0,h)} \gamma^h, \end{aligned}$$

where $T_{(0,h)} = \sum_{j=0}^1 \sum_{i=0}^2 ((2i+h, 3j+1)_6 + 2(2i+h, 3j+2)_6 - (2i, 3j+1)_6 - 2(2i, 3j+2)_6)$. According to Appendix D,

$$\begin{aligned} T_{(0,1)} &\equiv (1, 1)_6 + 2(1, 2)_6 - (0, 1)_6 - 2(0, 2)_6 + (3, 1)_6 \\ &\quad + 2(3, 2)_6 - (2, 1)_6 - 2(2, 2)_6 + (5, 1)_6 + 2(5, 2)_6 \\ &\quad - (4, 1)_6 - 2(4, 2)_6 + (1, 4)_6 + 2(1, 5)_6 - (0, 4)_6 \\ &\quad - 2(0, 5)_6 + (3, 4)_6 + 2(3, 5)_6 - (2, 4)_6 - 2(2, 5)_6 \\ &\quad + (5, 4)_6 + 2(5, 5)_6 - (4, 4)_6 - 2(4, 5)_6 \\ &\equiv B - F - H + I \pmod{3} \\ &= \begin{cases} b & \text{if } \text{Ind } 2 \equiv 0 \pmod{6} \\ b & \text{if } \text{Ind } 2 \equiv 2 \text{ or } 5 \pmod{6} \\ b & \text{if } \text{Ind } 2 \equiv 1 \text{ or } 4 \pmod{6} \end{cases} \\ &\equiv 0 \pmod{3} \text{ for all cases if } b \equiv 0 \pmod{3}, \end{aligned}$$

which implies that γ is a single root of $S(x)$ if and only if $b \equiv 0 \pmod{3}$.

Second, consider the case $R(\gamma) = 2$. From (15a), it follows that

$$\begin{aligned} S(\gamma)^{(1)} &= \sum_{k=1}^{3-1} k \sum_{i=1}^{3-1} \binom{i}{1} \sum_{h=0}^{2-1} \sum_{j=0}^{2-1} (u(i, h), 3j+k)_6 \gamma^h \\ &= \sum_{h=0}^1 \sum_{i=1}^2 \binom{i}{1} \sum_{j=0}^1 ((u(i, h), 3j+1)_6 + 2(u(i, h), 3j+2)_6) \gamma^h \\ &= \sum_{h=1}^1 T_{(1,h)} \gamma^h, \end{aligned}$$

where $u(i, h)$ is the unique integer u with $0 \leq u \leq 5$, $u - 1 \equiv h \pmod{2}$, and $u \equiv i \pmod{3}$; $T_{(1,h)} = \sum_{j=0}^1 ((u(1, h), 3j+1)_6 + 2(u(1, h), 3j+2)_6 + 2(u(2, h), 3j+1)_6 + 4(u(2, h), 3j+2)_6 - (1, 3j+1)_6 - 2(1, 3j+2)_6 - 2(5, 3j+1)_6 - 4(5, 3j+2)_6)$. According to Appendix D, it follows that

$$T_{(1,1)} \equiv \sum_{j=0}^1 ((4, 3j+1)_6 + 2(4, 3j+2)_6 + 2(2, 3j+1)_6$$

$$\begin{aligned} &+ 4(2, 3j+2)_6 - (1, 3j+1)_6 - 2(1, 3j+2)_6 \\ &- 2(5, 3j+1)_6 - 4(5, 3j+2)_6) \\ &\equiv J - G + E - F + C - B \pmod{3} \\ &= \begin{cases} 0 & \text{if } \text{Ind } 2 \equiv 0 \pmod{6} \\ 0 & \text{if } \text{Ind } 2 \equiv 2 \text{ or } 5 \pmod{6} \\ 0 & \text{if } \text{Ind } 2 \equiv 1 \text{ or } 4 \pmod{6} \end{cases} \\ &\equiv 0 \pmod{3}, \end{aligned}$$

which implies that if γ is a root of $S(x)$, it must be a double root of $S(x)$.

2) Consider the case $R(\gamma) = 3$. Similar to above,

$$\begin{aligned} S(\gamma)^{(2)} &= \sum_{k=1}^{3-1} k \sum_{i=2}^{3-1} \binom{i}{2} \sum_{h=0}^{2-1} \sum_{j=0}^{2-1} (u(i, h), 3j+k)_6 \gamma^h \\ &= \sum_{h=0}^1 \sum_{i=2}^2 \binom{i}{2} \sum_{j=0}^1 ((u(i, h), 3j+1)_6 + 2(u(i, h), 3j+2)_6) \gamma^h \\ &= \sum_{h=0}^1 \sum_{j=0}^1 ((u(2, h), 3j+1)_6 + 2(u(2, h), 3j+2)_6) \gamma^h \\ &= \sum_{h=1}^1 T_{(2,h)} \gamma^h, \end{aligned}$$

where $T_{(2,h)} = \sum_{j=0}^1 ((u(2, h), 3j+1)_6 + 2(u(2, h), 3j+2)_6 - (u(2, 0), 3j+1)_6 - 2(u(2, 0), 3j+2)_6)$. According to Appendix D, it follows that

$$\begin{aligned} T_{(2,1)} &= \sum_{j=0}^1 ((5, 3j+1)_6 + 2(5, 3j+2)_6 - (2, 3j+1)_6 - 2(2, 3j+2)_6) \\ &\equiv E - B + G - J \pmod{3} \\ &= \begin{cases} \frac{-2b}{3} & \text{if } \text{Ind } 2 \equiv 0 \pmod{6} \\ \frac{-a+b}{3} & \text{if } \text{Ind } 2 \equiv 2 \text{ or } 5 \pmod{6} \\ \frac{-2a-2b}{3} & \text{if } \text{Ind } 2 \equiv 1 \text{ or } 4 \pmod{6} \end{cases} \\ &\equiv \begin{cases} 0 \pmod{3} & \text{if } \text{Ind } 2 \equiv 0 \pmod{6} \text{ and } b \equiv 0 \pmod{9} \\ 0 \pmod{3} & \text{if } \text{Ind } 2 \equiv 2 \text{ or } 5 \pmod{6} \\ & \text{and } a \equiv b \pmod{9} \\ 0 \pmod{3} & \text{if } \text{Ind } 2 \equiv 1 \text{ or } 4 \pmod{6} \\ & \text{and } a \equiv -b \pmod{9}, \end{cases} \end{aligned}$$

which completes the proof. \square

Thirdly, it is worth noting from Proposition 1 that $q \equiv 1 \pmod{9}$ is one of necessary and sufficient conditions of 1 as a multiple root of $S(x)$. Then, we are interested in the multiplicity of γ (a primitive 3rd root of unity in an extension field of \mathbb{F}_3) as a root of $S(x)$.

Proposition 3: Let $q \equiv 1 \pmod{9}$ be a prime where $4q = c^2 + 27d^2$ and $c \equiv 7 \pmod{9}$.

$$R(\gamma) = 3 \text{ if and only if } \begin{cases} b \equiv 0 \pmod{9} & \text{if } \text{Ind } 2 \equiv 0 \pmod{6}, \\ a \equiv b \pmod{9} \text{ and } b \equiv 0 \pmod{3} & \text{if } \text{Ind } 2 \equiv 2 \text{ or } 5 \pmod{6}, \\ a \equiv -b \pmod{9} \text{ and } b \equiv 0 \pmod{3} & \text{if } \text{Ind } 2 \equiv 1 \text{ or } 4 \pmod{6}, \end{cases}$$

where $\text{Ind } 2$ means the index of 2 to a base g modulo q .

Let $q = (\sum_{i=0}^5 c_i \xi^i)(\sum_{i=0}^5 c_i \xi^{-i})$ where ξ is a primitive 9th root of unity of \mathbb{F}_q and $c_i (i = 1, 2, \dots, 5)$ are integers. Let $\gamma (\neq 1)$ be a primitive 3rd root of unity in an extension field of \mathbb{F}_3 . $S(x)$ is the generating function of a ternary Sidelnikov sequence $\{s_n\}_{0 \leq n \leq q-2}$. Then, the multiplicity (R) of γ as a root of $S(x)$ can be determined by

- 1) $R(\gamma) = 1$ if and only if $c \equiv 7 \pmod{18}$ and $d \equiv 1 \pmod{2}$, or $c \equiv 16 \pmod{18}$ and $d \equiv 0 \pmod{2}$;
- 2) $R(\gamma) = 2$ or 3 if and only if $c \equiv 7 \pmod{18}$ and $d \equiv 3 \pmod{6}$, or $c \equiv 16 \pmod{18}$ and $d \equiv 0 \pmod{6}$;
- 3) $R(\gamma) = 4$ if and only if $c \equiv 7 \pmod{18}$ and $d \equiv 3 \pmod{6}$ (or $c \equiv 16 \pmod{18}$ and $d \equiv 0 \pmod{6}$), and $c_4 \equiv 2c_1 \pmod{9}$ and $c_5 \equiv 3c_1 + 2c_2 \pmod{18}$;
- 4) $R(\gamma) = 5$ if and only if $c \equiv 7 \pmod{18}$ and $d \equiv 3 \pmod{6}$ (or $c \equiv 16 \pmod{18}$ and $d \equiv 0 \pmod{6}$), and $c_1 \equiv c_2 \equiv c_4 \equiv c_5 \equiv 0 \pmod{9}$.

Proof: Since $\gamma (\neq 1)$ is a primitive 3rd root of unity, $1 + \gamma + \gamma^2 = 0$.

1) From (14b), we get

$$\begin{aligned} S(\gamma) &= \sum_{k=1}^{3-1} k \sum_{h=0}^{3-1} (h, k)_3 \gamma^h = \sum_{h=0}^2 ((h, 1)_3 + 2(h, 2)_3) \gamma^h \\ &= \sum_{h=1}^2 T_{(0,h)} \gamma^h, \end{aligned}$$

where $T_{(0,h)} = (h, 1)_3 + 2(h, 2)_3 - (0, 1)_3 - 2(0, 2)_3$. Let $T_{(0,1)}, T_{(0,2)} \equiv 0 \pmod{3}$, and then according to Appendix C,

$$\begin{cases} T_{(0,1)} = \frac{2+c}{3} \equiv 0 \pmod{3} \\ T_{(0,2)} = \frac{2+c-9d}{6} \equiv 0 \pmod{3}. \end{cases}$$

Then, $S(\gamma) = 0$ if and only if $c \equiv 7 \pmod{18}$ and $d \equiv 1 \pmod{2}$, or $c \equiv 16 \pmod{18}$ and $d \equiv 0 \pmod{2}$.

2) According to Lemma 3, let $l = 1$ since $t = 1$. Then $M^{l-1} - \lceil \frac{k+1}{M} \rceil = 0$. From (15b),

$$\begin{aligned} S(\gamma)^{(1)} &= \sum_{k=1}^{3-1} k \sum_{i=1}^{3-1} \binom{i}{1} \sum_{h=0}^{3-1} (u(i, h), k)_3 \gamma^h \\ &= \sum_{h=0}^2 \sum_{i=1}^2 \binom{i}{1} ((u(i, h), 1)_3 + 2(u(i, h), 2)_3) \gamma^h \\ &= \sum_{h=1}^2 \sum_{i=1}^2 \binom{i}{1} ((u(i, h), 1)_3 + 2(u(i, h), 2)_3 \\ &\quad - (u(i, 0), 1)_3 - 2(u(i, 0), 2)_3) \gamma^h \\ &= \sum_{h=1}^2 T_{(1,h)} \gamma^h, \end{aligned}$$

where $u(i, h)$ is the unique integer u with $0 \leq u \leq 2$, $u - 1 \equiv h \pmod{3}$, and $u \equiv i \pmod{3}$. Let $T_{(1,1)} = 0$ and $T_{(1,2)} = 0$. It follows from Appendix C that

$$\begin{aligned} T_{(1,1)} &= -(1, 1)_3 - 2(1, 2)_3 + 2(2, 1)_3 + 4(2, 2)_3 \\ &\equiv B - C \pmod{3} \\ &\equiv -d \pmod{3}, \\ T_{(1,2)} &= -(1, 1)_3 - 2(1, 2)_3 \end{aligned}$$

$$\begin{aligned} &\equiv -C - 2D \pmod{3} \\ &\equiv \frac{-2q - c - 3d}{6} \pmod{3}. \end{aligned}$$

Thus, $S(\gamma)^{(1)} = 0$ if and only if $d \equiv 0 \pmod{3}$ and $-2q - c - 3d \equiv 0 \pmod{18}$. From $S(\gamma) = 0$ and $S(\gamma)^{(1)} = 0$, it follows that $R(\gamma) = 2$ if and only if $c \equiv 7 \pmod{18}$ and $d \equiv 1 \pmod{2}$, or $c \equiv 16 \pmod{18}$ and $d \equiv 0 \pmod{2}$.

Similarly, let $l = 1$ since $t = 2$. It follows from (15b) that

$$\begin{aligned} S(\gamma)^{(2)} &= \sum_{k=1}^2 k \sum_{i=2}^2 \binom{i}{2} \sum_{h=0}^2 (u(i, h), k)_3 \gamma^h \\ &= \sum_{h=0}^2 ((u(2, h), 1)_3 + 2(u(2, h), 2)_3) \gamma^h \\ &= \sum_{h=1}^2 T_{(2,h)} \gamma^h, \end{aligned}$$

where $T_{(2,h)} = (u(2, h), 1)_3 + 2(u(2, h), 2)_3 - (u(2, 0), 1)_3 - 2(u(2, 0), 2)_3$. Let $T_{(2,h)} = 0 (h = 1, 2)$. Then

$$\begin{aligned} T_{(2,1)} = T_{(2,2)} &= -(2, 1)_3 - 2(2, 2)_3 \\ &\equiv -D - 2B \pmod{3} \\ &\equiv -\frac{q-1-3d}{3} \pmod{3}. \end{aligned}$$

Thus, $S(\gamma)^{(2)} = 0$ if and only if $q - 1 - 3d \equiv 0 \pmod{9}$, that is, $S(\gamma)^{(2)} = 0$ if and only if $d \equiv 0 \pmod{3}$. It is clear that $S(\gamma)^{(1)} = 0$ implies $S(\gamma)^{(2)} = 0$.

3) According to Lemma 3, let $l = 2$ since $t = 3$. From (15b),

$$\begin{aligned} S(\gamma)^{(3)} &= \sum_{k=1}^2 k \sum_{i=3}^8 \binom{i}{3} \sum_{h=0}^2 \sum_{j=0}^{3-1} (u(i, h), 3j+k)_9 \gamma^h \\ &= \sum_{h=0}^2 \sum_{j=0}^2 \sum_{i=3}^8 \binom{i}{3} ((u(i, h), 3j+1)_9 + 2(u(i, h), 3j+2)_9) \gamma^h \\ &= \sum_{h=1}^2 T_{(3,h)} \gamma^h, \end{aligned}$$

where

$$\begin{aligned} T_{(3,h)} &= \sum_{j=0}^2 \sum_{i=3}^8 \binom{i}{3} ((u(i, h), 3j+1)_9 + 2(u(i, h), 3j+2)_9 \\ &\quad - (u(i, 0), 3j+1)_9 - 2(u(i, 0), 3j+2)_9), \end{aligned}$$

where $u(i, h)$ is the unique integer u with $0 \leq u \leq 8$, $u - 1 \equiv h \pmod{3}$, and $u \equiv i \pmod{9}$. Let $T_{(3,h)} = 0$ for $h = 1, 2$. According to Appendix E,

$$\begin{aligned} T_{(3,1)} &= \sum_{j=0}^2 ((\binom{4}{3})(4, 3j+1)_9 + (\binom{7}{3})(7, 3j+1)_9 + 2(\binom{4}{3})(4, 3j+2)_9 \\ &\quad + 2(\binom{7}{3})(7, 3j+2)_9 - (\binom{3}{3})(3, 3j+1)_9 - (\binom{6}{3})(6, 3j+1)_9) \end{aligned}$$

$$\begin{aligned}
& -2 \binom{3}{3} (3, 3j+2)_9 - 2 \binom{6}{3} (6, 3j+2)_9 \\
& \equiv 2C + F + J - 2M + N - P + Q \pmod{3} \\
& = \frac{-c_1 - 2c_2 - c_4 + c_5}{6} \\
& \equiv 0 \pmod{3},
\end{aligned}$$

$$\begin{aligned}
& T_{(3,2)} \\
& = \sum_{j=0}^2 \left(\binom{5}{3} (5, 3j+1)_9 + \binom{8}{3} (8, 3j+1)_9 + 2 \binom{5}{3} (5, 3j+2)_9 \right. \\
& \quad + 2 \binom{8}{3} (8, 3j+2)_9 - \binom{3}{3} (3, 3j+1)_9 - \binom{6}{3} (6, 3j+1)_9 \\
& \quad \left. - 2 \binom{3}{3} (3, 3j+2)_9 - 2 \binom{6}{3} (6, 3j+2)_9 \right) \\
& \equiv B + 2E + J + K + N + Q - R \pmod{3} \\
& = \frac{3c_1 - 2c_2 - 3c_4 + c_5}{6} \\
& \equiv 0 \pmod{3}.
\end{aligned}$$

Thus, $S(\gamma)^{(3)} = 0$ if and only if $c_4 \equiv 2c_1 \pmod{9}$ and $c_5 \equiv 3c_1 + 2c_2 \pmod{18}$.

4) Similar to 3), let $l = 2$ since $t = 4$. From (15b),

$$\begin{aligned}
& S(\gamma)^{(4)} \\
& = \sum_{k=1}^2 k \sum_{i=4}^8 \binom{i}{4} \sum_{h=0}^2 \sum_{j=0}^{3-1} (u(i, h), 3j+k)_9 \gamma^h \\
& = \sum_{h=0}^2 \sum_{j=0}^2 \sum_{i=4}^8 \binom{i}{4} ((u(i, h), 3j+1)_9 + 2(u(i, h), 3j+2)_9) \gamma^h \\
& = \sum_{h=1}^2 T_{(4,h)} \gamma^h.
\end{aligned}$$

Let $T_{(4,h)} = 0$ for $h = 1, 2$.

$$\begin{aligned}
& T_{(4,1)} \\
& = \sum_{j=0}^2 \left(\binom{5}{4} (5, 3j+1)_9 + \binom{8}{4} (8, 3j+1)_9 + 2 \binom{5}{4} (5, 3j+2)_9 \right. \\
& \quad + 2 \binom{8}{4} (8, 3j+2)_9 - \binom{4}{4} (4, 3j+1)_9 - \binom{7}{4} (7, 3j+1)_9 \\
& \quad \left. - 2 \binom{4}{4} (4, 3j+2)_9 - 2 \binom{7}{4} (7, 3j+2)_9 \right) \\
& \equiv 2B - 2C + E - F + J + 2K - M + N - 2P - 2Q + R \pmod{3} \\
& = \frac{c_1 + 4c_2 + c_4 - 5c_5}{6} \\
& \equiv 0 \pmod{3},
\end{aligned}$$

$$\begin{aligned}
& T_{(4,2)} \\
& = \sum_{j=0}^2 \left(\binom{6}{4} (6, 3j+1)_9 + 2 \binom{6}{4} (6, 3j+2)_9 - \binom{4}{4} (4, 3j+1)_9 \right.
\end{aligned}$$

$$\begin{aligned}
& \left. - \binom{7}{4} (7, 3j+1)_9 - 2 \binom{4}{4} (4, 3j+2)_9 - 2 \binom{7}{4} (7, 3j+2)_9 \right) \\
& \equiv -2C - F - J + 2M - N - 2P - Q \pmod{3} \\
& = \frac{2c_1 + c_2 - c_4 - 2c_5}{3} \\
& \equiv 0 \pmod{3}.
\end{aligned}$$

Thus, $S(\gamma)^{(4)} = 0$ if and only if $c_1 + c_4 - c_5 \equiv 4(c_5 - c_2) \pmod{18}$ and $2(c_1 - c_5) \equiv c_4 - c_2 \pmod{9}$. Solving the simultaneous equations $c_4 \equiv 2c_1 \pmod{9}$, $c_5 \equiv 3c_1 + 2c_2 \pmod{18}$, $c_1 + c_4 - c_5 \equiv 4(c_5 - c_2) \pmod{18}$ and $2(c_1 - c_5) \equiv c_4 - c_2 \pmod{9}$, we get $c_1 \equiv c_2 \equiv c_4 \equiv c_5 \equiv 0 \pmod{9}$, which completes the proof. \square

Combining Propositions 1 and 3 yields the following theorem.

Theorem 2: Let $q \equiv 1 \pmod{9}$ be a prime where $4q = c^2 + 27d^2$ and $c \equiv 7 \pmod{9}$. Let $q = (\sum_{i=0}^5 c_i \xi^i)(\sum_{i=0}^5 c_i \xi^{-i})$ where ξ is a primitive 9th root of unity of \mathbb{F}_q . Then, in the case of $\text{Ind } 3 \equiv 0 \pmod{3}$, the \mathbb{F}_3 -linear complexity of the ternary Sidelnikov sequence $\{s_n\}_{n \geq 0}$ of period $q - 1$ is

- 1) $LC \leq q - 4$ if $d \equiv 0 \pmod{3}$;
- 2) $LC \leq q - 10$ if $c \equiv 16 \pmod{18}$ and $d \equiv 0 \pmod{6}$, or $c \equiv 7 \pmod{18}$ and $d \equiv 3 \pmod{6}$;
- 3) $LC \leq q - 15$ if $c \equiv 16 \pmod{18}$ and $d \equiv 0 \pmod{6}$, or $c \equiv 7 \pmod{18}$ and $d \equiv 3 \pmod{6}$; and $c_4 \equiv 2c_1 \pmod{9}$ and $c_5 \equiv 3c_1 + 2c_2 \pmod{18}$;
- 4) $LC \leq q - 17$ if $c \equiv 16 \pmod{18}$ and $d \equiv 0 \pmod{6}$, or $c \equiv 7 \pmod{18}$ and $d \equiv 3 \pmod{6}$; and $c_1 \equiv c_2 \equiv c_4 \equiv c_5 \equiv 0 \pmod{9}$.

Proof: Let $q = 9f + 1$. Then $x^{q-1} - 1 \equiv (x^3 - 1)^{3f} \pmod{3}$. It is clear that $f \geq 2$. Thus, the multiplicities of the 3rd roots γ ($\gamma^3 = 1$) as roots of $x^{q-1} - 1$ are at least 6. Let $q = (\sum_{i=0}^5 c_i \xi^i)(\sum_{i=0}^5 c_i \xi^{-i})$ where ξ is a primitive 9th root of unity of \mathbb{F}_q .

- 1) It is clear from Proposition 1 3).
- 2) From Proposition 3 2), it follows by solving the simultaneous equations $c + 2 \equiv 0 \pmod{9}$, $d \equiv 0 \pmod{3}$, and $2 + c - 9d \equiv 0 \pmod{18}$.
- 3) From Proposition 1 5) and Proposition 3 3), we have $c_4 \equiv 2c_1 \pmod{9}$ and $c_5 \equiv 3c_1 + 2c_2 \pmod{18}$, which imply that $c_1 + c_5 \equiv 2(c_2 + c_4)$. Thus, the result is true.

4) It is clear from Proposition 1 5) and Proposition 3 4). \square

Example 6: Let $q = 73$ and $M = 3$. Then a ternary Sidelnikov sequence $\{s_n\}_{n \geq 0}$ of period 72 defined by $\alpha = 5$ of \mathbb{F}_{73} is represented as $\{2, 2, 1, 2, 2, 2, 1, 2, 0, 1, 0, 1, 0, 0, 0, 2, 1, 0, 1, 0, 2, 2, 2, 1, 2, 1, 2, 0, 1, 1, 2, 2, 0, 2, 1, 0, 0, 1, 0, 0, 1, 1, 2, 2, 0, 0, 0, 0, 0, 2, 1, 2, 0, 2, 1, 1, 0, 2, 1, 2, 0, 2, 2, 1, 1, 1, 1, 0, 1, 2, 2, 1\}$, and $2 = \alpha^8$ and $3 = \alpha^6$ over $\mathbb{F}_{73} = \{0, 1, 2, 3, \dots, 72\}$, i.e., $\text{Ind } 2 \equiv 2 \pmod{6}$ and $\text{Ind } 3 \equiv 0 \pmod{3}$. From $4q = c^2 + 27d^2$ and $c \equiv 1 \pmod{3}$, we have $c = 7$ and $d = -3$. Then, according to Proposition 1 3) and Theorem 2 1), 1 is a triple root of $S(x)$. In addition, although 2 is a root of $x^{72} - 1$, it is clear that 2 is not a root of $S(x)$ because $S(2) \equiv 2 \pmod{3}$, which can also be further confirmed from Proposition 2 since $73 = (-5)^2 + 3 \times 4^2$ and $4 \equiv 1 \pmod{3}$. Thus, the linear complexity of this sequence

TABLE I
THE \mathbb{F}_3 -LINEAR COMPLEXITIES OF TERNARY SIDELNIKOV SEQUENCES OF PERIOD $q - 1 = 2 \times 3^\lambda$ FOR $1 \leq \lambda \leq 20$

λ	$q = 2 \times 3^\lambda + 1$	LC	from Corollary 2	comments
1	7	5	1)	$b = d = 1$
2	19	16	2)	$b = d = 1$
3	55	-	-	no prime
4	163	160	2)	$b = 7, d = 1$
5	487	484	2)	$b = 1, d = 7$
6	1459	1454	3)	$b = 15, d = 10$ $\alpha = 3$, and $\text{Ind}_\alpha 2 \equiv 723$
7	4375	-	-	no prime
8	13123	-	-	no prime
9	39367	39362	3)	$b = 33, d = 22$ $\alpha = 3$, and $\text{Ind}_\alpha 2 \equiv 19674$
10	118099	-	-	no prime
11	354295	-	-	no prime
12	1062883	-	-	no prime
13	3188647	-	-	no prime
14	9565939	-	-	no prime
15	28697815	-	-	no prime
16	86093443	86093440	2)	$b = 1591, d = 2423$
17	258280327	$q^{1/2} - 1 \leq LC \leq q - 7$	5)	$b = 7281, d = 4854, \alpha = 5$ $\text{Ind}_\alpha 2 \equiv 159499944$, and $\text{Ind}_\alpha 3 \equiv 104564853$
18	774840979	-	-	no prime
19	2324522935	-	-	no prime
20	6973568803	-	-	no prime

Note: $\text{Ind}_\alpha x$ is the index of $x \in \mathbb{F}_q$ to the base α modulo q .

is $LC \leq 69$. In fact, we know that its linear complexity is exactly 69 since $\text{gcd}(x^{72} - 1, S(x)) = (x - 1)^3$, which means this upper bound is reachable.

Finally, for the special case of $q = 2 \times 3^\lambda + 1$ where λ is an integer, the following corollary can be obtained, and Table I lists all examples for $1 \leq \lambda \leq 20$. From this table, it is easy to see that the linear complexities of all sequences considered are extremely close to their periods.

Corollary 2: Let a prime $q = 2 \times 3^\lambda + 1$ ($\lambda \geq 1$). Let q have the decompositions $q = a^2 + 3b^2$ and $4q = c^2 + 27d^2$ where $a \equiv 1 \pmod{3}$ and $c \equiv 1$ or $7 \pmod{9}$. Let $q = (\sum_{i=0}^5 c_i \xi^i)(\sum_{i=0}^5 c_i \xi^{-i})$ where ξ is a primitive 9th root of unity of \mathbb{F}_q . Then, the \mathbb{F}_3 -linear complexity of ternary Sidelnikov sequence $\{s_n\}_{n \geq 0}$ of period $q - 1$ satisfies

- 1) $LC = q - 2$ if $q = 7$;
- 2) $LC = q - 3$ if $q \equiv 1 \pmod{9}$ and $b, d \not\equiv 0 \pmod{3}$;
- 3) $LC = q - 5$ if $q \equiv 1 \pmod{9}$, $b \equiv 0, d \not\equiv 0 \pmod{3}$, and $\text{Ind } 2 \equiv 0 \pmod{3}$;
- 4) $LC \leq q - 6$ if $q \equiv 1 \pmod{9}$, $b \equiv d \equiv 0 \pmod{3}$, and $\text{Ind } 2 \equiv \text{Ind } 3 \equiv 0 \pmod{3}$. More precisely, it follows that
 - 4-1) $LC = q - 7$ if $c_1 + c_5 \equiv 2(c_2 + c_4) \pmod{9}$,
 - 4-2) $LC = q - 9$ if $c_1 + c_5 \equiv 2(c_2 + c_4), c_4 \equiv 2c_1 \pmod{9}$,
 - 4-3) $LC = q - 10$ if $c_1 + c_5 \equiv 2(c_2 + c_4), c_4 \equiv 2c_1, c_3 + c_4 + c_5 \equiv 0 \pmod{9}$,
 - 4-4) $LC = q - 11$ if $c_1 + c_5 \equiv 2(c_2 + c_4), c_4 \equiv 2c_1, c_3 + c_4 + c_5 \equiv 0, 1 + c_0 + c_2 + 2c_4 \equiv 0 \pmod{9}$,

4-5) $LC \leq q - 12$ if $c_0 \equiv -1 + 2c_5, c_1 \equiv -c_5, c_2 \equiv 2c_5, c_3 \equiv c_5, c_4 \equiv -2c_5 \pmod{9}$;

5) $LC \leq q - 7$ if $q \equiv 1 \pmod{9}, b \equiv 0 \pmod{9}$ and $d \equiv 0 \pmod{3}$, and $\text{Ind } 2 \equiv \text{Ind } 3 \equiv 0 \pmod{3}$. More precisely, it follows that

- 5-1) $LC = q - 8$ if $c_1 + c_5 \equiv 2(c_2 + c_4) \pmod{9}$,
- 5-2) $LC = q - 10$ if $c_1 + c_5 \equiv 2(c_2 + c_4), c_4 \equiv 2c_1 \pmod{9}$,
- 5-3) $LC = q - 11$ if $c_1 + c_5 \equiv 2(c_2 + c_4), c_4 \equiv 2c_1, c_3 + c_4 + c_5 \equiv 0 \pmod{9}$,
- 5-4) $LC = q - 12$ if $c_1 + c_5 \equiv 2(c_2 + c_4), c_4 \equiv 2c_1, c_3 + c_4 + c_5 \equiv 0, 1 + c_0 + c_2 + 2c_4 \equiv 0 \pmod{9}$,
- 5-5) $LC \leq q - 13$ if $c_0 \equiv -1 + 2c_5, c_1 \equiv -c_5, c_2 \equiv 2c_5, c_3 \equiv c_5, c_4 \equiv -2c_5 \pmod{9}$.

IV. CONCLUSION

The main purpose of this paper is to give the Hasse derivative formulas to determine the multiplicity of γ , the primitive r th root of unity over \mathbb{F}_M or in an extension field of \mathbb{F}_M , as a root of $S(x)$ which is the generating function of an M -ary Sidelnikov sequence $\{s_n\}_{0 \leq n \leq q-2}$. In general, this proposed method can be used to determine the exact \mathbb{F}_M -linear complexity of the M -ary Sidelnikov sequence whenever the value of certain cyclotomic numbers and the factorization of $x^{q-1} - 1$ over \mathbb{F}_M are known. However, the well-known results on cyclotomic numbers are currently limited to the orders

$e \leq 24$. This limitation hinders our ability to calculate the multiplicity of γ if r is large. On the other hand, it is not easy to factorize the polynomial $x^{q-1} - 1$ over \mathbb{F}_M . Based on the above, it seems that the determination of the exact \mathbb{F}_M -linear complexity of the M -ary Sidelnikov sequence is a difficult problem, especially when the characteristic of the field is a factor of the period of Sidelnikov sequence [9]. Nevertheless, one may determine the \mathbb{Z}_4 -linear complexity of 4-ary Sidelnikov sequences when $q = 3 \cdot 4^\lambda + 1$ [28].

APPENDICES

APPENDIX A

THE CYCLOTOMIC NUMBERS OF ORDER 2

Let $q = ef + 1$ be a prime power. When $e = 2$, the cyclotomic numbers are given in [25] by

- (1) f even: $(0, 0)_2 = (f - 2)/2$; $(0, 1)_2 = (1, 0)_2 = (1, 1)_2 = f/2$;
- (2) f odd: $(0, 0)_2 = (1, 0)_2 = (1, 1)_2 = (f - 1)/2$; $(0, 1)_2 = (f + 1)/2$.

APPENDIX B

THE CYCLOTOMIC NUMBERS OF ORDER $2r$

Let $q = p^m \equiv 1 \pmod{2r}$ for any prime p such that $m = uv$, $u \geq 1$, v is the order of p modulo r and v is even. Then, the cyclotomic numbers of order $2r$ over \mathbb{F}_q are given in [27] by:

- (1) $(0, j)_{2r} = (j, 0)_{2r} = (j, j)_{2r}$;
- (2) $4r^2A := 4r^2(0, 0)_{2r} = q - 6r + 1 - (4r^2 - 6r + 2)(-1)^u q^{1/2}$;
- (3) $4r^2B := 4r^2(0, j)_{2r} = q - 2r + 1 + 2(r - 1)(-1)^u q^{1/2}$ for $j \not\equiv 0 \pmod{2r}$;
- (4) $4r^2C := 4r^2(i, j)_{2r} = q + 1 - 2(-1)^u q^{1/2}$ for $i, j, i - j \not\equiv 0 \pmod{2r}$.

APPENDIX C

THE CYCLOTOMIC NUMBERS OF ORDER 3

Let $q = ef + 1$ be a prime. When $e = 3$ and $4q = c^2 + 27d^2$ with $c \equiv 1 \pmod{3}$, the cyclotomic numbers are given in [25] by

- (1) $A := (0, 0)_3 = (q - 8 + c)/9$;
- (2) $B := (0, 1)_3 = (1, 0)_3 = (2, 2)_3 = (2q - 4 - c - 9d)/18$;
- (3) $C := (0, 2)_3 = (1, 1)_3 = (2, 0)_3 = (2q - 4 - c + 9d)/18$;
- (4) $D := (1, 2)_3 = (2, 1)_3 = (q + 1 + c)/9$.

APPENDIX D

THE CYCLOTOMIC NUMBERS OF ORDER 6

Let $q = ef + 1$ be a prime. When $e = 6$ and $q = a^2 + 3b^2$ with $a \equiv 1 \pmod{3}$, the cyclotomic numbers $(u, v)_6$ when f is even are given in Table II [29] by

- I. Case $\text{Ind } 2 \equiv 0 \pmod{6}$
 - (1) $36A := (0, 0)_6 = q - 17 - 20a$,
 - (2) $36B := (0, 1)_6 = q - 5 + 4a + 18b$,
 - (3) $36C := (0, 2)_6 = q - 5 + 4a + 6b$,
 - (4) $36D := (0, 3)_6 = q - 5 + 4a$,

TABLE II
CYCLOTOMIC NUMBERS OF ORDER 6 AND f EVEN

u	v					
	0	1	2	3	4	5
0	A	B	C	D	E	F
1	B	F	G	H	I	G
2	C	G	E	I	J	H
3	D	H	I	D	H	I
4	E	I	J	H	C	G
5	F	G	H	I	G	B

- (5) $36E := (0, 4)_6 = q - 5 + 4a - 6b$,
- (6) $36F := (0, 5)_6 = q - 5 + 4a - 18b$,
- (7) $36G := (1, 2)_6 = q + 1 - 2a$,
- (8) $36H := (1, 3)_6 = q + 1 - 2a$,
- (9) $36I := (1, 4)_6 = q + 1 - 2a$,
- (10) $36J := (2, 4)_6 = q + 1 - 2a$;

II. Case $\text{Ind } 2 \equiv 2 \pmod{6}$

- (1) $36A := (0, 0)_6 = q - 17 - 8a - 6b$,
- (2) $36B := (0, 1)_6 = q - 5 + 4a + 6b$,
- (3) $36C := (0, 2)_6 = q - 5 - 8a$,
- (4) $36D := (0, 3)_6 = q - 5 + 4a + 6b$,
- (5) $36E := (0, 4)_6 = q - 5 + 4a + 6b$,
- (6) $36F := (0, 5)_6 = q - 5 + 4a - 12b$,
- (7) $36G := (1, 2)_6 = q + 1 - 2a + 6b$,
- (8) $36H := (1, 3)_6 = q + 1 - 2a - 12b$,
- (9) $36I := (1, 4)_6 = q + 1 - 2a + 6b$,
- (10) $36J := (2, 4)_6 = q + 1 + 10a - 6b$;

III. Case $\text{Ind } 2 \equiv 1 \pmod{6}$

- (1) $36A := (0, 0)_6 = q - 17 - 8a + 6b$,
- (2) $36B := (0, 1)_6 = q - 5 + 4a + 12b$,
- (3) $36C := (0, 2)_6 = q - 5 + 4a - 6b$,
- (4) $36D := (0, 3)_6 = q - 5 + 4a - 6b$,
- (5) $36E := (0, 4)_6 = q - 5 - 8a$,
- (6) $36F := (0, 5)_6 = q - 5 + 4a - 6b$,
- (7) $36G := (1, 2)_6 = q + 1 - 2a - 6b$,
- (8) $36H := (1, 3)_6 = q + 1 - 2a - 6b$,
- (9) $36I := (1, 4)_6 = q + 1 - 2a + 12b$,
- (10) $36J := (2, 4)_6 = q + 1 + 10a + 6b$.

APPENDIX E

THE CYCLOTOMIC NUMBERS OF ORDER 9

Let $q = ef + 1$ be a prime. When $e = 9$ and $4q = c^2 + 27d^2$ with $c \equiv 7 \pmod{9}$, the cyclotomic numbers $(u, v)_9$ are given in Table III [21]. Each cyclotomic number is expressed as a constant plus a linear combination of $q, c, d, c_0, c_1, c_2, c_3, c_4$, and c_5 where

$$q = \left(\sum_{i=0}^5 c_i \xi^i \right) \left(\sum_{i=0}^5 c_i \xi^{-i} \right)$$

is a factorization of q in the field of 9th roots of unity, and ξ is a primitive 9th root of unity. The following cyclotomic numbers are only for $\text{Ind } 3 \equiv 0 \pmod{3}$.

TABLE III
CYCLOTOMIC NUMBERS OF ORDER 9

u	v								
	0	1	2	3	4	5	6	7	8
0	A	B	C	D	E	F	G	H	I
1	B	I	J	K	M	N	O	P	J
2	C	J	H	P	Q	R	S	Q	K
3	D	K	P	G	O	S	T	R	M
4	E	M	Q	O	F	N	R	S	N
5	F	N	R	S	N	E	M	Q	O
6	G	O	S	T	R	M	D	K	P
7	H	P	Q	R	S	Q	K	C	J
8	I	J	K	M	N	O	P	J	B

- (1) $162A := 2q - 52 + 2c + 108c_0 - 54c_3.$
- (2) $162B := 2q - 16 - c + 9d - 12c_0 + 42c_1 - 12c_2 + 24c_3 - 30c_4 + 24c_5.$
- (3) $162C := 2q - 16 - c - 9d - 12c_0 + 24c_1 + 42c_2 - 12c_3 - 12c_4 - 12c_5.$
- (4) $162D := 2q - 16 + 2c - 18c_0 + 36c_3.$
- (5) $162E := 2q - 16 - c + 9d - 12c_0 - 12c_1 + 24c_2 + 24c_3 + 42c_4 - 12c_5.$
- (6) $162F := 2q - 16 - c - 9d - 12c_0 - 12c_1 - 30c_2 - 12c_3 - 12c_4 + 42c_5.$
- (7) $162G := 2q - 16 + 2c - 18c_0 - 18c_3.$
- (8) $162H := 2q - 16 - c + 9d - 12c_0 - 30c_1 - 12c_2 + 24c_3 - 12c_4 - 12c_5.$
- (9) $162I := 2q - 16 - c - 9d - 12c_0 - 12c_1 - 12c_2 - 12c_3 + 24c_4 - 30c_5.$
- (10) $162J := 2q + 2 + 2c - 18c_1 + 18c_2.$
- (11) $162K := 2q + 2 - c + 9d + 6c_0 + 6c_1 - 12c_2 - 12c_3 + 6c_4 + 6c_5.$
- (12) $162M := 2q + 2 - c - 9d + 6c_0 - 12c_1 + 6c_2 + 6c_3 + 6c_4 + 6c_5.$
- (13) $162N := 2q + 2 + 2c + 18c_1 - 18c_4 - 18c_5.$
- (14) $162O := 2q + 2 - c + 9d + 6c_0 - 12c_1 + 6c_2 - 12c_3 + 6c_4 + 6c_5.$
- (15) $162P := 2q + 2 - c - 9d + 6c_0 + 6c_1 - 12c_2 + 6c_3 + 6c_4 + 6c_5.$
- (16) $162Q := 2q + 2 + 2c - 18c_2 + 18c_4 + 18c_5.$
- (17) $162R := 2q + 2 - c + 9d + 6c_0 + 6c_1 + 6c_2 - 12c_3 - 12c_4 - 12c_5.$
- (18) $162S := 2q + 2 - c - 9d + 6c_0 + 6c_1 + 6c_2 + 6c_3 - 12c_4 - 12c_5.$
- (19) $162T := 2q + 2 + 2c.$

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